

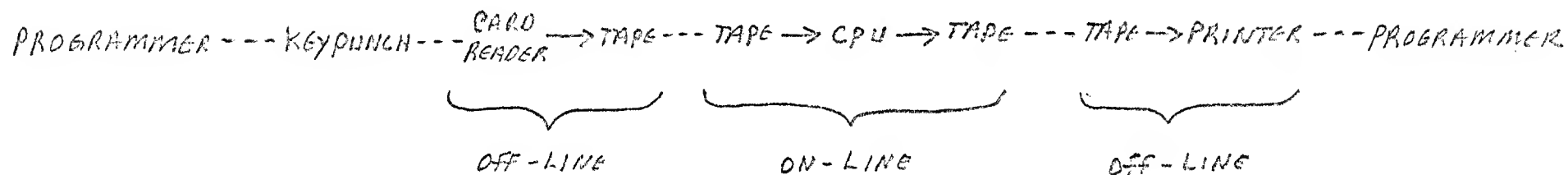
6000 SCOPE 3 HANDOUTS

<u>SECTION</u>	<u>TITLE</u>
I.	INTRODUCTION TO 6000 SYSTEMS
II.	GENERAL 6000 SCOPE 3 CHARACTERISTICS
III.	DEAD START LOAD AND PRELOADER (PLR)
IV.	POOL PROCESSOR RESIDENT (PPR)
V.	MONITOR (MTR)
VI.	DYNAMIC SYSTEM DISPLAY (DSD)
VII.	STACK PROCESSOR (1SP/1SQ)
VIII.	JOB FLOW
IX.	SCOPE LOADER
X.	CIRCULAR INPUT/OUTPUT (CIO)
XI.	CENTRAL PROGRAM CONTROL (CPC)
XII.	CHECKPOINT/RESTART
XIII.	JOB DISPLAY
XIV.	SYSTEM AND JOB MAINTENANCE

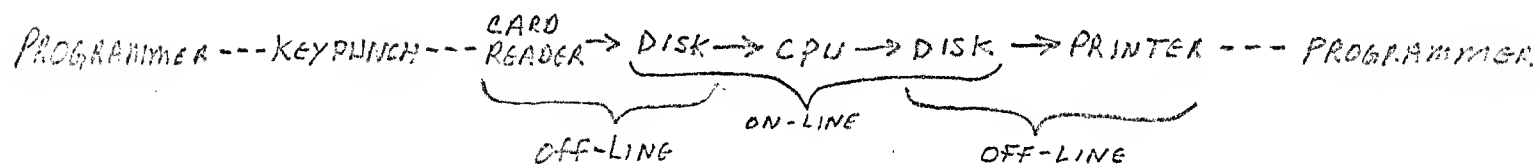
INTRODUCTION

this document is intended primarily as an instructional handout and is intended to be used in the 6000 Scope 3.X Operating System courses. It is not intended to replace the 6000 Scope 3.1 Reference Manual or the 6000 Scope 3.1 I. M. S. manual. Rather it is to be used with them.

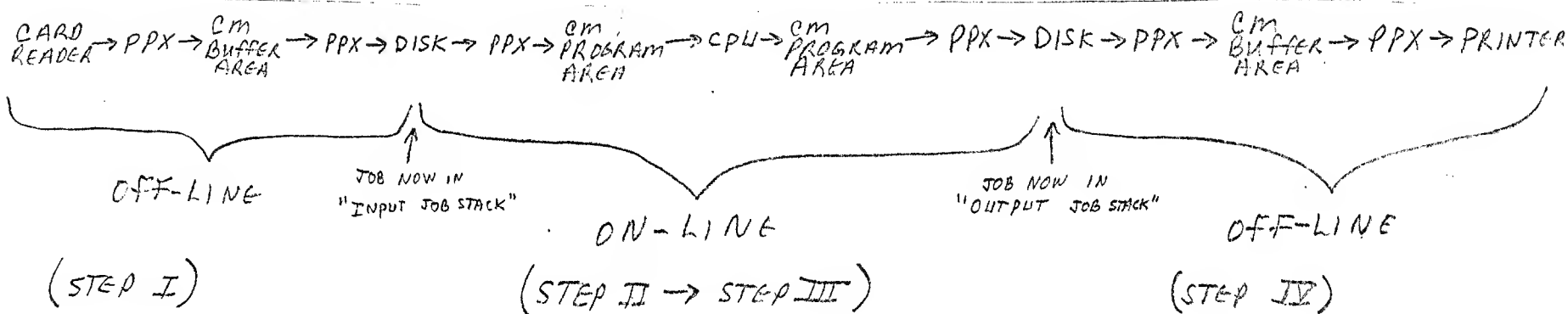
NON-6000 SYSTEM PROGRAM FLOW



6000 SYSTEM PROGRAM FLOW



⏟
(DETAILS)



PROGRAM FLOW CHART

PROGRAM FLOW STEPSSTEP I - LOAD JOB

- A. CARD READER TO DISK OPERATION.
- B. RECORD JOB NAME, PRIORITY, FL, AND WHERE PLACED ON DISK. IN CMR.
- C. READ NEXT JOB (RETURN TO A)

STEP II - BEGIN JOB

- A. SEARCH LIST OF JOBS PLACED ON DISK FOR THE HIGHEST PRIORITY JOB
- B. IF FOUND, LOAD CONTROL CARDS TO CONTROL POINT AREA IN CMR.
- C. TRANSLATE JOB CARD AND SET UP CONTROL POINT AREA.
- D. SET UP CONDITIONS WHICH WILL INDICATE TO MTR TO INITIATE STEP III

STEP III - ADVANCE JOB

- A. TRANSLATE NEXT CONTROL CARD AND CAUSE EXECUTION OF IT.
- B. WHEN EXECUTION FINISHED, MTR CAUSES A RETURN TO "A".
- C. LAST CARD EXECUTED, UNNEEDED FILES ARE DROPPED, EQUIPMENT IS DROPPED, JOB NAME AND PRIORITY ASSIGNED TO THE OUTPUT FILE, CLEAR OUT THE CONTROL POINT AREA.
- D. CALL FOR STEP II'S SYSTEM PROGRAM.
(RETURN TO STEP II)

STEP IV - DUMP JOB

- A. SEARCH FOR A JOB READY TO BE DUMPED
- B. IF ONE FOUND, ASSIGN TO CONTROL POINTS "BUFFER POINT" AND DUMP THE JOB
- C. DUMP THE DAYFILE FOR THE JOB ON THE PRINTER.
- D. DROP THE FILE.

I. 2 FILES MAKE UP THE OPERATING SYSTEM:

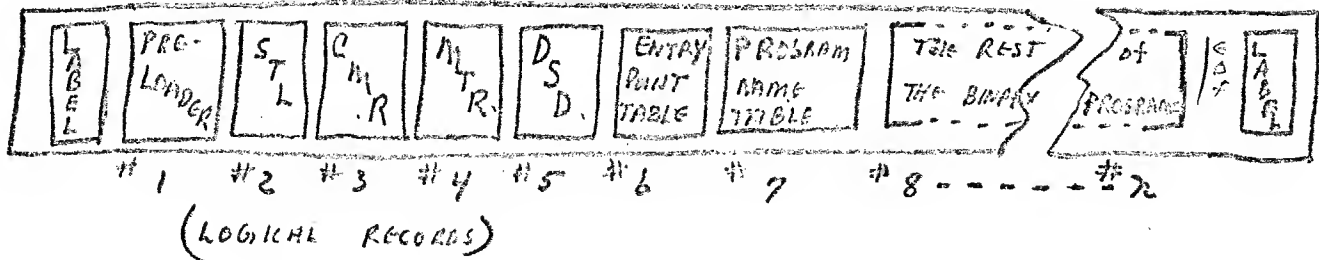
A. 1st FILE IS THE BINARY DEAD START PORTION

1. 2 VERSIONS

a. 6603 DISK FILE

b. 6638 DISK FILE

2. FORMAT OF BINARY FILE -



B. 2nd FILE IS PROGRAM LIBRARY, "PLDPL".

1. CONSISTING OF;

a. 1st RECORD HAS;

1) 1st WORD CONTAINS 2 COUNTERS; COUNT OF THE IDENTIFIERS IN THE IDENTIFIER SECTION (BITS 35-18) AND THE COUNT OF THE DECK NAMES (17-00).

2) A DIRECTORY WHICH CONTAINS A LIST OF ALL IDENTIFIERS USED. ("CORRECTION IDENTIFIERS")

3) A DECK LIST WHICH CONTAINS A LIST OF ALL DECKS WHICH OCCUR OR HAVE OCCURRED ON THE FILE. ("DECK NAMES CURRENTLY KNOWN")

4) THE TEXT STREAM WHICH CONTAINS THE CARD IMAGES AND CONTROL INFORMATION, KNOWN AS "CORRECTION HISTORY BYTES".

b. THE BINARY DECKS OF THE PRODUCT SET, EACH CONSTITUTING A RECORD.

...

ENDING IN E.O.F.

PROGRAM LIBRARY CONTENTS

COMMON DECKS

COMFILE	the deck containing definitions of system symbols
FNTSRCH	a routine to search the FNT for a specified file name
READPP	a routine to read a record directly into PP memory
BUFINIT	a routine to initialize PP memory for a call to READPP
RELOC	a set of macros which, if called, will cause a PP routine to be self-relocating
PPTIME	a routine to enter PP time into the system dayfile
COMCOMP	a micro which defines the COMPASS version number
TRNCOM	a set of declarations for use by the translator subroutines

TEXT DECKS

TAPEDS	program to generate the tape dead-start card
DISCDS	program to generate the disc dead-start card
DUMPDS	program to generate the dead-start dump cards
RECOV	program to generate the recovery dead-start card
PREPP	PP portion of the system pre-loader
PRECP	CP portion of the system pre-loader
STLPP	PP portion of the system dead-start loader
STLCP	CP portion of the system dead-start loader
CMR	central memory resident
MTR	system monitor
DSD	system display
CIO	central input/output package
LDR	relocatable loader
LOD	relocatable loader
MSG	routine to enter messages into the dayfile

RFL	routine to request field length
1AJ	routine to advance control point
1BJ	routine to bring a job to a control point
1LJ	routine to load a job from the card reader
1OT	routine to control printing and punching of output
1SP	6603 stack processor
2BP	routine to check buffer parameters (FET)
2EF	routine to process error flag
2ES	routine to enter a stack request
2RC	routine to read cards
2TR	routine to read 3000 magnetic tapes
2TS	routine to translate control card statements
2TW	routine to write 3000 magnetic tapes
ATS	STITCH routine
CHK	routine to check output file for RUN compiler
CKP	routine to take checkpoint dumps
CLL	STITCH routine
CLO	routine to close a file
CLS	STITCH routine
CPL	STITCH routine
CTS	routine to process COMMON function
DIS	job display
DMP	routine to dump central memory
EXU	STITCH routine

LBC	routine to load binary cards from INPUT
LOC	routine to load octal correctors from INPUT
MDI	routine to move system directory (EDITLIB)
MEM	routine to process MEMORY function
OPE	routine to open a file
PBC	routine to punch binary cards from central memory
RBR	routine to read a binary record
REQ	routine to process a REQUEST card or function
RST	routine to restart a job from a checkpoint file
SRB	routine to enter an expanded disc address into the directory (EDITLIB)
TIM	routine to process a TIME, DATE, CLOCK, or JDATE function
WBR	routine to write a binary record
1BT	routine to initialize a blank tape
1CO	routine to complete OUTPUT file
1CY	routine to copy a file to the checkpoint file
1DF	routine to dump the dayfile
1LT	routine to load a job from magnetic tape
1MF	routine to open a tape file (multi-file tape)
1MR	routine to open a tape file (multi-reel tape)
1MW	routine to open a tape file
1OD	routine to open a file
1PL	routine to process plotter output (dummy)
1PO	routine to process punch output
1RI	routine to process the ROLLIN type-in

1RO	routine to process the ROLLOUT type-in
1SX	routine to enter stack processor error messages into the dayfile
1TD	routine to dump output files to tape
2CA	checkpoint abort procedure
2CF	routine to close files in preparation for the output queue
2CJ	routine to append the control point dayfile to the primary output file
2DF	routine to drop a file from the system
2LA	relocatable loader overlay
2LB	relocatable loader overlay
2LE	relocatable loader overlay to process error conditions
2LP	routine to process on-line print output
2PC	routine to process punch output
2RT	routine to read 6000 magnetic tapes
2TB	routine to handle backward motion on 3000 magnetic tapes
2TF	routine to handle forward motion on 3000 magnetic tapes
2TJ	routine to translate job cards
2WT	routine to write 6000 magnetic tapes
3OT	routine to process print output from the OUTPUT package
4LB	routine to handle label-processing
7TP	routine to handle write parity errors on 3000 magnetic tapes

LOADER	central memory portion of the relocatable loader
OVERLOD	abbreviated central memory loader for loading overlays
QXX	STITCH routine
EDITLIB	program for producing and editing system libraries
EDITSYM	program for producing and editing program libraries
COPYBCD	program for writing discrete print line images on magnetic tape
CPC	program for communicating requests from central memory to the system
IORANDM	routine for processing SCOPE style indexes
IO	routine to handle blocking/deblocking
COMPARE	routine to compare contents of two files
RESTART	routine to initiate the restarting of a job from a checkpoint file
BKSP	routine to backspace a file one or more logical records
COPY	routine to create a job stack on magnetic tape
COPYBF	routine to copy one or more binary files
COPYN	routine to copy selected records or files
COPYSBF	routine to format a packed display code binary file for printing
REWIND	routine to rewind a file
SYSTEXT	source deck of the system macros

DEFUNCT PP OVERLAYS

<u>Name</u>	<u>Replacement</u>
HLP	EDITLIB includes HLP's function
PBS	Punching in I-mode not necessary
SOS	EDITLIB includes SOS's function
1DS	Function included in DIS
1RF	Dead-start recovery card
2BD	Stack processor
2DT	Stack processor/central memory tables
2RD	Stack processor
2WD	Stack processor
3SD	Dayfile search is not necessary
4SD	Dayfile search is not necessary
7DP	Stack processor

COMFILE

The first of the common decks contained in the SCOPE program library is COMFILE. The purpose of COMFILE is to gather all system parameters together in one centralized location so that they may be easily examined and/or altered. Each parameter is referenced by a system symbol which consists of three parts:

- an identifier of one or two characters denoting the category to which the symbol belongs;
- "." a period following the identifier to indicate that this symbol is part of COMFILE;
- a mnemonic of 1 to 6 characters suggesting the meaning of the symbol.

The definition of symbols of the above form should be avoided when COMFILE is to be called.

The system symbols in COMFILE reference the following categories of parameters:

- installation parameters
- system table lengths
- system locations, words, and bytes
- pointer words
- PP resident entry points
- monitor functions
- quantities of system elements (number of tables, devices, etc.)
- PP direct locations

The set of system symbols is easily expandable as new symbols are defined.

The set is callable from each system routine by use of a single card:

*CALL,COMFILE

When COMFILE is called, a listing of COMFILE may be obtained by defining the symbol called LISTCOMF before the call to COMFILE.

SYSTEM SYMBOL IDENTIFIER DEFINITIONS

- C. Most C.x symbols represent 12-bit byte positions within central memory words, where bytes are numbered from left to right as 0 through 4.
C.x symbols are also used to represent first word addresses of PP overlays.
- CP. CP.x symbols represent words or start of programs in the CP resident area.
- D. The D.x symbols represent PP direct locations (low core); they are equated to 6-bit values from 00 through 77₈.
- IP. All IP.x symbols represent installation parameters.
- L. The L.x symbols represent table lengths or lengths of miscellaneous quantities.
- LE. The LE.x symbols represent the lengths of entries within tables.
- M. M.x symbols are equated to the values representing monitor functions.
- F. F.x symbols are equated to error flag values.

- CH. CH.x symbols represent pseudo-channel assignments; e.g., CH.FNT is the File Name Table channel.
- N. N.x symbols represent quantities of things; e.g., N.DEVICE is equated to the number of allocatable devices within the system.
- Ø. Ø.x symbols represent stack processor orders (commands). It should be noted that these orders do not correspond to values used in the code and status FET field: stack processor orders are designed for ease of use by the stack processor.
- ØV. ØV.x symbols represent 3-character display code PP overlay names. There is one such symbol for each overlay; the mnemonic (x) is, in all cases, the 3-character overlay name. The appropriate ØV.x should be referenced whenever an overlay is to be loaded in order to force an entry in the cross reference table.
- P. P.x symbols represent locations of central memory pointer words.
- R. R.x symbols represent PP resident entry points.
- S. S.x symbols represent the right offset of a field within a PP word, i.e. the number of bit positions which must be right shifted to right-justify the field to bit 0.

T. The T.x symbols are equated to the first word addresses of central memory tables. In general, these addresses should be obtained from the contents of the pointer words rather than direct use of T.x symbols.

W.x The W.x symbols are equated to values representing the relative positions of central memory words within tables. For example, assume that the address of a control point area is contained in the PP A-register; then to obtain the word containing the job name, the following code should be written

ADN W.CPJNAM

CRD D.TO

* 64/65/6600 SCOPE 3.0			COMFILE00001
*****			COMFILE00002
W,CPRES2	EQU	160B	COMFILE00003
*			COMFILE00004
* SCOPE VERSION 3.0 COMMON SYMBOL DEFINITION			COMFILE00005
* IN ORDER TO LIST THE CONTENTS OF OF COMFILE, THE			COMFILE00006
* SYMBCL #LISTCOMF# SHOULD BE DEFINED PRIOR TO CALLING COMFILE			COMFILE00007
*			COMFILE00008
*****			COMFILE00009
IF =DEF,LISTCOMF,1			COMFILE00010
LIST =L			COMFILE00011
EJECT			COMFILE00012
PCPADR	EQU	22000B	COMFILE00013
LCPADR	EQU	42000B	COMFILE00014
T,CPZ	EQU	20B	COMFILE00015
T,MON	EQU	21B	COMFILE00016
T,STO	EQU	22B	COMFILE00017
T,CIDLE	EQU	23B	COMFILE00018
T,PIDLE	EQU	24B	COMFILE00019
T,PPR	EQU	25B	COMFILE00020
T,CLK	EQU	30B	COMFILE00021
T,SLAB1	EQU	31B	COMFILE00022
T,DATE	EQU	31B	COMFILE00023
T,JDATE	EQU	27B	COMFILE00024
T,SLAB2	EQU	32B	COMFILE00025
T,SLAB3	EQU	33B	COMFILE00026
T,SLAB4	EQU	34B	COMFILE00027
T,SLAB5	EQU	35B	COMFILE00028
T,SLAB6	EQU	36B	COMFILE00029
T,MSP	EQU	37B	COMFILE00030
T,MSC	EQU	40B	COMFILE00031
T,CPS	EQU	40B	COMFILE00032
T,PPS0	EQU	41B	COMFILE00033
T,PPS1	EQU	42B	COMFILE00034
T,PPS2	EQU	43B	COMFILE00035
T,PPS3	EQU	44B	COMFILE00036
T,PPS4	EQU	45B	COMFILE00037
T,PPS5	EQU	46B	COMFILE00038
T,PPS6	EQU	47B	COMFILE00039
T,PPS7	EQU	50B	COMFILE00040
T,PPS8	EQU	51B	COMFILE00041
T,PPS9	EQU	52B	COMFILE00042
T,TMP	EQU	55B	COMFILE00043
T,CPT1	EQU	56B	COMFILE00044

LE,DFB3	EQU	100B		COMFILE00565
LE,DFB4	EQU	100B		COMFILE00566
LE,DFB5	EQU	100B		COMFILE00567
LE,DFB6	EQU	100B		COMFILE00568
LE,DFB7	EQU	40B		COMFILE00569
L,DFB	EQU	LE,DFB0+LE,DFB1+LE,DFB2+LE,DFB3+LE,DFB4+LE,DFB5+LE,DFB6+LE		COMFILE00570
,,DFB7+100B				
T,LIB	EQU	T,DFB+L,DFB/10B+10B+10B	11/1	COMFILE00571
C,DIRPTR	EQU	0		COMFILE00572
C,DIRRBA	EQU	2		COMFILE00573
C,DIRRBN	EQU	3		COMFILE00574
C,DIRPRU	EQU	4		COMFILE00575
C,DIRCMA	EQU	1		COMFILE00576
C,DIRUNT	EQU	1	10/25	COMFILE00577
S,DIRPT	EQU	4	10/25	COMFILE00578
S,DIRPR	EQU	8		COMFILE00579
C,RBTWPL	EQU	0		COMFILE00580
C,RATRBR	EQU	1	11/16	COMFILE00581
S,RATRBR	EQU	3	11/16	COMFILE00582
C,RATFB	EQU	1	11/16	COMFILE00583
S,RATRND	EQU	6	11/16	COMFILE00584
S,RBTREL	EQU	7		COMFILE00585
C,RBTAL	EQU	2		COMFILE00586
C,RBTPRU	EQU	3	11/16	COMFILE00587
C,RBTLRB	EQU	4	11/16	COMFILE00588
C,RBTLPR	EQU	0	11/16	COMFILE00589
EJECT			11/16	COMFILE00590
LIST	L			COMFILE00591
EJECT				COMFILE00592
				COMFILE00593

The organization of the dead-start loader is as follows:

The first two records on the system tape are concerned with dead-starting.

The first one is PLR (called preloader) and has the job of copying the tape to disc (or some allocatable device). The second is STL which actually loads the system from the allocatable device.

The programs are subdivided in the following way:

PLR The two major divisions are the portion that runs in the peripheral processors which is called PLRPP and the part that is executed by the central processor which is called PLRCP.

PLRPP

This is divided up into a tape-read program which reads central memory, and a disc-write program which reads the information from central memory and writes it on the disc. At the end of the write program is a set-up program which is written on the disc at the end of preloading and is used to bootstrap in STL from disc.

PLRCP

-- This program accepts the information from the read program, checksums it, and arranges it in the proper format for the write program.

STL This record contains the PP resident and the recovery package as well as the programs to load. There are also two of these, STLPP and STLCP.

STLPP

This reads the information from disc and transfers it to central memory as well as sending MTR to PPO, DSD to PP9, and the resident to each of the other PP's.

STLCP

This receives the information from the STLPP program, checksums it, compares the checksum, and arranges all programs that are central memory resident in their proper place in central memory. It also inserts the addresses of all library programs into the directory.

Depending on the channel of the card reader, three different settings of the dead-start panel are necessary. One setting is used by people with card reader on channel 0 (if any such exist), another must be used for card reader on channels 1-11B, and the third for channels 12B or 13B.

The setting for channel 0 is:

<u>Word on</u> <u>Panel</u>	<u>Setting</u>	
0001	7700	
0002	E000	select card reader
0003	7700	
0004	1400	select input for EOR
0005	7400	activate channel
0006	7100	input on channel into location zero
0007	0000	

Words 0010 through 0014 are irrelevant.

Note: E is the equipment number of the card reader.

The setting for channels 12B or 13B is:

<u>Word on</u> <u>Panel</u>	<u>Setting</u>	
0001	75xx	
0002	77xx	
0003	E000	select card reader
0004	77xx	
0005	1400	select input to EOR
0006	74xx	activate channel
0007	71xx	input on channel into location zero
0010	0000	

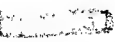
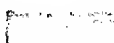
Words 0011 through 0014 are irrelevant.

Note: xx = the card reader channel
E = the equipment number of the card reader

The setting for channel 1-11B is:

<u>Word on Panel</u>	<u>Setting</u>	
0001	1410	load A with 10B
0002	73xx	send part of panel out to another PP
0003	0006	
0004	75xx	disconnect channel
0005	7112	leave this PP on channel 12
0006	0000	
0007	77xx	
0010	E000	select card reader
0011	77xx	
0012	1400	select input to EOR
0013	74xx	activate channel
0014	71xx	input into location zero

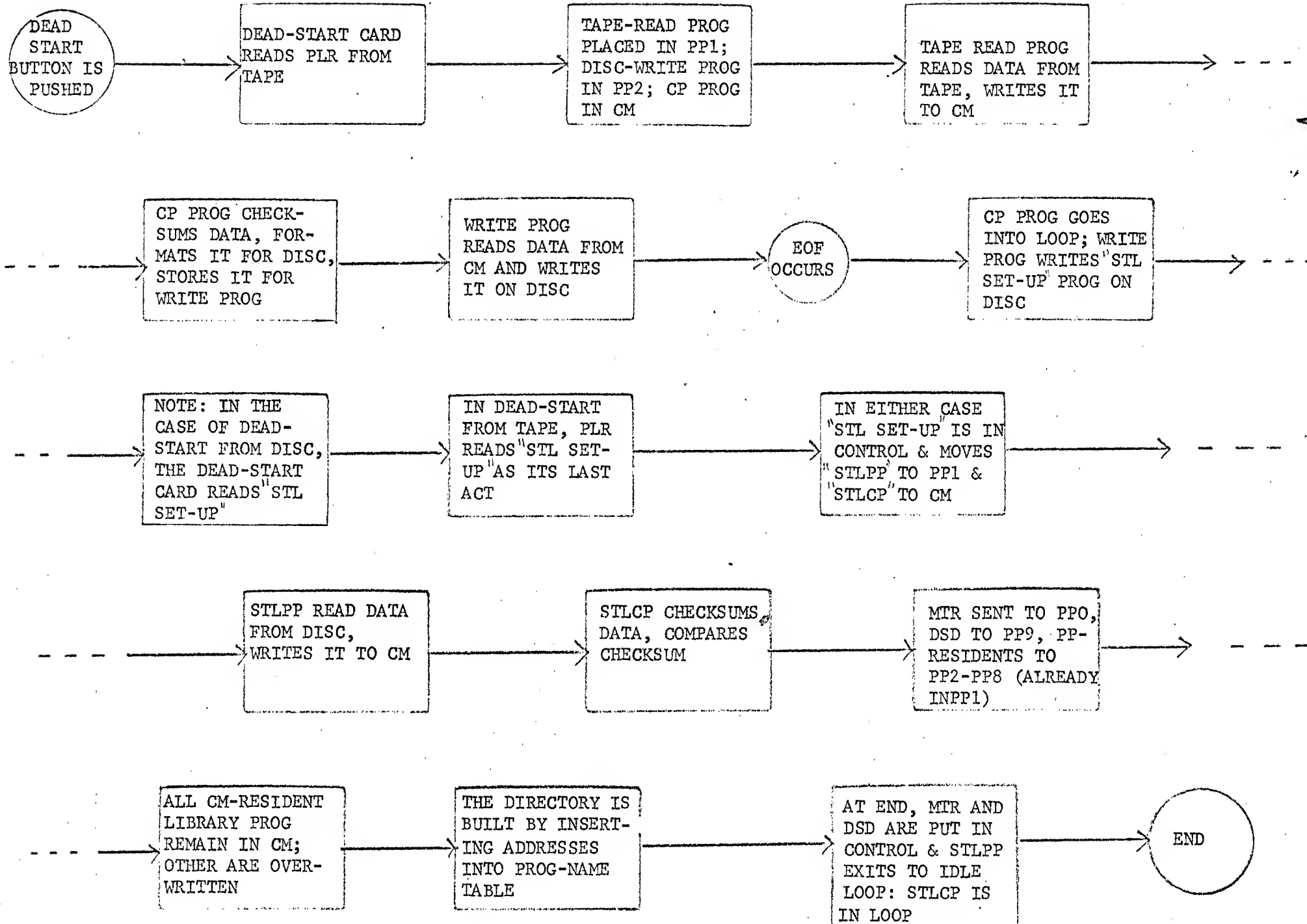
Note: The dead-start process adds a word of zeros at the end of the panel so this is why the last instruction will read into location zero.

A  setting for dead-starting directly from tape without the use of cards is also provided. 

<u>Word on Panel</u>	<u>Setting</u>	
0001	75xx	disconnect channel
0002	77xx	
0003	E00u	select tape unit
0004	77xx	
0005	0010	rewind tape
0006	77xx	
0007	1400	select input to EOR
0010	74xx	activate channel
0011	71xx	
0012	0073	read record into location 73B
0013	0112	jump to beginning of record + 7
0014	0007	

Note: This can only be used if the tapes are on channels 12B or 13B.

xx = the channel of the tape
 E = the equipment number of the tape
 u = the unit of the tape



SUMMARY OF CM RESIDENT AREAS

The Central Memory Resident Pointer Area contains pointers to larger tables contained elsewhere in CM, small tables, and various flags.

The PP Communication Area contains ten 8-word areas, one for each PP, through which PP's communicate with each other and Monitor. The Monitor communication area (T.PPC10) is not used.

The Control Point Area contains seven 200₈-word areas, one for each control point. This area is used to contain the exchange package, job name, and information about the job which is running at that control point.

The CP Resident area contains two programs which run at control point zero (the storage move program, and the idle package); their exchange packages are also kept in the CP Resident area.

The Equipment Status Table contains one entry for each device (allocatable or non-allocatable) attached to the system. Non-allocatable devices are those which may be assigned to a single control point, e.g. magnetic tape unit; allocatable devices may be used by many control points simultaneously.

The File Name Table is composed of as many 3-word entries as there are files in the system. An FNT entry is set up at the time that a file is created; it is not accessible to the user. This table provides a linkage between the user program and all required I/O tables and functions.

The Record Block Reservation area contains at least one Record Block Reservation table (RBR) for each allocatable device attached to the system. Each table contains a series of bits which denote whether or not the record block which it represents is assigned.

The Request Stack area is actually composed of two tables: The Device Status Table (DST) and the request stack itself. The DST contains one 2-word entry for each allocatable device within the system; it is static and its contents are defined at assembly time. The request stack contains entries which are requests for data transfers, device positioning, or logical operations on a file. Each entry is two words in length. The table grows from high memory to low.

The Catalogue and Security Password Index are currently of zero length; they are reserved for future use.

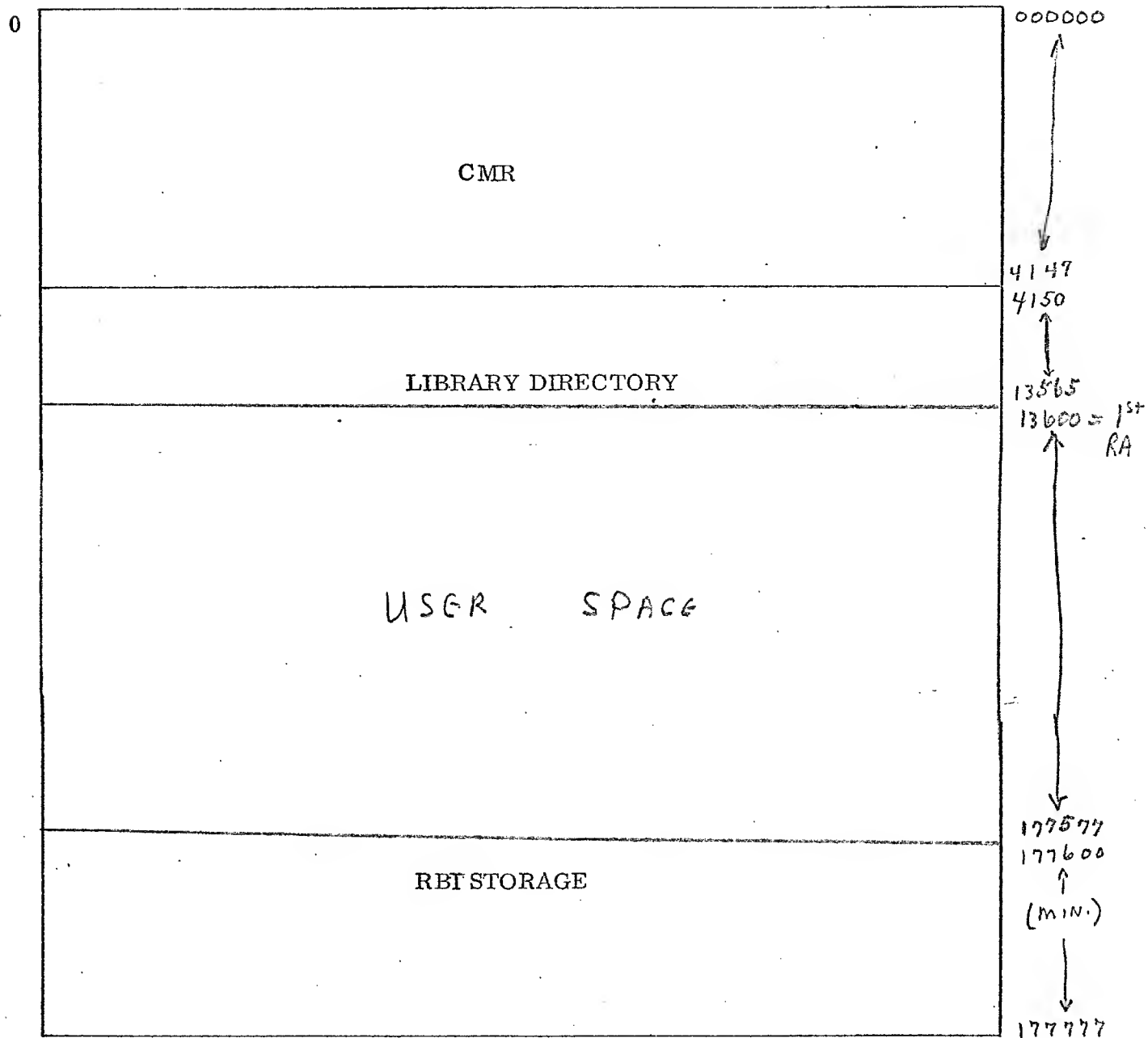
The Installation Area is currently of zero length; it is reserved for installation use.

The Dayfile Buffer area contains eight File Environment Tables and eight buffers, one for the system dayfile and one for each of the seven control points.

The library directory is composed of three sections: The entry point table, containing 1-word entries; the program name table, containing 2-word entries; and the bodies of the CM resident programs. The directory may be expanded or contracted as programs are added or deleted or as program residence is changed.

The Record Block Table area (RBT) is a collection of individual file chains, one for each file on an allocatable device currently recognized by the system. When a file is initiated, a single two-word RBT entry is assigned to that file; additional entries are assigned as needed. Each entry is divided into ten 12-bit bytes, some of which are used as pointers to additional entries, other tables, etc. The remaining bytes each contain the number of a particular record block assigned to the file in the physical order of their assignment. A record block number in an RBT is the ordinal of the bit in an RBR which represents that record block. The RBT empty chain is a pool from which words may be extracted to construct file chains and to which words are returned when the chains are discarded. The RBT expands and contracts by 100g-word blocks as files are created and released.

Overview Central Memory Allocation (65K)



CMR

0	POINTERS
60	PP COMMUNICATION AREA
200	CONTROL POINT AREA
2000	CP RESIDENT
2040	EQUIPMENT STATUS TABLE (EST)
2140	FILE NAME/STATUS TABLE (FNT/FST)
3050	RECORD BLOCK RESERVATION TABLE (RBR)
3120	REQUEST STACK
3430	CATALOGUE*
3430	SECURITY PASSWORD INDEX*
3430	INSTALLATION AREA (INS)
3430	DAYFILE BUFFER (DFB)
4440	

* ZERO LENGTH -- TO BE USED IN FUTURE DEVELOPMENT

CENTRAL MEMORY RESIDENT

P.ZERO	ZEROS					0
P.LIB	C.DIRFWA	LWA+1 of directory		dead-start load flag		1
P.RBR/P.RBT	C.RBRAD	FWA of directory	FWA/2 of empty chain	length/100B of RBT	(LWA+1)/100B of memory	2
P.DFB	FWA/10 of dayfile					3
P.FNT	FWA of FNT	LWA+1 of FNT				4
P.EST	FWA of EST	LWA+1 of EST				5
						6
P.INS	installation use					7
						10
P.CAT	future development					11
P.SPI	future development					12
P.RQS	C.RQSCS stack entry word pair ct.		C.RQSFS FWA/2 of request stack	number of devices	FWA/10B of DSTs	13
T.CST	Channel 0	Channel 4	Channel 10	FST Channel		14
P.CST2	Channel 1	Channel 5	Channel 11	FNT Channel		15
P.CST3	Channel 2	Channel 6	Channel 12	LIB CHANNEL		16
P.CST4	Channel 3	Channel 7	Channel 13	RBT Channel		17
T.CPZ	0003		Storage Move Flag		Machine Field Length	20
T.MON	* MONITOR * 55					21
T.STD						22
T.CIDLE			seconds		milliseconds	23
T.PIDLE			seconds		milliseconds	24
T.PPR						25
						26
T.JDATE	0000	0000	00	Y Y	D D D	27

T.CLK	* H H . M M . S S *										30									
T.DATE	b M M / D D / Y Y b										31									
T.SLAB2	SCOPE VERSION 3.X										32									
:																				33
:																				
T.SLAB6											36									
T.MSP										Step Flag	37									
T.MSC	future development	seconds	milliseconds mod 1000 ₁₀	millisecond clock							40									
T.PPS0			(MILLISEC. SINCE LAST SECOND UPDATE)	seconds	milliseconds						41									
:											42									
:											43									
:											44									
T.PPS9											45									
											52									
											53									
T.CPZECS											54									
T.TMP	MTR Scratch										55									
T.CPT1/T.UAS	Unassigned GM/100B	Unassigned ECS/1000B	NN C X	CPU-B ACTIVITY	CPU-A Activity						56									
											57									

NN = CP #

C = 1 ⇒ CPU OFF

2 ⇒ DEDICATED

3 ⇒ DOESN'T EXIST

X = 1 ⇒ CPU-A

2 ⇒ CPU-B

2000 ⇒ STORAGE MOVE

0000 ⇒ IDLE PROGRAM

0200 ⇒ ASSIGNED TO CP1

0400 ⇒ ASSIGNED TO CP2

⋮ ⋮ ⋮ ⋮ ⋮

1600 ⇒ ASSIGNED TO CP7

PERIPHERAL PROCESSOR COMMUNICATION AREAS

T.PPC1	PP NAME OR ZERO					60 W.PPIR
						61 W.PPOR
						62 W.PPMES1
						63 W.PPMES2
						64 W.PPMES3
						65 W.PPMES4
						66 W.PPMES5
T.PPC2						67 W.PPMES6
						70
T.PPC10	T.PPC10 is not presently used					170

CONTROL POINT AREAS

T.CPA1

EXCHANGE PACKAGE

200

FLAGS

217

220

CONTROL CARD OVERFLOW BUFFER

240

241

CONTROL CARD BUFFER

250

251

FLAGS

350

351

377

400

T.CPA2

T.CPA7

1777

EXCHANGE PACKAGE

				Words
	Program Address(P)	A0(Address Registers)	0 0 0 0 0 0	0
	Reference Address(RA)	A1	B1(Increment Register)	1
	Field Length(FL)	A2	B2	2
	Exit Mode(EM)	A3	B3	3
	RA - ECS	A4	B4	4
	FL - ECS	A5	B5	5
	<i>M A</i>	A6	B6	6
		A7	B7	7
X0 (Operand Registers)				10
X1				11
X2				12
X3				13
X4				14
X5				15
X6				16
X7				17

CONTROL POINT AREA

	57	48 47	36 35	24 23	12 11	20
W.CPSTAT/ W.CPEF	C.CPSTAT (Status)	C.CPEF (Error Flag)	C.CPSM (Storage Move)	C.CPRA (RA/100B)	C.CPFL (FL/100B)	20
W.CPJNAM	JOB NAME				Next Control Statement	21
W.CPECS/W.CPPRI/ W.CPECS	Priority		Time Limit	ECS RA/1000B	ECS FL/1000B	22
W.CPTIME	millisecond count since last priority <i>RE-ENTRY</i>		C.CPTIML CPU Time (seconds)	milliseconds		23
W.PPTIME		Loader Flag	PPU Time (seconds)	milliseconds		24
W.CPRCL	PP Recall (Obsolescent)					25
W.SSW	Sense Switches/Lights					26
W.EQP	Equipment Assignments					27
W.CPDM	Last Dayfile Message					
W.CPERT		C.CPFST (FST Address)	C.CPFP	C.CPERT		40
W.CPTBUF	Control Card Overflow Buffer					41
W.CPTBUL						50
W.CPCAF	Control Card Buffer					51
W.CPCAL						150
W.FSTCC	FST Entry for Next Control Card PRU					151
W.FSTUR						152
W.CPOAE					Op. Assigned Equipment	153
W.CPVRNO	Visual Reel Number					154
W.CPENC	Stack Entry/ Exit Count	PP Job Entry Count				155
W.CPRCS1	<i>(USED BY RESPOND)</i>					156
W.CPAR	Auto-Recall Pointer					157
W.RES2	<i>(USED BY RESPOND)</i>					160
						161
						162
						163
						164
						165
						166
						167

↓ FST ADDRESS FOR
CHECKING STATUS CODE

CONCORDANCE OF CONTROL POINT FIELDS, 2.0/3.0

<u>Word</u>	<u>Mnemonic</u>	<u>Changed</u>	<u>Comments</u>
20	W.CPSTAT, etc.	no	
21	W.CPJNAM	no	
22	W.CPPRI, etc.	yes	Msg. count removed, Time Limit moved left one byte Operator assigned equipment move, RA/1000B of ECS in byte 3, FL/1000B of ECS in byte 4
23	W.CPTIME	yes	bytes 0 and 1 contain a msec count since last priority re-evaluation (altered by MTR only)
24	W.PPTIME	yes	express flag removed
25	W.CPRCL	no	this work is being phased out in favor of MTR function 37 (see IMS)
26	W.SSW	no	
27	W.EQP	no	
30-37	W.CPDEFM	no	
40	W.CPERT	yes	now contains pointers and flags for control card processing in bytes 2, 3, 4
41-50	W.CPTBUF	no	Control Card overflow buffer
51-150	W.CPTBUL	no	Control Card buffer
151	W.FSTCC	yes	FST entry for next control card PRU
152	unused	yes	
153	W.CPOAE	yes	Operator assigned equipment (see 22) in byte 4
154	W.CPVARN unused	yes	
155	W.CPENC	yes	I/O stack entry count in byte 0, PP Job Buffer Entry Count in byte 1
156	unused	yes	

CONCORDANCE OF CONTROL POINT DISPOSITION

<u>Word</u>	<u>Monomelic</u>	<u>Classed</u>	<u>Comments</u>
157	W.CPAR	yes	Auto Recall pointer in bytes 0-1
160-177	unused	yes	

CP. SM	SB7	I
	NG	B3, DOWN
	SB6	-2
	SA1	B2-B7
	SA2	B2+B6
	EQ	LOOP
DOWN	SB6	2
	SA1	B1
	SA2	B1+B7
LOOP	SA4	A1+B6
	SA5	A2+B6
	BX6	X1
	LX7	X2
	SA6	A1+B3
	SA7	A2+B3
	SB1	B1+4
	SA1	A4+B6
	SA2	A5+B6
	BX6	X4
	LX7	X5
	SA6	A4+B3
	SA7	A5+B3
	LT	B1, B2, LOOP
	JP	0

3-15-1

STORAGE MOVE PROGRAM

IFNE	IP.MECS,0		CMR	00337
	MOVE EXTENDED CORE STORAGE		CMR	00338
	PARAMETERS PASSED IN EXCHANGE PACKAGE		CMR	00339
			CMR	00340
			CMR	00341
P	CP.ECSM		CMR	00342
RA(CM)	0		CMR	00343
FL(CM)	400000B		CMR	00344
EM	0		CMR	00345
RA(ECS)	0		CMR	00346
FL(ECS)	10000000B		CMR	00347
			CMR	00348
R1	RA(ECS) OF CTL PT TO BE MOVED /100B		CMR	00349
B2	RA+FL (ECS) /100B		CMR	00350
R3	DISPLACEMENT /100B		CMR	00351
R4	LENGTH OF CM BUFFER AREA		CMR	00352
A0	ADDRESS OF CM BUFFER AREA		CMR	00353
			CMR	00354
CP.FCSM	SX1	R1	INITIALIZE THE X REGISTERS	CMR
	LX1	6		CMR
	SX2	B2		CMR
	LX2	6		CMR
	SX3	B3		CMR
	LX3	6		CMR
	SX5	B4		CMR
	LX5	6		CMR
	NG	X3,SM34	IF DISPLACEMENT IS NEG, SHUTTLE DOWN	CMR
	EQ	SM24	ELSE, SHUTTLE UP	CMR
	SPACE	3		CMR
		SHUTTLE UP LOOP (ECS)		CMR
SM18	IX6	X5-X4	WHEN THE REMAINING PORTION TO BE	CMR
	NG	X6,SM20	MOVED IS LESS THAN THE BUFFER	CMR
	RX5	X4	REDUCE THE BUFFER SIZE	CMR
SM20	SB5	X5		CMR
	IX0	X2-X5	ECS ADDRESS OF THE MOVE	CMR
	RX2	X0		CMR
	RE	B5+0	READ INTO THE CM BUFFER	CMR
	IP	*		CMR
	IX0	X0+X3	ADJUST THE ECS ADDRESS	CMR
	WE	B5+0	WRITE BACK INTO ECS	CMR
	JP	*		CMR
SM24	IX4	X2-X1	CONTINUE TO LOOP UNTIL THE ENTIRE	CMR
	NZ	X4,SM18	BLOCK HAS BEEN MOVED,	CMR
	JP	0	THEN EXIT	CMR

3-16-2


```

SPACE      2
*          SHUTTLE DOWN LOOP (ECS)
SM26      IX6      X5-X4
          NG      X6,SM28
          RX5      X4
SM28      SB5      X5
          BX0      X1
          IX1      X1+X5
          RE      B5+0
          JP      *
          IX0      X0+X3
          WE      B5+0
          JP      *
SM34      IX4      X2-X1
          NZ      X4,SM26
          JP      0
          ENOIF
    
```

```

CMR      00381
CMR      00382
CMR      00383
CMR      00384
CMR      00385
CMR      00386
CMR      00387
CMR      00388
CMR      00389
CMR      00390
CMR      00391
CMR      00392
CMR      00393
CMR      00394
CMR      00395
CMR      00396
CMR      00397
    
```

BSSZ T.EST-*

CMR 00399

3-16

EQUIPMENT STATUS TABLE

POINTER

P. EST

T. EST	T. EST + L. EST	0 ————— 0
--------	--------------------	-----------

5

1600

EST

59 47 35 23 11 0

T. EST

3000 Type
nonallocatable

Z	bb	aa	dd	cc	o	h	s	e	u
---	----	----	----	----	---	---	---	---	---

2040 → 00

6000 Type
nonallocatable

Z		cc	e Ouu	o	h	0000
---	--	----	-------	---	---	------

L. EST
WORDS
IN
LENGTH

Allocatable

4 0 0 0 Z	O	O	o	h	O
--------------	---	---	---	---	---

59 47 35 23 11 0

- Z - Status of Entry
- aa, bb, cc, dd - Channel Numbers
- o - ON/OFF Bit
- h - hardware type
- s - 6681 number
- e - equipment
- u - unit

4000 ⇒ allocatable device
6000 ⇒ "non allocatable, and unassigned" or "no entry"
0200
0400
0600
0800
1000
1200
1400
1600
} nonallocatable and assigned.

~~2000~~

(3.2)

6000 + 3000
allocatable.

FILE NAME TABLE ENTRY

3-17

WORD "1"
(FNT)

FILE NAME										TYPE		LOCK		CONTROL POINT			
										18	17	16	15	14	12	11	ED
														PRIORITY			

WORD
"2"

<p>READ CARD FILE</p> <p>INPUT FILE</p> <p>DISC FILES RANDOM (ALLOCATABLE) FILE</p> <p>Tape FILES INDIVIDUAL SEQUENTIAL FILE</p> <p>MULTI-FILE FNT/FST ENTRY</p> <p>PUNCH CARD FILE</p>	EQUIP MENT CODE	RECORD COUNT	RECORD TYPE	EST ORDINAL	NOT USED	E O J	CARD COUNT WITHIN RECORD
	59 54	52 112	51 26	35	24 23	18 17 16	15 12 11 001
	EQUIP MENT CODE	NOT USED	FIRST RBT WORD PAIR	ECS FIELD ENTRY/1000	CURRENT RBT ENTRY	CURRENT RBT ENTRY	CM FIELD LENGTH/100
	59 54	53 48 47		35	24 23		11 00
	EQUIP- MENT CODE	NOT USED	FIRST RBT WORD PAIR	CURRENT RBT WORD PAIR	CURRENT RBT ENTRY	CURRENT RBT ENTRY	CURRENT PRU NO.
	59 54	53 48 47	36	35	24 23	15 14 13 12 11	00
	EQUIPMENT CODE	D E N S I T Y	EST ORDINAL OF SECONDARY DEVICE	EST ORDINAL OF PRIMARY DEVICE	CURRENT PRU NO.		
	59 54	53 52 51 50 49 48 47	36 35	24 23	00		
	EQUIP. CODE	T D A P N E S S I T Y	SECONDARY DEVICE ORDINAL	PRIMARY DEVICE ORDINAL	POSITION # OF THE LAST FILE WRITTEN IN THE SEQUENCE OF FILES.		CURRENT FILE POSITION
	59 54	53 52 51 50 49 48 47	36 35	24 23	12 11		00
	EQUIP MENT CODE		EST ORDINAL	FLAG USED BY 2PC	E O J	CARD COUNT	
	59 54 53		36 35	24 23	18 17 16	00	

WORD "3"
(FST 2)

FUTURE USE	LAST USED FET ADDRESS	DISPOSITION CODE	SE- CURITY CODE	LAST CODE AND STATUS
59 58 54 53	36 35	24 23	18 17	00

NOTE: On Multitrac files, the FNT word "3", bits 18+19 are always set (File always "CLOSED"), bits 24-35 = 0002 (multi file dispositions), and bits 48 and 49 are used for logical security code.

Record types continued.

002 - 12.0₈ ⇒ "Free-form" word record
(See Ref. Man. Appendix E)
177 ⇒ Special mode ended.

Record Count within this file.

Equipment Code, see table on Page 3-4
and 3-5 of Ref. Manual.
(60₈ ⇒ card reader, etc.)

CM Field Length / 100₈ for this file.

ECS Field Length / 1000₈ for this file.

Current RBT entry. This is a count of the number of full RBT word pairs up to the current position of the file. For example, if the file is positioned in the second RBT word pair, the current RBT entry will be 1. Just as soon as the second RBT pair fills up, this number will roll over to 2.

Current RBT Byte indicates which byte in the RBT word pair contains the RB address.

Current RBT word pair position. If this number = N, the address of the first word of this word pair is $400,000B - 2 \times N$.

First RBT word pair associated with this file. If this number = M, the address of the first word of this word pair is $400,000B - 2 \times M$.

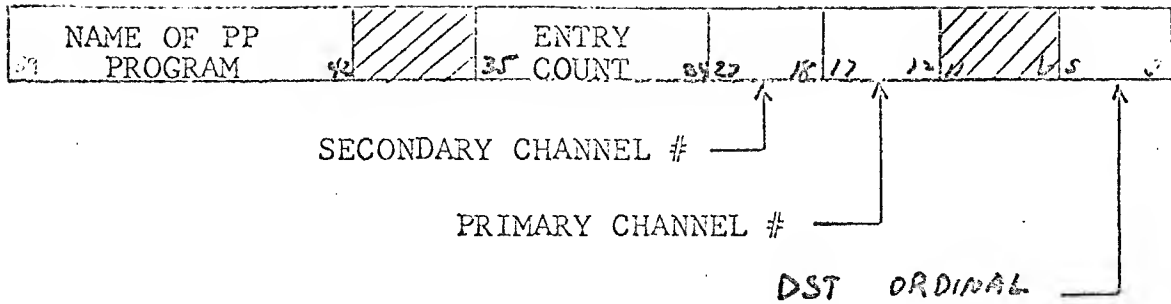
Check point flag, is used with checkpoint/restart files.

Current PRU # indicates which sector is currently being read.

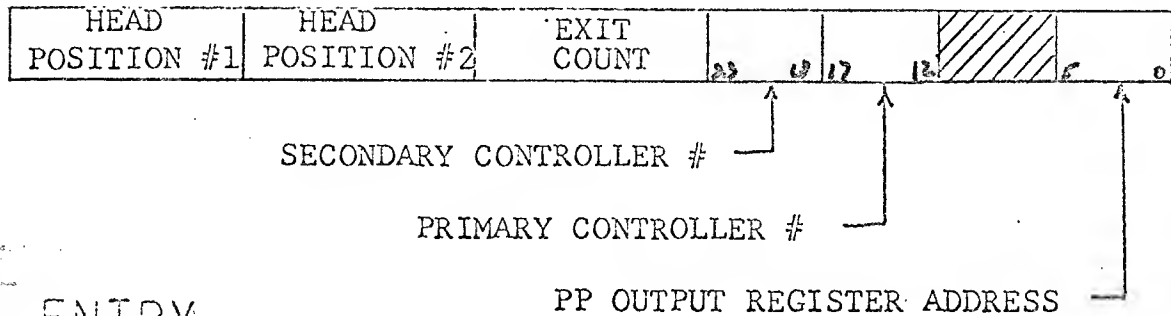
DEVICE STATUS TABLE ENTRY

3-18 TNDM
3-20

WORD 1

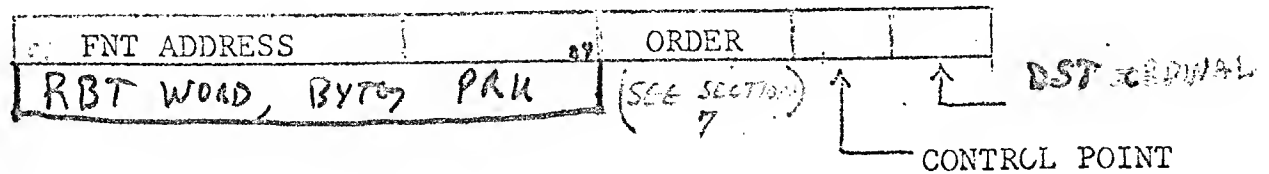


WORD 2

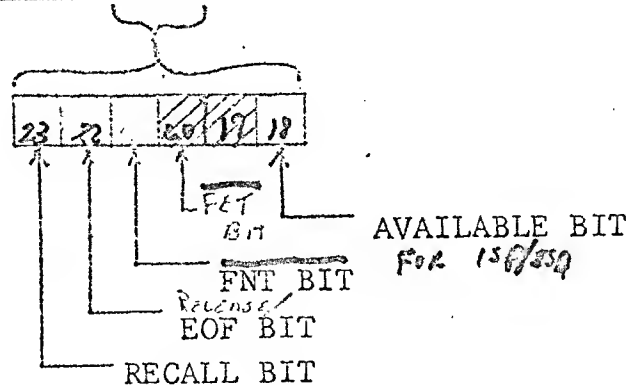
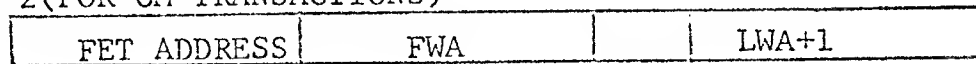


STACK ENTRY

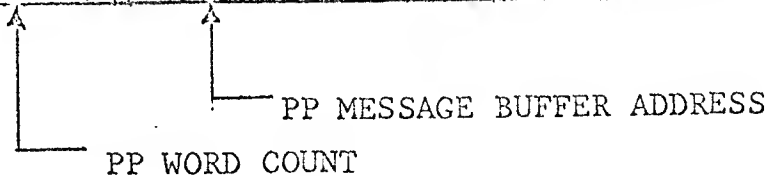
WORD 1



WORD 2 (FOR CM TRANSACTIONS)



WORD 2 (FOR PP TRANSACTIONS)



A Record Block Reservation (RBR) table serves a dual purpose:

1. It defines record blocks in a specific allocatable area.
2. It identifies those record blocks within that area as available or not available.

A maximum of 2048 record blocks may be defined by one RBR. The area described by an RBR must all be on one device.

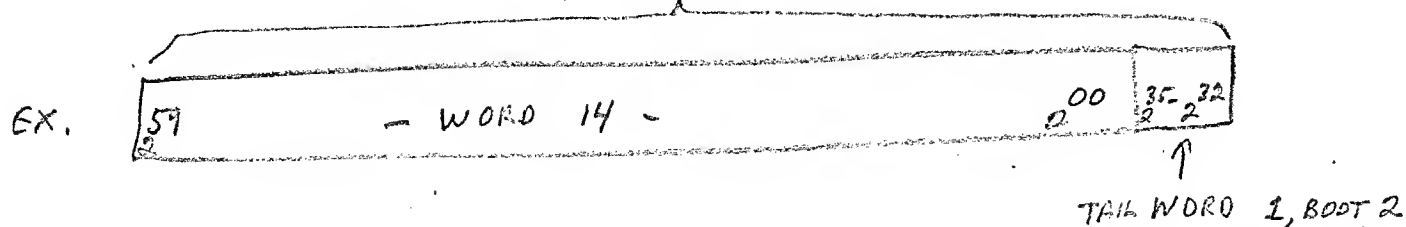
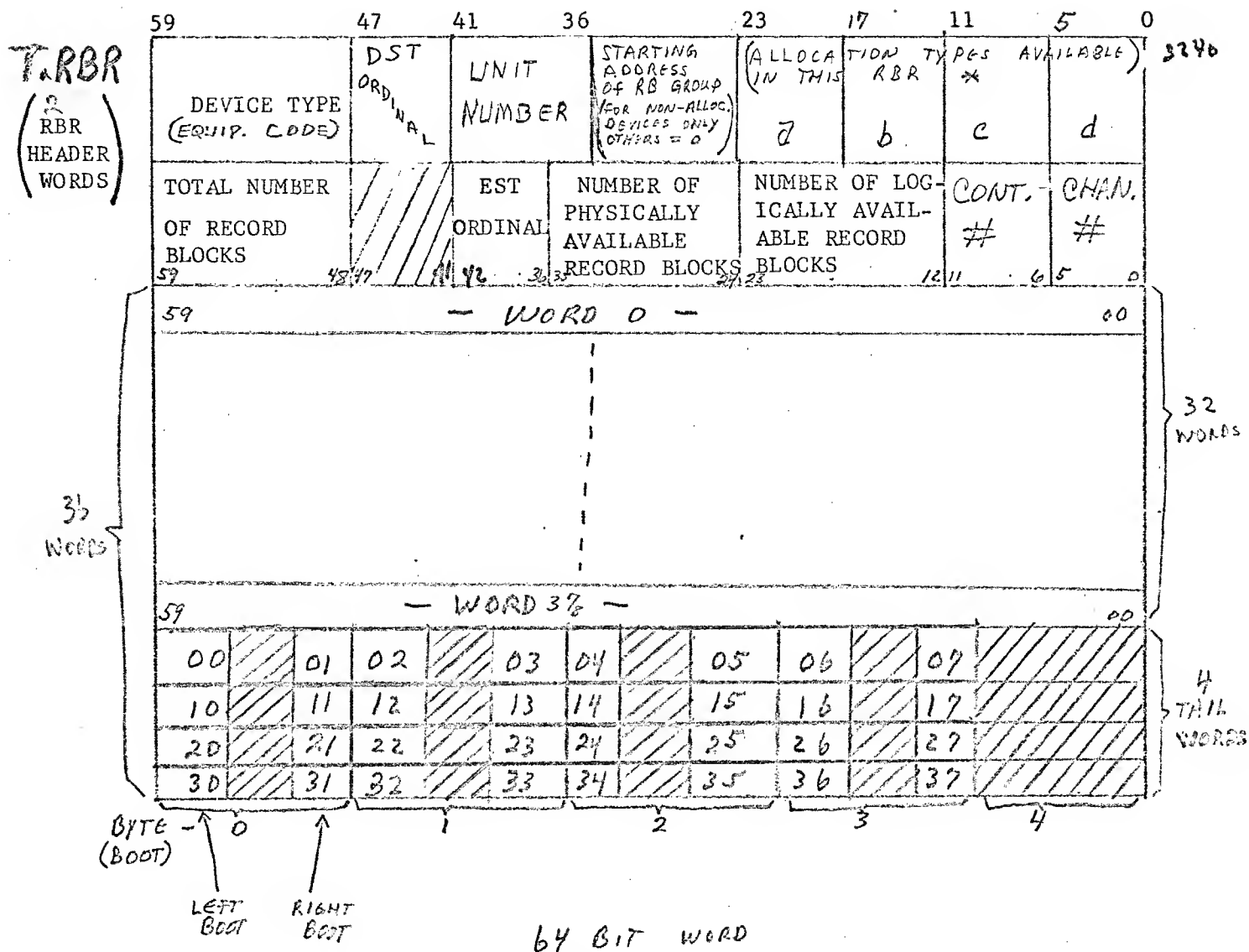
The RBR tables in the SCOPE Operating System follow immediately after the FNT/FST area. The pointer to the first word of the first RBR table is found in Word 2, Byte 1 (counting 0-4, left to right) of low core. The second RBR table starts immediately after the first. The total length of each table is 38_{10} CM words. Thus, adding 46 (octal) to the pointer in Word 2, Byte 1 will give the starting address of the second RBR table.

The header words (first two) of each RBR table are comprised of the following fields:

Word

- Bits 0 -11: Allocation type link. The ordinal of an RBR which references a device of the same type and which uses the same record block definition. The linkage should be circular.
- Bits 12-23: Device link. The ordinal of an RBR which references a device of the same type. The linkage should be circular so that all RBR's referencing the same device type are linked.

RECORD BLOCK RESERVATION TABLE (RBR)



SEE PAGE 3-21-4 for order of reservations.

* If standard allocation for the device is included, it is listed first, and the last 6-bit field is left zero. If all allocations are non-standard, all 4 fields must be non-zero. (Repeat last allocation type to end).

The allocation type link and device link may be broken up into 2 bytes of the form: * \Rightarrow (see SYSTEMS BULLETIN #10, pages 25-27)

aabbccdd=ALLOC

where each term ii specifies a legitimate allocation within the portion of the device indexed by this RBR.

01=50 PRU's/RB.

02=64 PRU's/RB.

00=Free allocation; when the user does not specify allocation in an OPEN call, the file may be assigned to any RBR with some ii=00. In this case, allocation aa is assigned to the file after the RBR is chosen, thus, aa should not be 00.

03=Free allocation; this allocation style is used to permit continuation of a file from one device to another (theoretically).

10=8 PRU's/RB (Used by RESPOND).

aa determines the actual RBR type for physical allocation purposes. ALLOC need not be supplied. Default values are:

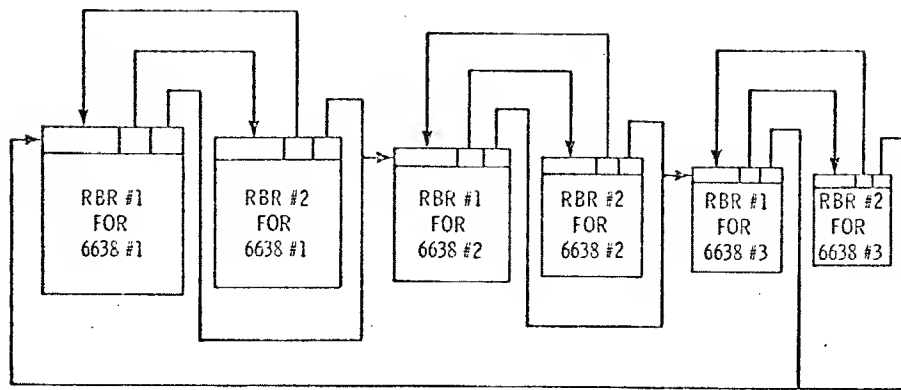
for TYPE=1 (bits 48-50)

ALLOC=03020100

for TYPE=2

ALLOC=03010100

- Bits 24-35: RB group start address - the half-track on which the described area starts.
- Bits 36-41: Unit number.
- Bits 42-47: Device number. The ordinal in the device status table (DST) of the device to which this RBR refers.
- Bits 48-59: Device group. The device type and allocation type used in the area described by this RBR.



6638 RBR LINKAGES FOR 3 6638's

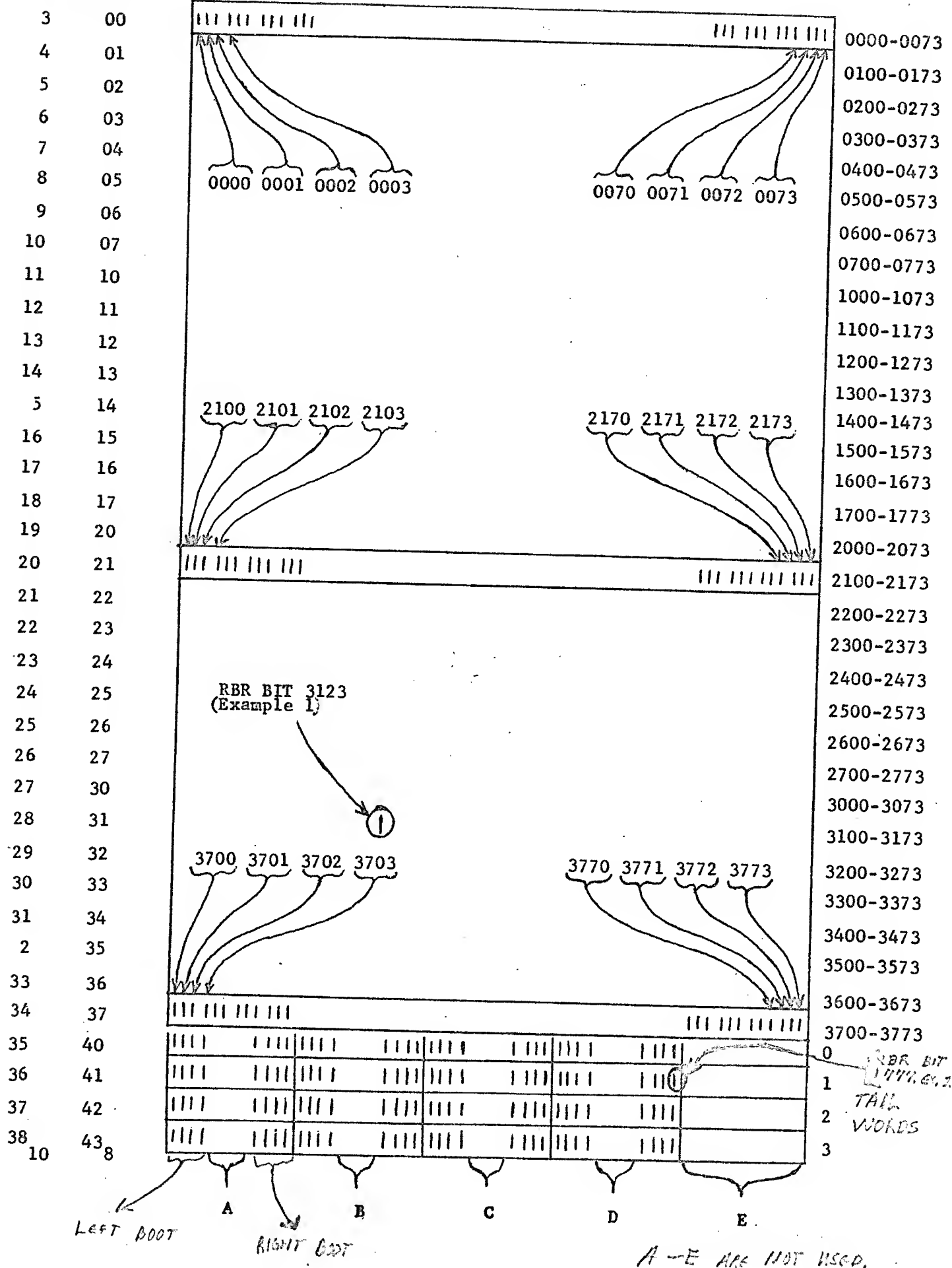
Word TWO

- Bits 0 - 5: Channel number. The channel used by the device referenced by this RBR.
- Bits 6 -11: Controller number.
- Bits 12-23: Logical availability. The number of record blocks not assigned to files.
- Bits 24-35: Physical availability. The number of usable record blocks. This is normally the same as the total RB count.
- Bits 36-41: EST ordinal.
- Bits 42-47: Not used.
- Bits 48-59: Total RB count. The number of record blocks defined by this RBR.

WORDS

RBR WORDS 0-378

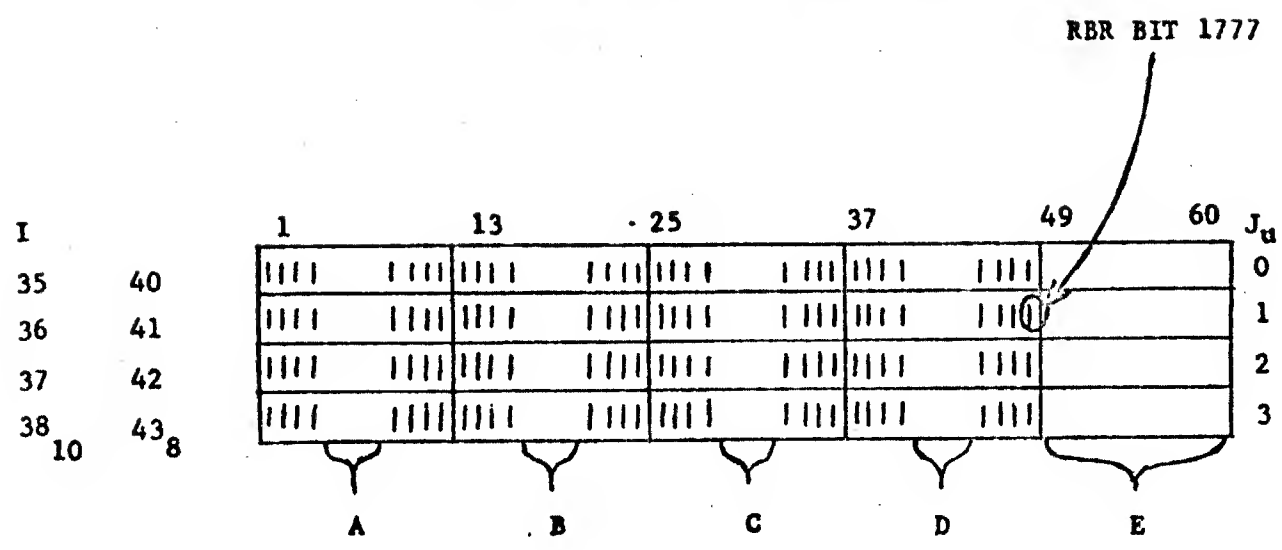
3-21-4



A-E ARE NOT USED.

RBR "TAIL WORDS"

(The last 4 words in an RBR Table.)



The bits designated by A, B, C, D, and E are not used.

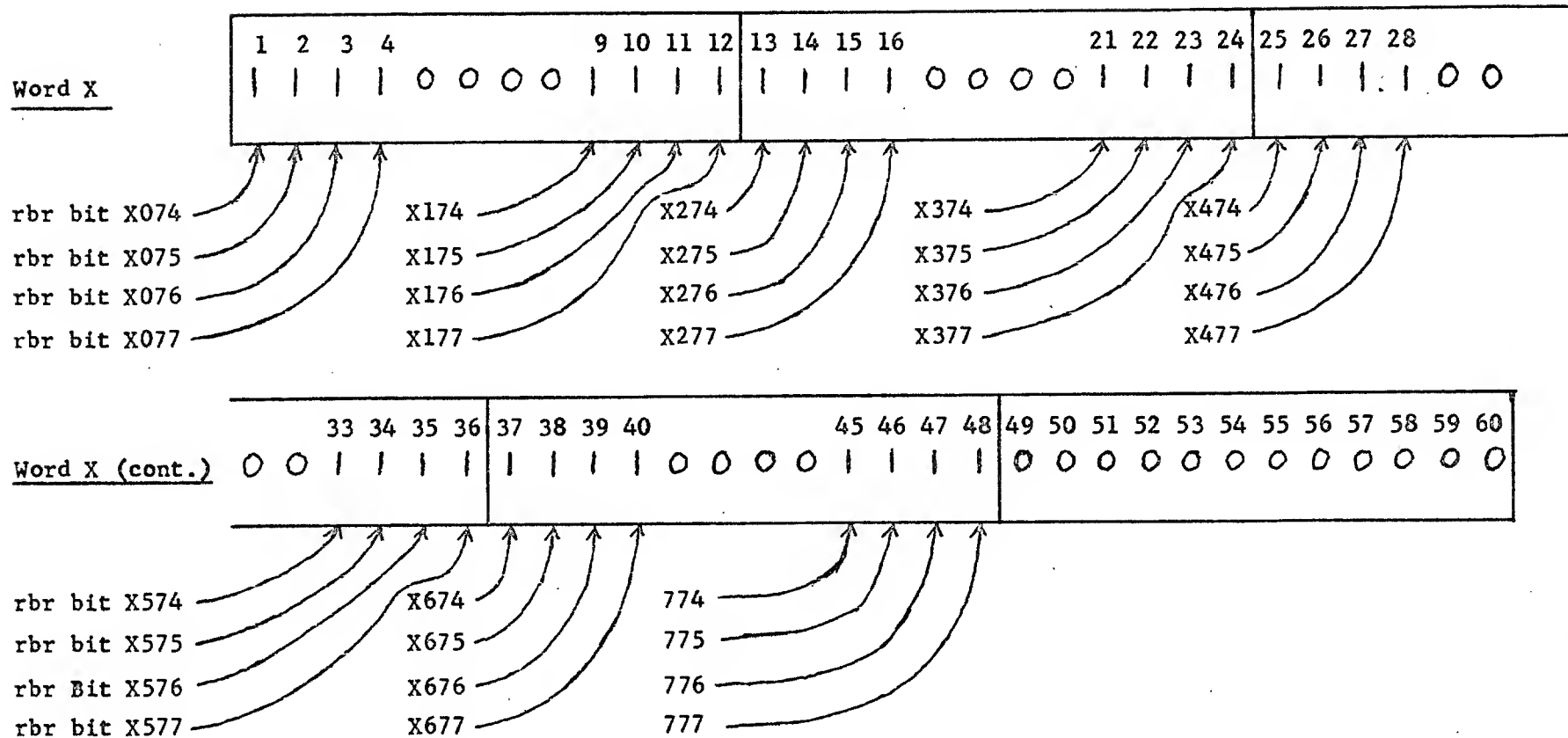
- A = Bits 52 - 55 (as usually numbered) or 5 - 8 (as numbered above)
- B = Bits 40 - 43 (as usually numbered) or 17 - 20 (as numbered above)
- C = Bits 28 - 31 (as usually numbered) or 29 - 32 (as numbered above)
- D = Bits 16 - 19 (as usually numbered) or 41 - 44 (as numbered above)
- E = Bits 0 - 11 (as usually numbered) or 49 - 60 (as numbered above)

A number corresponding to a bit in the RBR Table may be calculated from and RB number found in one of the 0-7 bytes of an RBT word pair. Suppose the RB number is N (N is always odd). Then the RBR bit number, NRBR, is given by the equation (using truncated integer arithmetic):

$$NRBR = N/2 = J_u \times 8^3 + J_L \times 8^2 + L_1 \times 8^1 + L_2 \times 8^0 \quad (\text{eqn. 1})$$

$J = J_u \times 8^3 + J_L \times 8^2 =$ Main table pointer to words 3-34 (decimal) or 0-37 (octal) of the RBR table. The entire J pointer is used when the last 2 octal digits of NRBR are from 0-73.

DETAIL OF REPRESENTATIVE TAIL WORD



X is 0, 1, 2, or 3, depending on whether we are in tail word 40, 41, 42, or 43.

3-21-6

$L = L_1 \times 8^1 + L_2 \times 8^0$ = Lower 2 octal digits of RBR bit number. If L is 74-77 (octal) then J_u determines which of the tail words (0-3) the designated RBR bit falls in. Also, J_L determines in which byte in the tail word the bit falls. For $J_L = 0$ to 1, byte number = 1 (counting bytes left to right, 1-5).
 For $J_L = 2$ to 3, byte number = 2.
 For $J_L = 4$ to 5, byte number = 3.
 For $J_L = 6$ to 7, byte number = 4.

A formula for obtaining the exact position of a given bit in one of the tail words (numbering the bits from left to right, 1-60) is given by:

$$\text{BIT POSITION} = \text{LPOSN}(J_L + 1) + L_2 - 4 \quad (\text{eqn. 2})$$

where:

$$\text{LPOSN}(1) = 1$$

$$\text{LPOSN}(2) = 9$$

$$\text{LPOSN}(3) = 13$$

$$\text{LPOSN}(4) = 21$$

$$\text{LPOSN}(5) = 25$$

$$\text{LPOSN}(6) = 33$$

$$\text{LPOSN}(7) = 37$$

$$\text{LPOSN}(8) = 45$$

These values are decimal.

Sample calculation 1:

One of the 0-7 bytes in an RBT word pair contains the RB number 6247. Find the corresponding bit position in the RBR table.

Applying equation 1:

$$NRBR = 6247/2 = 3123 \quad (\text{dropping the fraction})$$

L = last 2 digits = 23, which is between 0 and 73, and thus the desired bit is in one of the RBR table words from 0 to 37 (octal).

$J = 31$, which points to the 31st (octal) word in the RBR table.

Thus, the desired bit is bit 23 (octal - starting with zero and counting from left to right) of word 31 of the RBR table. The location of this bit is circled on the diagram on Page 12.

Sample calculation 2:

One of the 0-7 bytes in an RBT word pair contains the RB number 3777.

Locate the corresponding bit in the RBR table.

Applying equation 1 from page 13:

$$J_u \ J_L \ L_1 \ L_2$$

$$NRBR = 3777/2 = 1777$$

$L = 77$, which implies that the bit is in one of the tail words.

$J_u = 1$, so the bit is in the second tail word (word 41₈).

$J_L = 7$ = second digit from left.

L_2 = low order digit = 7.

Applying equation 2 from page 14:

$$\text{BIT POSITION} = \text{LPOSN}(J_L + 1) + L_2 - 4$$

$$\text{BIT POSITION} = \text{LPOSN}(8) + 7 - 4$$

$$\text{BIT POSITION} = 45 + 3 = 48$$

Example #2.

- [illegible]

4. The lower 2 bits of the lower 6 bits indicate the 2^{xx} position in the boot. In this case 11, or the right-most bit of the boot. (Remember they are reserved from left to right).
5. Thus actual RB address # 1777 is tail word # 1, boot 3, right boot, bit 3. Or 2^{12} in tail word # 1.

Thus, the desired bit is bit 48 (decimal - starting with one and counting from left to right) of the second tail word (word 41). Or, if we count in octal, and start with zero, the bit is number 57 (still counting from left to right). This bit is circled on the diagram on Page 13.

DISCOVERING PARTY ERRORS ON DISK

Sample:

RBR	RB	PRH
01	0741	012

RBR

00	FILE 0	6638
01	FILE 1	6638

Then - RBR RB PKU
01 0741 012

would be decoded as -

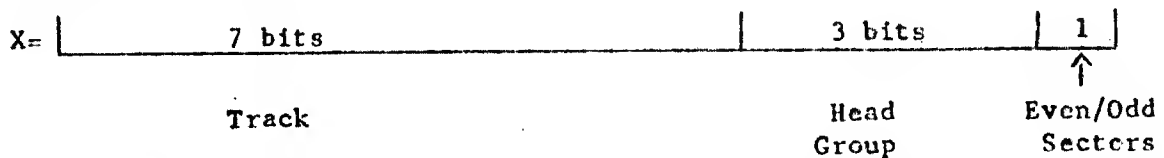
- 6638, file 1
- position 7
- odor half block
- Head group 20.
- sector 12.

Note - the RB # is actual RB address,
not as shown in RBT system.

Physical Position on the Disk

Another question which arises is how one computes the track number, head group, and whether we are writing on odd or even sectors. These quantities may also be derived from the RB numbers found in the 0-7 bytes of the RET word pairs. The steps are as follow:

1. Assume the RB number is N (12 bits).
2. Divide N by 2 (shift right by 1), leaving us with an eleven bit quantity, X.
3. Divide X up into three groups as indicated below, thus giving us the desired quantities:



Thus, the track numbers can go from 0 to 177₈ (0 to 128₁₀) and the head groups can go from 0 to 7. From these numbers, we can calculate the total possible number of half tracks per RBR Table:

$$\text{Total} = (128 \text{ tracks/RBR}) \times (8 \text{ head groups}) \times (2 \text{ half-tracks/track})$$

so:

$$\text{Total number of half-tracks} = 2048.$$

Figures 1-4 on Pages 1-9 of the 6638 manual is somewhat misleading in that it seems to indicate that head groups are specified by two digits. Each RBR Table is for one stack. Thus, if we are going to use only one octal digit to specify a head group, it would seem necessary to supply the File

Unit head r (0 or 1), also. However, this is easily taken care of by supplying the track number. Thus, there are head groups 0-7 for File Unit 0 and 0-7 for File Unit 1 (for stack 0, as well as stack 1). In order to specify any unique location on one stack, it is sufficient to give:

1. The 3 bit head group number,
2. The 7 bit track number, and
3. The 1 bit even/odd sector indicator (0 for even sectors, 1 for odd sectors).

6638 System for Discerning Parity Errors on Disc File

RBR

00 File 1 (Bryant Disk File)
 01
 02 File 0 (6638)
 03 File 1 (6638)

RB

00 00

Upper 6 (six) bits signify position. Lowest bit signifies odd or even halftrack. The remaining bits will be right shifted one and this will give you the Hd group.

PRU

Signifies the sector #.

Example: RBR RB PRU
 03 0741 012

would be decoded as 6638 File 1
 Position 7
 odd halftrack
 Hd group 20
 sector 12

0741 = 2^4 + 2^3 + 2^2 + 2^1 = 2^6 Position

DAYFILE BUFFERS

DAYFILE POINTER

P.DFB

T.DFB/ _{10B}	0 ————— 0
-----------------------	-----------

3

DFB

T.DFB

SYSTEM DAYFILE FET (LFN = DAYFILE)	12 ₈
CONTROL POINT 1 FET (LFN = DFILE1)	10 ₈
CONTROL POINT 2 FET (LFN = DFILE2)	10 ₈
CONTROL POINT 3 FET (LFN = DFILE3)	10 ₈
CONTROL POINT 4 FET (LFN = DFILE4)	10 ₈
CONTROL POINT 5 FET (LFN = DFILE5)	10 ₈
CONTROL POINT 6 FET (LFN = DFILE6)	10 ₈
CONTROL POINT 7 FET (LFN = DFILE7)	10 ₈
SYSTEM DAYFILE BUFFER	100 ₈
CONTROL POINT 7 DAYFILE BUFFER	

3140

100₈ - 2K24
CM
WORDS

3240

3340

100₈

THE FORMAT OF THE CMR DIRECTORY IS THIS --
IN CELL P.DIR (I.E. ABSOLUTE 1)

VFD 24/A,24/B,12/O
WHERE A IS THE ADDRESS OF THE FIRST WORD IN THE DIRECTORY,
AND B IS THE ADDRESS OF THE LAST+1 WORD.
IN CELL A

VFD 60/C
THEN THE ENTRY POINT TABLE LIES IN CELLS A+1 TO C-1 INCLUSIVE.
EACH ENTRY IN THIS TABLE HAS THE FORM

VFD 42/D,18/E
WHERE D IS THE NAME OF THE ENTRY POINT, AND E IS THE NUMBER OF
THE PROGRAM, WHICH WILL BE A CP PROGRAM NOT AN OVERLAY, TO WHICH
IT BELONGS. THE PROGRAM NAME TABLE ENTRY FOR THE PROGRAM CAN BE
FOUND IN CELLS C+1+2E AND C+1+2E+1.
IN CELL C

VFD 36/F,24/G
THEN THE PROGRAM NAME TABLE LIES IN CELLS C+1 TO G-1 INCLUSIVE.
F IS THE ADDRESS OF THE FIRST WORD OF THE FIRST PROGRAM NAME
TABLE ENTRY FOR A NON-PP PROGRAM. THIS HELPS THE PP RESIDENT,
WHICH SEARCHES ONLY FOR PP PROGRAMS IN THE PROGRAM NAME TABLE,
AND NEED SEARCH ONLY FROM C+1 TO F-2

EACH PROGRAM NAME TABLE ENTRY IS TWO WORDS. THE FIRST WORD IS
VFD 42/H,18/K

WHERE H IS THE NAME OF THE PROGRAM, AND K IS ITS LENGTH --
NOT ITS LENGTH AFTER LOADING, BUT ITS LENGTH IN THE LIBRARY.
K DOES NOT INCLUDE THE PREFIX OF THE PROGRAM RECORD, SO
THIS RECORD ON FILE SYSTEM IS ACTUALLY K+3 WORDS LONG.
BUT IF THE PROGRAM IS CENTRAL MEMORY RESIDENT, ITS BODY IN
THE DIRECTORY WILL BE K WORDS LONG.

NOTE THAT NO TWO NON-STITCH PROGRAMS MAY HAVE THE SAME NAME,
EVEN IF OF DIFFERENT TYPES. BUT STITCH PROGRAMS ARE ABSOLUTELY
IGNORED IN CHECKING FOR DUPLICATION.
THE SECOND WORD OF THE PROGRAM NAME TABLE ENTRY HAS TWO POSSIBLE
FORMATS. IF THE PROGRAM IS CENTRAL MEMORY RESIDENT,

VFD 4/M,4/N,1/P,6/Q,21/R,9/S,3/W,12/T
WHERE M IS THE RESIDENCE, 0 FOR CENTRAL MEMORY, N IS THE
PROGRAM TYPE, 0 FOR PP, 1 FOR CP, 2 FOR OVERLAY, AND
4 FOR STITCH, P IS THE DISK FILE IN WHICH THE PROGRAM
RECORD IS TO BE FOUND, 0 FOR ((SYSTEM)) AND 1 FOR
((SSSSSSU)), Q IS THE EDITION NUMBER, R IS THE ADDRESS IN
CMR AT WHICH THE BODY OF THE PROGRAM BEGINS
(IT ENDS AT R+K-1), S IS THE RBT ORDINAL OF THE PROGRAM
RECORD IN DISK STORAGE, W IS THE BYTE NUMBER IN THE RBT WORD
PAIR, AND T IS THE PRU NUMBER AT WHICH IT BEGINS IN DISK STORAGE.
IF THE PROGRAM IS DISK RESIDENT,

VFD 4/M,4/N,1/P,6/Q,3/U,12/V,9/S,3/W,12/T
WHERE M IS 1 FOR DISK RESIDENCE, N, P, Q, S, T, AND W ARE THE
SAME AS ABOVE, U IS THE PHYSICAL UNIT NUMBER OF THE
PROGRAM RECORD IN DISK STORAGE, AND V IS THE RBT NUMBER,
NOT ORDINAL, OF THE RECORD.
THE PROGRAM NAME TABLE ENDS WITH THE WORD AT G-1, AND THE BODIES
OF CENTRAL MEMORY RESIDENT PROGRAMS, UNLESS THERE ARE NONE AND
G=B, EXTEND FROM G THROUGH B-1.

3-23-1

ENTRY POINT TABLE

1st word NOT
TYPICAL ENTRY
↓

FWA OF
PNT
↓

TYPICAL
ENTRY
FORMAT

0	...	0	4566	4360
59	...	17	...	0
03 17	ENTRY POINT NAME	03 0	000147	PROGRAM NUMBER

147
x 2
314
+ 1
315
4566
5705

PROGRAM NAME TABLE

1st word
NOT
TYPICAL
↓

FWA OF 1st
PP ENTRY IN PNT
↓

FWA OF
RESIDENT FILE
↓

TYPICAL
ENTRY FORMAT

0	...	005017	0	...	05335	4566
59	...	17	...	0
23 11 16 03 17 23 0	PROGRAM NAME	00000001	PROGRAM LENGTH	00000101101	CM	5105
RES- IDEN	PROG- TYPE	ED. NO.	FWA in CM	RBT ORDINAL	B Y T E	STARTING PRU NUMBER
59	36	55	51	50	44	41
00	01	00	00	00	00	00
<p>DISC FILE RECORD IN DISK STORAGE</p> <p>RECORD IS FOUND ON</p> <p>FILE 00 = SYSTEM</p> <p>1 = RESIDENT</p> <p>0 = PP</p> <p>1 = CPH (CANDIDATE)</p> <p>2 = CPH (NON-CANDIDATE)</p>						

00 = CM
01 = DISK

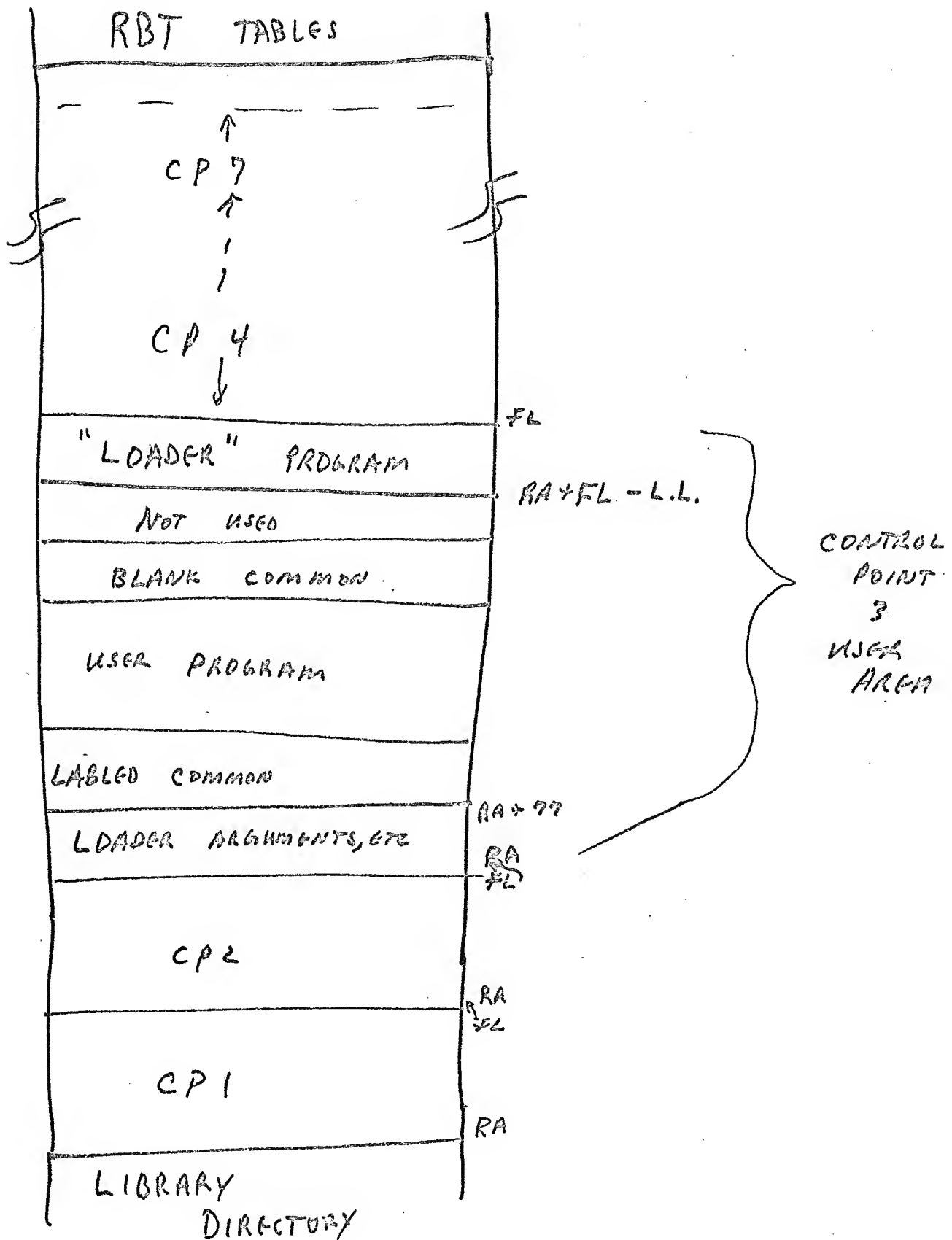
CM RESIDENT PROGRAMS

03 11 17 00	...	0773	0000	0144	5335
NAME OF PROGRAM IN DFC		PP START ADDRESS	1st of CM WORDS		

1st instruction, etc.

USER SPACE ALLOCATION

3-24



Record Block Tables (RBT)

An RBT consists of a series of 12 bit bytes identifying, in logical sequence, the record blocks assigned to a logical file residing on an allocatable device.

A maximum of 8,192 CM words may be occupied by all the RBT's active at any one time. Two CM words at a time are assigned to each RBT as required (Thus, we speak of RBT word pairs); the two words may reference 1 to 8 record blocks (RB's). Two consecutive RBT pairs for a particular file do not have to be contiguous. The first byte of the first word of each RBT pair contains a pointer to the next RBT in the chain. This pointer is the ordinal or relative position of the next RBT. If the pointer is N, the address of the word pair is calculated as:

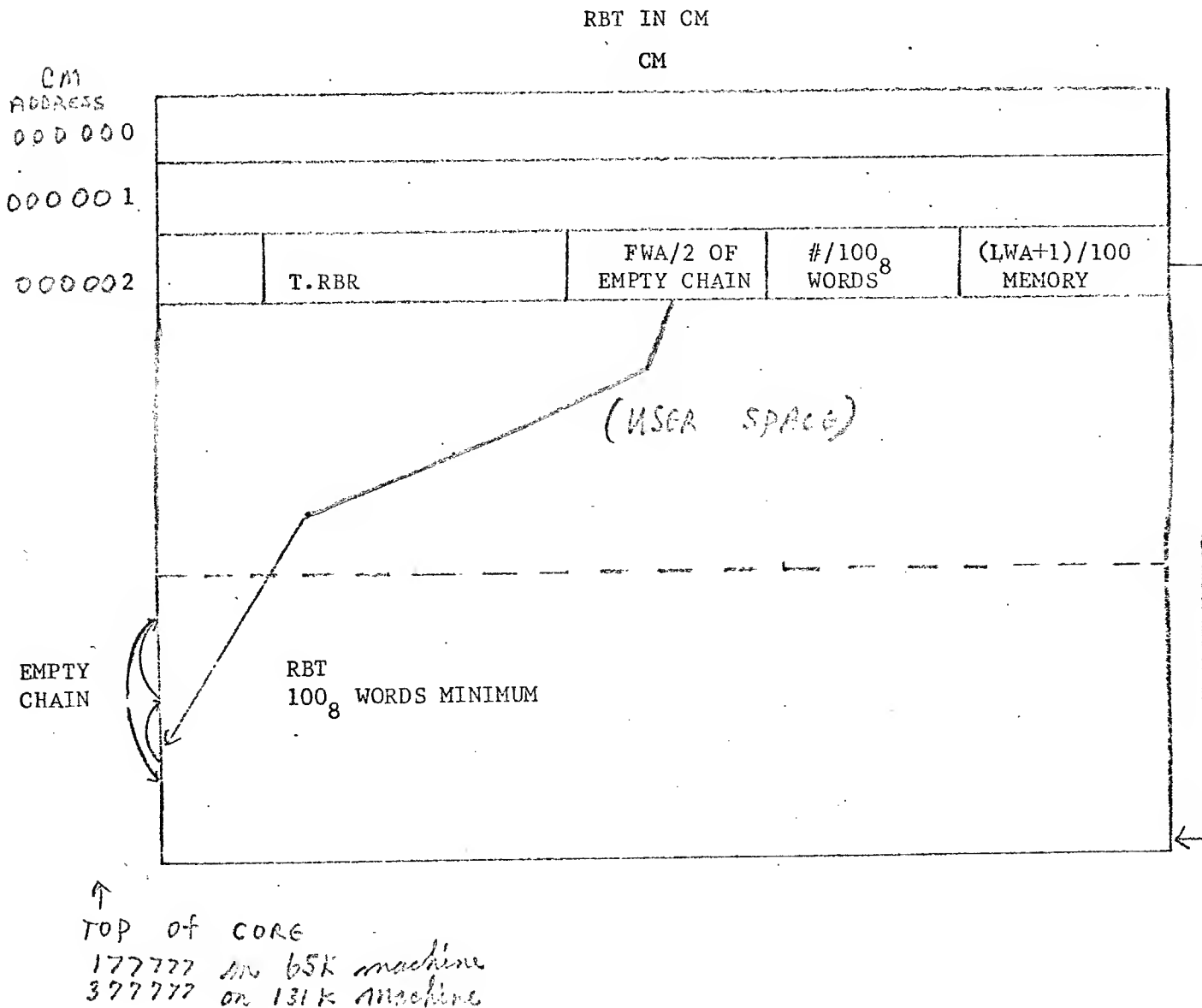
$$\text{ADDRESS} = (\text{LWA} + 1 \text{ of memory}) - 2 * N = 400,000\text{B} - 2 * N.$$

If the pointer is zero, the current RBT ends the chain.

The area reserved for RBT use may be reduced by control point space requirements. The maximum is defined by the number of word pairs referenceable by a twelve bit byte.

The RBT area starts at high core with its first word pair, and works down. The first word of the first RBT is at location 377,776; the first word of the second RBT is at location 377,774; etc. The number of RBT word pairs/100B is found in word 2, byte 3, (starting with 0, left to right) of low core.

The unoccupied RBT pairs are also chained together (blocks of word pairs are added or subtracted in groups of 100B words). The position of the start of the empty chain is found in the middle byte of word 2 of low core.



Byte Types That Can Occur in an RBT

RBT Link - (X-byte) The first byte of each RBT CM word pair. If blank, the chain of RBT word pairs assigned to a logical file terminates in the current word pair. If occupied, its contents multiplied by 2 and subtracted from the LWA+1 of central memory gives the address of the next RBT word pair in the chain.

RBR Link - (Y-byte) The high order 9 bits of the 2nd byte of each RBT CM word pair. It identifies the RBR related to the record block group in which the file resides. It may also appear in the third through the tenth byte of each word pair, if the file is contained partially in more than one RB group (We have not found an example of this.).

First RBT Byte - (Y-byte) The first byte used in the RBT word pair is indicated in the low order 3 bits of the 2nd byte of each word pair.

Byte 2 (middle byte) of the first word of a pair is counted as RBT byte 0. Byte 3 is RBT byte 1; byte 4 is RBT byte 2. Byte 0 of word 2 is RBT byte 3, etc. Byte 4 of word 2 is RBT byte 7. See Figure I on Page 25.

If the word pair is the first in an RBT chain the 3-bit "first RBT byte" field (low order 3 bits of the Y-byte) will be set to 2 for sequenced files and 4 for random files.

Flag and Allocation Type - (word 1, first pair - "0 byte") Bit 31 contains the release flag; if set, record blocks are to be released after reading.

Bit 30 is not clearly defined. ^(RANDOM BIT) (See 2.6.3 (4), Page 26, of SCOPE 3.0 IRS.)

Bits 24-29 contain the allocation type as defined in the ERS. 00 and 03 imply no restriction. 01 implies 50 sector record blocks. 02 implies 64 sector blocks.

Next PRU - (bits 12-23, word one - "1 byte") The first word pair in an RBT chain contains the ordinal of the next PRU in which writing is to occur for the referenced logical file. This field may differ from the same field in the FTN, which references the PRU following the last one read or written. This field occurs only in the first word pairs for a logical file.

Last Assigned Record Block - (bits 0-11, word one - "2 byte") - The first word pair for a random file contains the ordinal of the last record block assigned when the file was last written. This field occurs only in the first word pair of a randomly accessed logical file.

Last Assigned PRU - (bits 48-59, word 2 - "3 byte") - The first word pair for a random file contains the ordinal of the last record block assigned when the file was last written. This field occurs only in the first word pair of a randomly accessed logical file.

RB Link - A 12 bit quantity, appearing in one of the 0-7 bytes in an RBT word pair. The upper eleven bits, taken together with the last RBR link, define a unique position on the disk. The lower bit of an RB link is always set to distinguish it from an RBR link (which may also occur in bytes 0-7 when the file is contained partially in more than one RB group). This condition has not been observed yet, however, in our operations at the Los Angeles Data Center. Upon filling one RBR Table, whereupon we expected

the system to continue reserving half-tracks, for all files in the process of being written, in the other RBR Table, the system instead hung up. Dividing the RB link by 2 gives the position in the corresponding RBR Table. (See sample calculations on Page 13.)

See figures on following pages.

RBT WORD PAIR FORMATS

Figure I.

59	47	35	23	11	0
X-BYTE	Y-BYTE	0-BYTE	1-BYTE	2-BYTE	
3-BYTE	4-BYTE	5-BYTE	6-BYTE	7-BYTE	

SAMPLE
WORD PAIR

Figure II.

RBT LINK	RBR LINK	B Y T E	FLAGS AND ALLOCATION TYPE	NEXT PRU	1ST RB LINK
2ND RB LINK	3RD RB LINK		4TH RB LINK	5TH RB LINK	6TH RB LINK

FIRST RBT
PAIR FOR A
SEQUENTIAL
FILE

Figure III.

RBT LINK	RBR LINK	B Y T E	FLAGS AND ALLOCATION TYPE	NEXT PRU	LAST ASSIGNED RB
LAST ASSIGNED PRU	1ST RB LINK		2ND RB LINK	3RD RB LINK	4TH RB LINK

FIRST RBT
PAIR FOR A
RANDOM
FILE

Figure IV.

RBT LINK	RBR LINK	0	NTH RB LINK	N+1ST RB LINK	N+2ND RB LINK
N+3RD RB LINK	N+4TH RB LINK		N+5TH RB LINK	N+6TH RB LINK	N+7TH RB LINK

RBT PAIR INTER-
MEDIATE IN THE
CHAIN (I.E., NOT
FIRST OR LAST)

Figure V.

0 0 0 0	RBR LINK	0	MTH RB LINK	M+1ST RB LINK	M+2ND RB LINK
M+3RD RB LINK	LAST RB LINK		0 0 0 0	0 0 0 0	0 0 0 0

LAST RBT
PAIR OF A
PARTICULAR
CHAIN

NOTE: In the last example above, the X-BYTE is zero. The LAST RB LINK may fall in any of the bytes 0-7. The bytes after the LAST RB LINK in the last word pair are all zero.

SAMPLE RBT CHAIN

377654

377655

0	0	0	0	0	0	1	0	0	3	4	5	0	3	4	7	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

 52 } Last Pair
of
Sample Chain

377704

377705

377706

377707

0	0	5	2	0	0	1	0	0	3	0	1	0	3	0	7	0	3	1	5
0	3	1	7	0	3	2	1	0	3	2	3	0	3	3	7	0	3	4	3
0	0	3	6	0	0	1	0	0	2	0	7	0	2	2	5	0	2	3	1
0	2	5	7	0	2	6	1	0	2	6	3	0	2	6	5	0	2	7	7

36

35

377734

377735

377736

377737

0	0	3	5	0	0	1	0	0	1	2	5	0	1	2	7	0	1	3	1
0	1	5	5	0	1	5	7	0	2	0	1	0	2	0	3	0	2	0	5
0	0	2	2	0	0	1	2	0	0	0	3	0	0	0	4	0	0	3	7
0	0	4	1	0	0	5	7	0	0	6	3	0	1	0	1	0	1	2	3

22

 21 } First Pair
of
Sample Chain

Designates RER
Table 01

Means file is sequential because
first RB LINK of first word pair
starts in the 2-byte.

The RER Table for
"System" is 00.

"System" is a random file.

377776

377777

0	0	0	2	0	0	0	4												

 1 } First Pair
of
"System" Chain

HYPOTHETICAL FNT/FST ENTRY
ASSOCIATED WITH RBT CHAIN ON

PREVIOUS PAGE

59	47	35	23	11	0
FILE NAME				TYPE	P.
EQP. CODE	LOCK	FIRST RBT PAIR	CURRENT RBT PAIR	CURRENT RBT ENTRY	SECTOR
	FET ADDRESS		DISPOSITION CODE	SEC. CODE	LAST CODE AND STATUS

The designations of the FNT/FST field contents are repeated here for convenience. See pages ff. for details.

Assume that we have rewound the file after writing it and then have read up to RB LINK 0337, SECTOR 52. (This is in word pair 36, byte 6.) The FNT/FST entry might look like that below:

T		A		P		E		1											
2	4	0	1	2	0	0	5	3	4	0	0	0	0	3	2	0	0	0	0
0	2	0	0	0	0	2	1	0	0	3	6	0	0	3	6	0	0	5	2
0	0	0	0	2	2	7	6	0	0	0	0	0	2	0	0	0	0	1	1

TAPE 1

From the FNT/FST entry, we can get the following information:

1. The file name is TAPE 1.
2. The lock bit is not set.
3. The type is local (03).
4. The control point is 2.
5. There is no priority associated with the file at this time.
6. The equipment code specifies 6638 Disk, alternate sector half-tracks (02).
7. The first RBT pair is the 21st in the table.

8. The current RET pair (where file is currently positioned) is 36B.
9. The current RET entry is 3, i.e., there are 3 completely full RET word pairs up to (preceding) the current pair in which the file is positioned (the current position is in word pair array 36B).
10. The current byte position in RET pair 36 is 6.
11. The sector number is 52.
12. The address of the start of the File Environment Table (FET) associated with this file is 2276.
13. There is no disposition code associated with this file.
14. The security code is 02.
15. The last code and status is 11, which says the last operation completed was a read.

Pertinent Disk Tables

The FNT/FST area, RBT (Record Block Table) word pair area, and RBR (Record Block Reservation) tables play important roles in the disk reservation scheme. Low core contains pointers to and information related to these tables.

On the diagrams on the following pages:

A = Pointer to the start of the FNT/FST area.

B = Pointer to the LWA+1 of the FNT/FST area.

C = Length of RBT area/100B. The address of the first word of the RBT area = $400\ 000B - C * 100B$.

D = Start of RBR tables area (=B in the current LADC system).

E = Start of the empty chain in the RBT word pairs area (ie., the ordinal of the first word pair of the empty chain). The address associated with the first word of this word pair is $400\ 000B - 2 * E$.

F = The contents of the second byte of the second word of an FNT/FST entry. F is the ordinal of the first RBT word pair for that file.

P - Refers to a 3 word FNT/FST entry.

Q - Refers to a 38 word RBR table.

R - Refers to an RBT word pair.

Program Disk

Purpose

To perform on-line dynamic verification of the SCOPE 3.1 Disk Linkage. Reservation tables or the corresponding tables that are placed in an image area after each dead start by the bootstrap programs L99 and B00.

Method

This program prints out 20 words of low core, the FNT/FST entries, the RBT area, RBR tables, and the RBT empty chain. It then checks to be sure that there is a bit set in the appropriate RBR Table corresponding to each RBR link in the RBT area. An error message is printed out for each discrepancy found. Each FNT/FST entry is printed out followed by the proper disk number and all of the RB links associated with that file. The final result of the test is displayed on the scope and also goes to the dayfile.

Usage

1. To verify the current DISK Tables - DISK.
2. To verify the IMAGE DISK Tables - DISK (IMAGE)

Restrictions

The CM routines MIORC, MIORF, MBYTE, LBYT, MEMGET, and MEMPUT, and the PP routine MIO are used.

The program currently assumes two disks (2 RBR tables).

LOW CORE AREA

000000	0000	0000	0000	0000	0000
000001	0000	6700	0003	1036	0000
000002	0000	4420	0027	0003	4000
000003	0560	0000	0000	0000	1206
000004	2140	4420	0000	0000	0000
000005	2040	2140	0000	0000	0000
000006	0000	0000	0000	0000	0003
000007	5340	0000	0000	0000	0000
000010	0000	0000	0000	0000	0000
000011	0000	0000	0000	0000	0000
000012	0000	0000	0000	0000	0000
000013	0276	0000	2262	0002	0454
000014	0000	0000	0011	0000	0000
000015	0000	0000	0000	0001	0000
000016	0000	0000	0000	0000	0000
000017	0000	0000	0000	0000	0000

FNT/FST AREA

002140	0401 3106 1114 0530 0000	DAYFILEX
002141	0000 0041 0041 0002 0007	6 6 B G
002142	0000 0000 0000 0000 0016	N
002143	2331 2324 0515 0020 0000	SYSTEM P
002144	0200 0001 0001 0004 0000	B A A D
002145	0000 0626 0000 0000 0053	FV S
002146	0406 1114 0534 0030 0000	DFILE1 X
002147	0000 0114 0114 0002 0005	ALAL B E
002150	0000 0000 0000 0000 0017	O
002151	0406 1114 0535 0030 0000	DFILE2 X
002152	0000 0000 0000 0002 0000	B
002153	0000 0010 0000 0074 1073	H <H>
002154	0406 1114 0536 0030 0000	DFILE3 X
002155	0000 0000 0000 0002 0000	B
002156	0000 0010 0000 0074 1073	H <H>
002157	0406 1114 0537 0030 0000	DFILE4 X
002160	0000 0026 0026 0002 0001	V V B A
002161	0000 0010 0000 0000 0017	H O
002162	0406 1114 0540 0030 0000	DFILE5 X
002163	0000 0000 0000 0002 0000	B
002164	0000 0010 0000 0074 1073	H <H>
002165	0406 1114 0541 0030 0000	DFILE6 X
002166	0000 0000 0000 0002 0000	B
002167	0000 0010 0000 0074 1073	H <H>
002170	0406 1114 0542 0030 0000	DFILE7 X
002171	0000 0000 0000 0000 0000	
002172	0000 0000 0000 0000 0001	A
002173	2205 0104 0100 0031 0000	READA Y
002174	6000 0000 0000 0000 0000	=
002175	0000 0010 0000 0000 0031	H Y
002176	2205 0104 0300 0031 0000	READC Y
002177	6000 0000 0000 0000 0000	=
002200	0000 0030 0000 0000 0031	X Y
002201	1116 2025 2400 0034 0000	INPUT 1
002202	0200 0021 0021 0002 0053	B Q Q B S
002203	0000 0100 0000 0000 0021	A Q
002204	0130 3132 0000 0020 0000	AXYZ P
002205	0200 0020 0020 0003 0060	B P P C =
002206	0000 2112 0000 0200 0025	QJ B U
002207	1407 1700 0000 0034 0000	LGO 1
002210	0200 0022 0022 0002 0003	B R R B C
002211	0000 0626 0000 0200 0023	FV B S
002212	2022 1116 2403 0011 0003	PRINTC I C
002213	0200 0025 0030 0011 0026	B U X I V

002215	1725	2426	2524	0034	0000
002216	0200	0023	0023	0004	0025
002217	0000	0114	0000	0200	0025
002220	0315	2023	0322	0034	0000
002221	0000	0000	0000	0000	0000
002222	0000	3022	0000	0000	0051
002223	2323	2323	2323	2634	0000
002224	0000	0024	0024	0002	0001
002225	0000	0100	0000	0000	0021
002226	2323	2323	2323	3034	0000
002227	0000	0000	0000	0000	0000
002230	0000	0635	0000	0000	0053
002234	2323	2323	2323	2420	0000
002235	0000	0034	0034	0004	0056
002236	0000	0626	0000	0000	0027
002237	2323	2323	2323	2520	0000
002240	0000	0040	0040	0006	0057
002241	0000	0626	0000	0000	0027

OUTPUT 1
B S S D U
AL B U

CMPSCR 1

XR (

SSSSSSV1
T T B A
A Q

SSSSSSX1

F2 S

SSSSSSTP
1 1 D
FV W

SSSSSSUP
5 5 F
FV W

CURRENT LENGTH OF RRT AREA = N = 0300B

377500	0000 0000 0022 0023 0000	0140
377501	0000 0000 0000 0000 0000	
377502	0140 0000 0022 0023 0000	0137
377503	0000 0000 0000 0000 0000	
377504	0137 0000 0022 0023 0000	0136
377505	0000 0000 0000 0000 0000	
377506	0136 0000 0022 0023 0000	0135
377507	0000 0000 0000 0000 0000	
377510	0135 0000 0000 0000 0000	0134
377511	0000 0000 0000 0000 0000	
377512	0134 0000 0000 0000 0000	0133
377513	0000 0000 0000 0000 0000	
377514	0133 0000 0000 0000 0000	0132
377515	0012 0012 0012 0012 0012	
377516	0132 0000 0000 0000 0000	0131
377517	0002 0002 0002 0002 0000	
377520	0131 0000 0000 0000 0000	0130
377521	0002 0002 0002 0002 0002	
377522	0130 0000 0000 0000 0000	0127
377523	0002 0002 0002 0002 0002	
377524	0127 0000 0000 0000 0000	0126
377525	0000 0000 0000 0000 0000	
377526	0126 0000 0000 0000 0000	0125
377527	0002 0002 0002 0002 0002	
377530	0125 0000 0000 0000 0000	0124
377531	0002 0002 0002 0002 0002	
377532	0124 0000 0000 0000 0000	0123
377533	0002 0002 0000 0000 0000	
377534	0123 0000 0000 0000 0000	0122
377535	0002 0002 0002 0002 0002	
377536	0122 0000 0000 0000 0000	0121
377537	0002 0002 0002 0002 0002	
377540	0121 0000 0000 0000 0000	0120
377541	0002 0002 0002 0002 0002	
377542	0120 0000 0000 0000 0000	0117
377543	0002 0002 0002 0002 0002	
377544	0117 0000 0000 0000 0000	0116
377545	0012 0000 0000 0000 0000	

377546	0116 0000 0000 0000 0000	0115
377547	0002 0002 0002 0002 0002	
377550	0000 0002 0000 0005 0667	0114
377551	0000 0000 0000 0000 0000	
377552	0115 0000 0000 0000 0000	0113
377553	0000 0000 0000 0000 0000	
377554	0113 0000 0000 0000 0000	0112
377555	0012 0012 0012 0012 0012	
377556	0112 0000 0000 0000 0000	0111
377557	0012 0012 0012 0012 0012	
377560	0111 0000 0000 0000 0000	0110
377561	0012 0012 0012 0012 0012	
377562	0110 0000 0000 0000 0000	0107
377563	0012 0012 0012 0012 0012	
377564	0107 0000 0000 0000 0000	0106
377565	0012 0012 0012 0012 0012	
377566	0106 0000 0000 0000 0000	0105
377567	0012 0012 0012 0012 0012	
377570	0105 0000 0000 0000 0000	0104
377571	0002 0002 0002 0002 0002	
377572	0104 0000 0000 0000 0000	0103
377573	0000 0000 0000 0000 0000	
377574	0103 0000 0000 0000 0000	0102
377575	0012 0012 0012 0012 0012	
377576	0102 0000 0000 0000 0000	0101
377577	0012 0012 0012 0012 0012	
377600	0101 0000 0000 0000 0000	0100
377601	0012 0012 0012 0012 0012	
377602	0100 0000 0000 0000 0000	0077
377603	0012 0012 0012 0012 0012	
377604	0077 0000 0000 0000 0000	0076
377605	0012 0012 0012 0012 0012	
377606	0076 0000 0000 0000 0000	0075
377607	0000 0000 0000 0000 0000	
377610	0075 0000 0000 0000 0000	0074
377611	0002 0002 0000 0000 0000	
377612	0074 0000 0000 0000 0000	0073
377613	0002 0002 0002 0002 0002	
377614	0073 0000 0000 0000 0000	0072
377615	0002 0002 0002 0002 0002	
377616	0072 0000 0000 0000 0000	0071
377617	0002 0002 0002 0002 0002	
377620	0071 0000 0000 0000 0000	0070
377621	0002 0002 0002 0002 0002	

2
3-41

377622	0070 0000 0000 0000 0000	0057
377623	0012 0000 0000 0000 0000	
377624	0067 0000 0000 0000 0000	0066
377625	0002 0002 0002 0002 0002	
377626	0066 0000 0000 0000 0000	0065
377627	0002 0002 0002 0002 0002	
377630	0065 0000 0000 0000 0000	0064
377631	0002 0002 0002 0002 0002	
377632	0064 0000 0000 0000 0000	0063
377633	0002 0002 0002 0002 0002	
377634	0063 0000 0000 0000 0000	0062
377635	0002 0002 0002 0002 0002	
377636	0062 0000 0000 0000 0000	0061
377637	0002 0002 0002 0002 0002	
377640	0061 0000 0000 0000 0000	0060
377641	0002 0002 0002 0002 0002	
377642	0060 0000 0000 0000 0000	0057
377643	0002 0002 0002 0000 0000	
377644	0057 0000 0000 0000 0000	0056
377645	0002 0002 0002 0002 0002	
377646	0056 0000 0000 0000 0000	0055
377647	0000 0000 0000 0000 0000	
377650	0055 0000 0000 0000 0000	0054
377651	0000 0000 0000 0000 0000	
377652	0054 0000 0000 0000 0000	0053
377653	0002 0002 0002 0002 0002	
377654	0053 0000 0000 0000 0000	0052
377655	0000 0000 0000 0000 0000	
377656	0052 0000 0000 0000 0000	0051
377657	0000 0000 0000 0000 0000	
377660	0051 0000 0000 0000 0000	0050
377661	0012 0012 0012 0012 0012	
377662	0050 0000 0000 0000 0000	0047
377663	0012 0012 0012 0012 0012	
377664	0047 0000 0000 0000 0000	0046
377665	0012 0012 0000 0000 0000	
377666	0046 0000 0000 0000 0000	0045
377667	0012 0012 0012 0012 0012	
377670	0045 0000 0000 0000 0000	0044
377671	0012 0012 0012 0012 0012	
377672	0044 0000 0000 0000 0000	0043
377673	0012 0012 0012 0012 0012	
377674	0043 0000 0000 0000 0000	0042
377675	0000 0000 0000 0000 0000	

377676	0000 0002 0000 0007 0365	0041
377677	0000 0000 0000 0000 0000	
377700	0000 0014 0101 0057 0004	0040
377701	0000 0055 0071 0133 0000	
377702	0042 0000 0000 0000 0000	0037
377703	0012 0012 0012 0012 0012	
377704	0037 0000 0000 0000 0000	0036
377705	0000 0000 0000 0000 0000	
377706	0036 0000 0000 0000 0000	0035
377707	0012 0012 0012 0012 0012	
377710	0000 0012 0003 0056 0073	0034
377711	0045 0047 0000 0000 0000	
377712	0035 0000 0000 0000 0000	0033
377713	0012 0012 0012 0012 0012	
377714	0000 0010 0065 0063 0067	0032
377715	0000 0000 0000 0000 0000	
377716	0033 0000 0000 0000 0000	0031
377717	0000 0000 0000 0000 0000	
377720	0032 0010 0027 0033 0035	0030
377721	0041 0043 0051 0053 0057	
377722	0031 0000 0000 0000 0000	0027
377723	0000 0000 0000 0000 0000	
377724	0000 0002 0000 0001 0003	0026
377725	0000 0000 0000 0000 0000	
377726	0030 0012 0003 0051 0011	0025
377727	0015 0017 0021 0023 0025	
377730	0000 0012 0003 0001 0037	0024
377731	0000 0000 0000 0000 0000	
377732	0000 0012 0003 0025 0005	0023
377733	0007 0013 0000 0000 0000	
377734	0000 0012 0003 0004 0001	0022
377735	0000 0000 0000 0000 0000	
377736	0000 0012 0003 0053 0001	0021
377737	0000 0000 0000 0000 0000	
377740	0000 0012 0003 0060 0003	0020
377741	0031 0000 0000 0000 0000	
377742	0000 0000 0335 0337 0341	0017
377743	0343 0000 0000 0000 0000	
377744	0017 0000 0315 0317 0321	0016
377745	0323 0325 0327 0331 0333	
377746	0016 0000 0275 0277 0301	0015
377747	0303 0305 0307 0311 0313	
377750	0015 0000 0255 0257 0261	0014
377751	0263 0265 0267 0271 0273	

3-43

377752	0014	0000	0235	0237	0241	0013
377753	0243	0245	0247	0251	0253	
377754	0013	0000	0215	0217	0221	0012
377755	0223	0225	0227	0231	0233	
377756	0012	0000	0175	0177	0201	0011
377757	0203	0205	0207	0211	0213	
377760	0011	0000	0155	0157	0161	0010
377761	0163	0165	0167	0171	0173	
377762	0010	0000	0135	0137	0141	0007
377763	0143	0145	0147	0151	0153	
377764	0007	0000	0115	0117	0121	0006
377765	0123	0125	0127	0131	0133	
377766	0006	0000	0075	0077	0101	0005
377767	0103	0105	0107	0111	0113	
377770	0005	0000	0055	0057	0061	0004
377771	0063	0065	0067	0071	0073	
377772	0004	0000	0035	0037	0041	0003
377773	0043	0045	0047	0051	0053	
377774	0003	0000	0015	0017	0021	0002
377775	0023	0025	0027	0031	0033	
377776	0002	0004	0100	0024	0163	0001
377777	0024	0005	0007	0011	0013	

RBR TABLE 00

004420	0002	0100	0000	0301	0100
004421	4000	0001	4000	3614	0000

004422	7777	7777	7777	7777	7777
004423	7777	7777	7777	7777	6002
004424	0000	0000	0000	0000	0000
004425	0000	0000	0400	0000	0000
004426	0000	0000	0000	0000	0000
004427	0000	0000	0000	0000	0000
004430	0000	0000	0000	0000	0000
004431	0000	0000	0000	0000	0000
004432	0000	0000	0000	0000	0000
004433	0000	0000	0000	0000	0000
004434	0000	0000	0000	0000	0000
004435	0000	0000	0000	0000	0000
004436	0000	0000	0000	0000	0000
004437	0000	0000	0000	0000	0000
004440	0000	0000	0000	0000	0000
004441	0000	0000	0000	0000	0000
004442	0000	0000	0000	0000	0000
004443	0000	0000	0000	0000	0000
004444	0000	0000	0000	0000	0000
004445	0000	0000	0000	0000	0000
004446	0000	0000	0000	0000	0000
004447	0000	0000	0000	0000	0000
004450	0000	0000	0000	0000	0000
004451	0000	0000	0000	0000	0000
004452	0000	0000	0000	0000	0000
004453	0000	0000	0000	0000	0000
004454	0000	0000	0000	0000	0000
004455	0000	0000	0000	0000	0000
004456	0000	0000	0000	0000	0000
004457	0000	0000	0000	0000	0000
004460	0000	0000	0000	0000	0000
004461	0000	0000	0000	0000	0000

004462	7400	0000	0000	0000	0000
004463	0000	0000	0000	0000	0000
004464	0000	0000	0000	0000	0000
004465	0000	0000	0000	0000	0000

RBR TABLE 01

004466	0002	0200	0000	0301	0100
004467	4000	0002	4000	3741	0002

004470	7777	7777	7700	0004	0000
004471	0000	0000	0000	0000	0000
004472	0000	0000	0000	0000	0000
004473	0000	0000	0000	0000	0000
004474	0000	0000	0000	0000	0000
004475	0000	0000	0000	0000	0000
004476	0000	0000	0000	0000	0000
004477	0000	0000	0000	0000	0000
004500	0000	0000	0000	0000	0000
004501	0000	0000	0000	0000	0000
004502	0000	0000	0000	0000	0000
004503	0000	0000	0000	0000	0000
004504	0000	0000	0000	0000	0000
004505	0000	0000	0000	0000	0000
004506	0000	0000	0000	0000	0000
004507	0000	0000	0000	0000	0000
004510	0000	0000	0000	0000	0000
004511	0000	0000	0000	0000	0000
004512	0000	0000	0000	0000	0000
004513	0000	0000	0000	0000	0000
004514	0000	0000	0000	0000	0000
004515	0000	0000	0000	0000	0000
004516	0000	0000	0000	0000	0000
004517	0000	0000	0000	0000	0000
004520	0000	0000	0000	0000	0000
004521	0000	0000	0000	0000	0000
004522	0000	0000	0000	0000	0000
004523	0000	0000	0000	0000	0000
004524	0000	0000	0000	0000	0000
004525	0000	0000	0000	0000	0000
004526	0000	0000	0000	0000	0000
004527	0000	0000	0000	0000	0000

004530	0000	0000	0000	0000	0000
004531	0000	0000	0000	0000	0000
004532	0000	0000	0000	0000	0000
004533	0000	0000	0000	0000	0000

002140 0401 3106 1114 0530 0000 DAYFILEX
 002141 0000 0041 0041 0002 0007 6 6 B G
 002142 0000 0000 0000 0000 0016 N

THE FOLLOWING RRS ARE FOR DISK 00

0365

002143 2331 2324 0515 0020 0000 SYSTEM P
 002144 0200 0001 0001 0004 0000 B A A D
 002145 0000 0624 0000 0000 0053 FV S

THE FOLLOWING RRS ARE FOR DISK 00

0005 0007 0011 0013 0015 0017 0021 0023 0025 0027 0031 0033 0035 0037 0041 0043 0045 0047 0051 0053
 0055 0057 0061 0063 0065 0067 0071 0073 0075 0077 0101 0103 0105 0107 0111 0113 0115 0117 0121 0123
 0125 0127 0131 0133 0135 0137 0141 0143 0145 0147 0151 0153 0155 0157 0161 0163 0165 0167 0171 0173
 0175 0177 0201 0203 0205 0207 0211 0213 0215 0217 0221 0223 0225 0227 0231 0233 0235 0237 0241 0243
 0245 0247 0251 0253 0255 0257 0261 0263 0265 0267 0271 0273 0275 0277 0301 0303 0305 0307 0311 0313
 0315 0317 0321 0323 0325 0327 0331 0333 0335 0337 0341 0343

002146 0406 1114 0534 0030 0000 DFILE1 X
 002147 0000 0114 0114 0002 0005 ALAL B E
 002150 0000 0000 0000 0000 0017 O

THE FOLLOWING RRS ARE FOR DISK 00

0667

002151 0406 1114 0535 0030 0000 DFILE2 X
 002152 0000 0000 0000 0002 0000 B
 002153 0000 0010 0000 0074 1073 H SH>

THERE ARE NO RRT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002154 0406 1114 0536 0030 0000 DFILE3 X
 002155 0000 0000 0000 0002 0000 B
 002156 0000 0010 0000 0074 1073 H SH>

THERE ARE NO RRT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002157 0406 1114 0537 0030 0000 DFILE4 X
 002160 0000 0026 0026 0002 0001 V V B A
 002161 0000 0010 0000 0000 0017 H O

THE FOLLOWING RRS ARE FOR DISK 00

3-47

3-48

002162	0406 1114 0540 0030 0000	DFILE5 X
002163	0000 0000 0000 0002 0000	B
002164	0000 0010 0000 0074 1073	H SH>

THERE ARE NO RBT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002165	0406 1114 0541 0030 0000	DFILE6 X
002166	0000 0000 0000 0002 0000	B
002167	0000 0010 0000 0074 1073	H SH>

THERE ARE NO RBT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002170	0406 1114 0542 0030 0000	DFILE7 X
002171	0000 0000 0000 0000 0000	
002172	0000 0000 0000 0000 0001	A

THERE ARE NO RBT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002173	2205 0104 0100 0031 0000	READA Y
002174	6000 0000 0000 0000 0000	E
002175	0000 0010 0000 0000 0031	H Y

THIS FNT/FST ENTRY IS FOR A NON-ALLOCATABLE DEVICE

THERE ARE NO RBT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002176	2205 0104 0300 0031 0000	READC Y
002177	6000 0000 0000 0000 0000	E
002200	0000 0030 0000 0000 0031	X Y

THIS FNT/FST ENTRY IS FOR A NON-ALLOCATABLE DEVICE

THERE ARE NO RBT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002201	1116 2025 2400 0034 0000	INPUT 1
002202	0200 0021 0021 0002 0053	B Q Q B S
002203	0000 0100 0000 0000 0021	A Q

THE FOLLOWING RBS ARE FOR DISK 01

0001

002204	0130 3132 0000 0020 0000	AXYZ P
002205	0200 0020 0020 0003 0060	B P P C E
002206	0000 2112 0000 0200 0025	QJ B U

THE FOLLOWING RBS ARE FOR DISK 01

0003 0031

002207	1407 1700 0000 0034 0000	LGO 1
002210	0200 0022 0022 0002 0003	B R R B C
002211	0000 0624 0000 0200 0023	FV B S

002212	2022 1116 2403 0011 0003	PRINTC I C
002213	0200 0025 0030 0011 0026	B U X I V
002214	0000 0030 0040 0200 0011	X 58 I

THE FOLLOWING RBS ARE FOR DISK 01

0011 0015 0017 0021 0023 0025 0027 0033 0035 0041 0043 0051 00

002215	1725 2420 2524 0034 0000	OUTPUT 1
002216	0200 0023 0023 0004 0025	B S S D U
002217	0000 0114 0000 0200 0025	AL B U

THE FOLLOWING RBS ARE FOR DISK 01

0005 0007 0013

002220	0315 2023 0322 0034 0000	CMPSCR 1
002221	0000 0000 0000 0000 0000	
002222	0000 3022 0000 0000 0051	XR (

THERE ARE NO RBT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002223	2323 2323 2323 2634 0000	SSSSSSV1
002224	0000 0024 0024 0002 0001	T T B A
002225	0000 0100 0000 0000 0021	A Q

THE FOLLOWING RBS ARE FOR DISK 01

0037

002226	2323 2323 2323 3034 0000	SSSSSSX1
002227	0000 0000 0000 0000 0000	
002230	0000 0635 0000 0000 0053	F2 S

THERE ARE NO RBT WORD PAIRS FOR THE ABOVE FNT/FST ENTRY

002234	2323 2323 2323 2420 0000	SSSSSSTP
002235	0000 0034 0034 0004 0056	1 1 D
002236	0000 0626 0000 0000 0027	FV W

THE FOLLOWING RBS ARE FOR DISK 01

0073 0045 0047

002237	2323 2323 2323 2520 0000	SSSSSSUP
002240	0000 0040 0040 0006 0057	5 5 F
002241	0000 0626 0000 0000 0027	FV W

THE FOLLOWING RBS ARE FOR DISK 01

0055 0071 0133

THE R81 PAIRS IN THE EMPTY CHAIN FOLLOW

377722	0031 0000 0000 0000 0000	0027
377723	0000 0000 0000 0000 0000	
377716	0033 0000 0000 0000 0000	0031
377717	0000 0000 0000 0000 0000	
377712	0035 0000 0000 0000 0000	0033
377713	0012 0012 0012 0012 0012	
377706	0036 0000 0000 0000 0000	0035
377707	0012 0012 0012 0012 0012	
377704	0037 0000 0000 0000 0000	0036
377705	0000 0000 0000 0000 0000	
377702	0042 0000 0000 0000 0000	0037
377703	0012 0012 0012 0012 0012	
377674	0043 0000 0000 0000 0000	0042
377675	0000 0000 0000 0000 0000	
377672	0044 0000 0000 0000 0000	0043
377673	0012 0012 0012 0012 0012	
377670	0045 0000 0000 0000 0000	0044
377671	0012 0012 0012 0012 0012	
377666	0046 0000 0000 0000 0000	0045
377667	0012 0012 0012 0012 0012	
377664	0047 0000 0000 0000 0000	0046
377665	0012 0012 0000 0000 0000	
377662	0050 0000 0000 0000 0000	0047
377663	0012 0012 0012 0012 0012	
377660	0051 0000 0000 0000 0000	0050
377661	0012 0012 0012 0012 0012	
377656	0052 0000 0000 0000 0000	0051
377657	0000 0000 0000 0000 0000	
377654	0053 0000 0000 0000 0000	0052
377655	0000 0000 0000 0000 0000	
377652	0054 0000 0000 0000 0000	0053
377653	0002 0002 0002 0002 0002	
377650	0055 0000 0000 0000 0000	0054
377651	0000 0000 0000 0000 0000	
377646	0056 0000 0000 0000 0000	0055
377647	0000 0000 0000 0000 0000	
377644	0057 0000 0000 0000 0000	0056
377645	0002 0002 0002 0002 0002	

377642	0060	0000	0000	0000	0000	0057
377643	0002	0002	0002	0000	0000	
377640	0061	0000	0000	0000	0000	0060
377641	0002	0002	0002	0002	0002	
377636	0062	0000	0000	0000	0000	0061
377637	0002	0002	0002	0002	0002	
377634	0063	0000	0000	0000	0000	0062
377635	0002	0002	0002	0002	0002	
377632	0064	0000	0000	0000	0000	0063
377633	0002	0002	0002	0002	0002	
377630	0065	0000	0000	0000	0000	0064
377631	0002	0002	0002	0002	0002	
377626	0066	0000	0000	0000	0000	0065
377627	0002	0002	0002	0002	0002	
377624	0067	0000	0000	0000	0000	0066
377625	0002	0002	0002	0002	0002	
377622	0070	0000	0000	0000	0000	0067
377623	0012	0000	0000	0000	0000	
377620	0071	0000	0000	0000	0000	0070
377621	0002	0002	0002	0002	0002	
377616	0072	0000	0000	0000	0000	0071
377617	0002	0002	0002	0002	0002	
377614	0073	0000	0000	0000	0000	0072
377615	0002	0002	0002	0002	0002	
377612	0074	0000	0000	0000	0000	0073
377613	0002	0002	0002	0002	0002	
377610	0075	0000	0000	0000	0000	0074
377611	0002	0002	0000	0000	0000	
377606	0076	0000	0000	0000	0000	0075
377607	0000	0000	0000	0000	0000	
377604	0077	0000	0000	0000	0000	0076
377605	0012	0012	0012	0012	0012	
377602	0100	0000	0000	0000	0000	0077
377603	0012	0012	0012	0012	0012	
377600	0101	0000	0000	0000	0000	0100
377601	0012	0012	0012	0012	0012	
377576	0102	0000	0000	0000	0000	0101
377577	0012	0012	0012	0012	0012	
377574	0103	0000	0000	0000	0000	0102
377575	0012	0012	0012	0012	0012	
377572	0104	0000	0000	0000	0000	0103
377573	0000	0000	0000	0000	0000	
377570	0105	0000	0000	0000	0000	0104
377571	0002	0000	0000	0002	0002	

377566	0106 0000 0000 0000 0000	0105
377567	0012 0012 0012 0012 0012	
377564	0107 0000 0000 0000 0000	0106
377565	0012 0012 0012 0012 0012	
377562	0110 0000 0000 0000 0000	0107
377563	0012 0012 0012 0012 0012	
377560	0111 0000 0000 0000 0000	0110
377561	0012 0012 0012 0012 0012	
377556	0112 0000 0000 0000 0000	0111
377557	0012 0012 0012 0012 0012	
377554	0113 0000 0000 0000 0000	0112
377555	0012 0012 0012 0012 0012	
377552	0115 0000 0000 0000 0000	0113
377553	0000 0000 0000 0000 0000	
377546	0116 0000 0000 0000 0000	0115
377547	0002 0002 0002 0002 0002	
377544	0117 0000 0000 0000 0000	0116
377545	0012 0000 0000 0000 0000	
377542	0120 0000 0000 0000 0000	0117
377543	0002 0002 0002 0002 0002	
377540	0121 0000 0000 0000 0000	0120
377541	0002 0002 0002 0002 0002	
377536	0122 0000 0000 0000 0000	0121
377537	0002 0002 0002 0002 0002	
377534	0123 0000 0000 0000 0000	0122
377535	0002 0002 0002 0002 0002	
377532	0124 0000 0000 0000 0000	0123
377533	0002 0002 0000 0000 0000	
377530	0125 0000 0000 0000 0000	0124
377531	0002 0002 0002 0002 0002	
377526	0126 0000 0000 0000 0000	0125
377527	0002 0002 0002 0002 0002	
377524	0127 0000 0000 0000 0000	0126
377525	0000 0000 0000 0000 0000	
377522	0130 0000 0000 0000 0000	0127
377523	0002 0002 0002 0002 0002	
377520	0131 0000 0000 0000 0000	0130
377521	0002 0002 0002 0002 0002	
377516	0132 0000 0000 0000 0000	0131
377517	0002 0002 0002 0002 0000	
377514	0133 0000 0000 0000 0000	0132
377515	0012 0012 0012 0012 0012	
377512	0134 0000 0000 0000 0000	0133

377510	0135	0000	0000	0000	0000	0134
377511	0000	0000	0000	0000	0000	
377506	0136	0000	0022	0023	0000	0135
377507	0000	0000	0000	0000	0000	
377504	0137	0000	0022	0023	0000	0136
377505	0000	0000	0000	0000	0000	
377502	0140	0000	0022	0023	0000	0137
377503	0000	0000	0000	0000	0000	
377500	0000	0000	0022	0023	0000	0140
377501	0000	0000	0000	0000	0000	

63
3-53

LOWING ARE THE RESULTS OF A TEST TO CHECK IF THERE IS AN RBR BIT SET FOR EACH RB INDICATED IN THE RBT WORD PAIRS TABLE.

R TEST INITIATED FOR FILE DAYFILEX
R TEST INITIATED FOR FILE SYSTEM P
R TEST INITIATED FOR FILE DFILE1 X
R TEST INITIATED FOR FILE DFILE4 X
R TEST INITIATED FOR FILE INPUT 1
R TEST INITIATED FOR FILE XYZ P
R TEST INITIATED FOR FILE LGO 1
R TEST INITIATED FOR FILE PRINTC I C
R TEST INITIATED FOR FILE OUTPUT 1
R TEST INITIATED FOR FILE SSSSSSV1
R TEST INITIATED FOR FILE SSSSSSTP
R TEST INITIATED FOR FILE SSSSSSUP

9 FILES HAD NO RBT WORD PAIRS ASSOCIATED WITH THEM

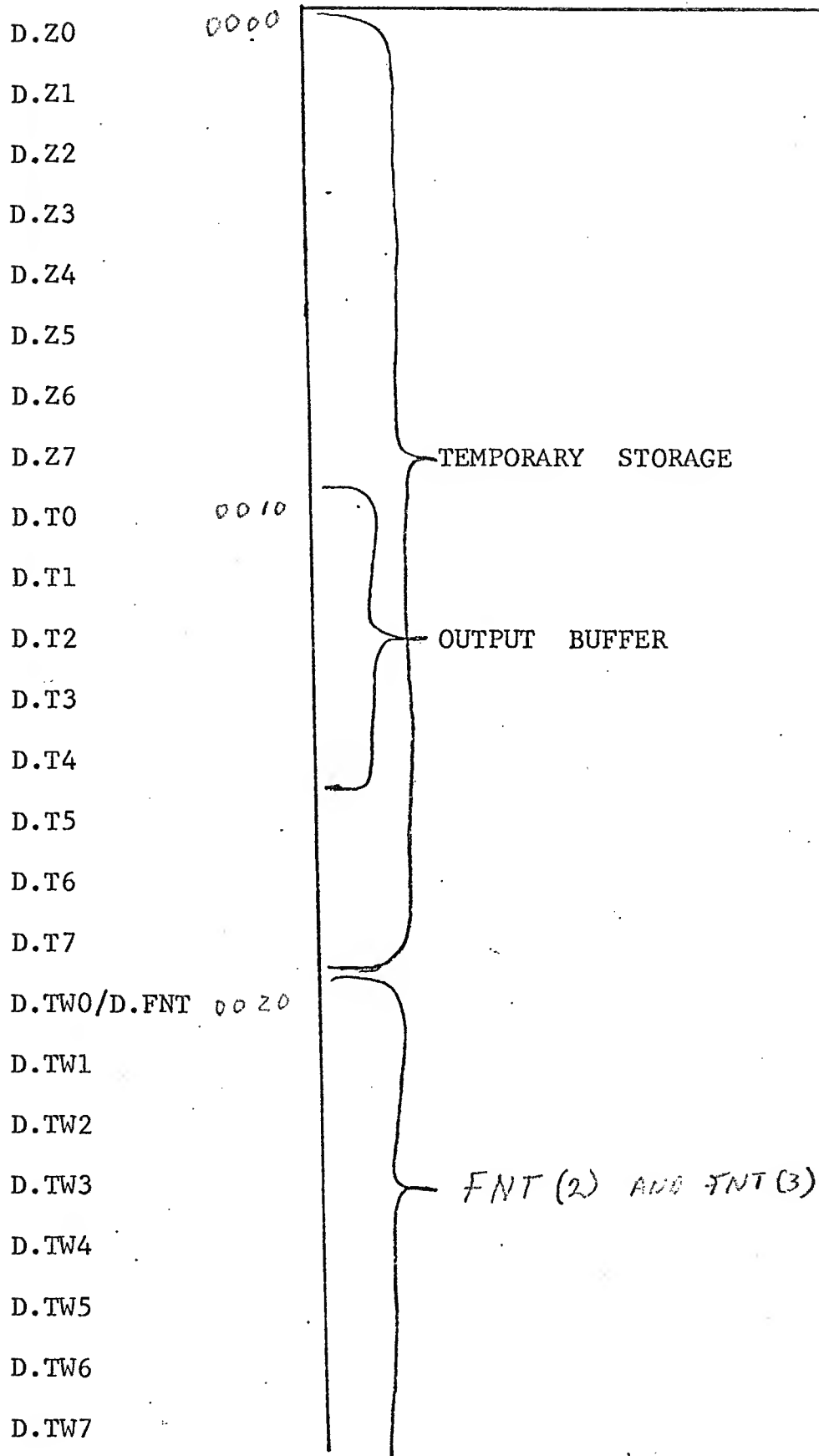
12 FILES TESTED FOR RB - RBR CORRESPONDENCE

THE NUMBER OF FILES IN THE FNT/FST AT THE TIME OF TESTING WAS 21

0 ERRORS IN RB - RBR CORRESPONDENCE TEST.

PP RESIDENT

4-1-1



D.TH0		
D.TH1		
D.TH2/D.EST 0032	}	EST
D.TH3		
D.TH4		
D.TH5		
D.TH6		
D.TH7/D.DTS/D.JFL 0037		
D.FR0/D.BA 0040	}	FET ADDRESS RELATIVE TO LOWER 18 BITS OF 2A.
D.FR1		
D.FR2		
D.FR3		
D.FR4		
D.FR5/D.JECS 0045		ECS FIELD LENGTH REQUIREMENT
D.FR6/D.JPR		COMPARE PRIORITY
D.FR7/D.JTL/ 10 ₈		JOB TIME LIMIT
D.FF0/D.PPIRB 0050	}	INPUT REGISTER BUFFER
D.FF1		
D.FF2		
D.FF3		
D.FF4		
D.FF5/D.RA 0055		CENTRAL MEMORY RA/100B
D.FF6/D.FL		CENTRAL MEMORY FL/100B
D.FF7/D.FA		FILE #ST ADDRESS (FNT(2) ADDRESS)

D.SX0/D.FIRST	0060	}	
D.SX1			
D.SX2/D.IN	0062		
D.SX3			
D.SX4/D.OUT	0064		
D.SX5			
D.SX6/D.LIMIT	0066		
D.SX7		}	
D.SV0/D.PPONE	0070		1
D.SV1/D.HN	0071		100B
D.SV2/D.TH	0072		1000B
D.SV3/D.TR	0073		3
D.SV4/D.CPAD	0074		CONTROL POINT ADDRESS
D.SV5/D.PPIR	0075		INPUT REGISTER ADDRESS
D.SV6/D.PPOR	0076		OUTPUT REGISTER ADDRESS
D.SV7/D.PPMES1	0077	MESSAGE BUFFER ADDRESS	

4-1-4

R.IDLE 0100

R.OVLJ 0111

R.OVL ~~0111~~ 124

IDLE LOOP

LOAD A PRIMARY OVERLAY & JUMP TO IT

LOAD OVERLAYS

R.EREQS 0300

ENTER A REQUEST IN STACK

R.WAIT 0410

WAIT FOR OUTPUT REG. CLEAR

R.PAUSE 0430

PAUSE FOR CM MOVE

R.MTR 0450 (B. PROCESS)

ISSUE A MONITOR REQUEST

R.READP 0460

DISK I/O ROUTINES

R.WRITEP 0470

R.STBMSK 0611

MASK FOR R.STB = 7700B

R.STB 0620

MASK BITS IN A LIST OF WORDS

4-1-5

R.CPFL 0627

CONTROL POINT FIELD LENGTH/100g

R.CPRA 0631

CONTROL POINT RELATIVE ADDRESS/100g

R.TFL 0634

CHECK ADDRESS FOR IN RANGE

R.DFM 0650

ISSUE DAYFILE MESSAGE

R.RCH 0704

REQUEST CHANNEL

R.DCH 0714

DROP CHANNEL

0772

1772

Rec'd
by
P. D. HSE

PP RESIDENT ENTRY POINTS

<u>Comfile Tag</u>	<u>3.0</u>	<u>3.1</u>	<u>Function</u>
R. IDLE	100	100	Idle loop
R.OVLJ	115	111	Enter to load overlay at 1000
R.OVL	124	124	Load PP overlay
R.EREQS	300	300	Enter request stack
R.WAIT	410	410	Wait output register clear
R.PAUSE	430	430	Pause for relocation
R.PROCESS (R.MTR)	450	450	Issue MTR function
R.READP	460	460	Xmit data to stack processor
R.WRITEP	470	470	Get data from stack processor
R.STBMSK	611	611	Channel table mask
R.STB	620	620	Process channel table mode
R.CPFL	627	627	FL
R.CPRA	631	631	RA
R.TFL	634	634	Add RA to ACC, test vs FL
R.DFM	650	650	Process dayfile message
R.RCH	704	704	Reserve channel
R.DCH	714	714	Drop channel
R.STEP	734	-	

100	Idle	(4)
117		
123	OVL	(4)
273		
277	EREQS	(3)
404		
407	WAIT	(1)
420		
421	PAUSE	
447	MTR	
457	READ/W	
503		(3)
506	STB	
525	TFL	(3)
544		
547	DFM	(3)
700		
703	RCH	
713	DCM	
714		(55)
771	MAIN	

PP RESIDENT

4-1-7

4-1-7

Resident Idle Loop

R.IDLE

Calling Sequence

LJM R.IDLE

R.IDLE is the idle loop; PP resident continually scans its input register for something to do.

Overlay Load

R.OVL

Calling Sequence

Load A register Load Address

RJM R.OVL

R.OVL causes an overlay whose name appears in D.T6 & D.T7 (right justified) to be loaded into the PP beginning at the address specified in the A register. R.OVL is used both by PP overlays to load higher level overlays & by PP resident to load the overlay named in the input register. PP resident does not reference the disk directly to load disk resident overlays but makes a call to the stack processor.

Enter Request Stack

R.EREQS

Calling Sequence

Store L (request) in D.T0

RJM R.EREQS

R.EREQS adds the control point number to the already formatted

4-2

request & searches the central memory request stack for an empty entry. The monitor function, M.EREQS, is called & PP resident iterates until the monitor accepts the request.

Wait

R.WAIT

Calling Sequence

RJM R.WAIT

R.WAIT will cause the PP to idle until the output register is clear.

Pause

R.PAUSE

Calling Sequence

RJM R.PAUSE

STD D.RA

R.PAUSE will exit if the PP is attached to control point zero or if the storage move flag is not set. Otherwise, the monitor function, M.PAUSE, will be issued & the PP will pause until monitor has completed the storage move for that control point. In any event, before an exit is made from R.PAUSE, the following information will be set:

(D.TO + C.CPST)	=	control point status
(D.TO + C.CPEF)	=	control point error flag
(D.TO + C.CPRA)	=	control point RA (hundreds)
(D.TO + C.CPFL)	=	control point FL (hundreds)
A register	=	control point RA

D.RA should always be reset after a jump to R.PAUSE.

Process Monitor Function

R.MTR

Calling Sequence

Store function parameters in
D.T1 to D.T4

Load function code

RJM R.MTR

R.MTR places the function code in D.T0, writes D.T0 through D.T4 to the output register, & waits for the output register to clear.

Transmit Data Via Channel
From/To Stack Processor

R.READP/R.WRITEP

Calling Sequence

Load L (request)

RJM R.READP/R.WRITEP

R.READP/R.WRITEP computes the PP word count from the first & last word addresses given in the already formatted request & adds the computed word count, the address of the PP message buffer, & the control point number to the request. The request is entered in the stack & data is transmitted via channel directly to/from PP memory. Upon exit from R.READP/R.WRITEP, the following information will be set:

(D.T3 + C.RWPPPLW)	=	LWA + 1 of data transmitted
(D.T3 + C.RWPPST)	=	status
(D.T3 + C.RWPPWT)	=	number of PP words transmitted

Calling Sequence

4-41

Load L (list)

RJM R.STB

where "list" has the form

L (byte)

L (word 1)

L (word 2)

⋮

L (word n)

zero

An entry point to R.STB called R.STBMSK is the address of the mask "ended" with each word in the list before the word is "exclusive orred" with the byte. This mask is initially 7700B and this value should be restored by any routine which substitutes an alternate mask. R.STB is used primarily to substitute channel numbers in driver overlays.

Test Field Length

R.TFL

Calling Sequence

Load relative address

RJM R.TFL

R.TFL is used to insure that a relative address is within the field length. The 18-bit address is added to the control point reference address (RA) & compared with the field length. If the address is out of range, R.TFL will exit with a negative A register; if the

address is legal, the A register will contain the absolute CM address (RA + relative address) upon exit. The control point RA & FL are kept locally within PP resident at R.CPRA & R.CPFL, respectively; these locations are initialized when an entry to R.PAUSE is made.

Enter Dayfile Message

R.DFM

Calling Sequence

Load L (message) + flag bits

RJM R.DFM

R.DFM will cause a message to be written from PP memory to the dayfile and/or the console. The flag bits are contained in the high order 6-bits of the A register upon entry to R.DFM & are used to determine the destination(s) of the message. Possible values of the flag bits are described below; one or more bits may be on; all are optional.

- 1 = Dayfile only ("A" display)
- 2 = control point 0 (system) message
- 4 = no "A" display

Drop Channel

R.DCH

Calling Sequence

Load channel number

RJM R.DCH

R.DCH will cause the specified channel to be dropped.

4-6

Request ChannelR.RCH

Calling Sequence

Load channel number

RJM R.RCH

The channel number(s) contained in the "A" register will be stored in byte D.T1, monitor function M.RCH inserted in D.T0, and D.T0 - D.T4 written to the output register for that PP. Channels will be assigned by MTR on the following priority basis:

D.T0	D.T1	D.T2	D.T3	D.T4
	2 1	4 3		

If alternate channels are specified MTR will stop looking for alternate channels upon sensing 6 bits of zero. Thus if desiring one alternate channel the programmer must clear D.T2 before entering R.RCH so the search will be terminated at that point. The procedure for requesting channel 12 with alternate channel 13 would be:

LDN	0
STD	D.T2
LDC	1312B
RJM	R.RCH

Monitor will stop looking for alternate channels after 4 channels have been investigated.

When R.RCH is used, D.T4 is automatically set nonzero; in this case, the function is not considered complete (i.e., output register is

4-7

not cleared) until a channel can be assigned. When complete, byte 0 of the output register is cleared & byte 4 is set to 7777B. A channel request may be made directly to the monitor (M.RCH); one other option is allowed in this case. If byte 4 of the output contains zero, the monitor will notify the requesting PP whether or not the channel could be assigned.

Peripheral Services Request Codes

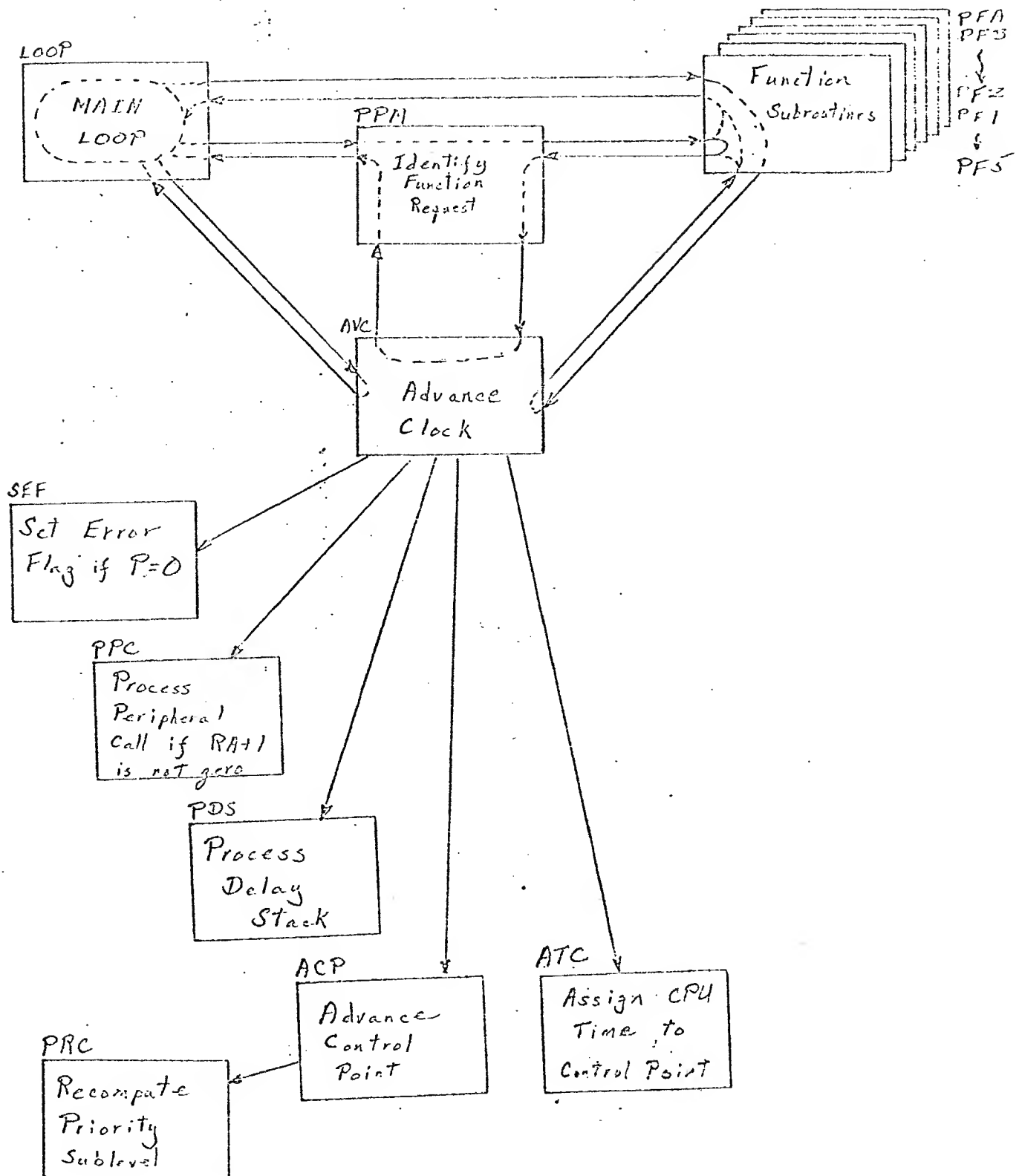
SCOPE 3.1

OPCODE	MNEMONIC	DESCRIPTION	OUTPUT REGISTER
01	M.DFM	Dayfile Message	0001 000X 0000 0000 0000 X=0 Normal X=1 Dayfile Only X=2 CP0 Message X=4 No A Display
02	M.RCH	Request Chanel	0002 BBAA CCDD 0000 ZZZZ ZZZZ = 2 Hang Till Free AA = Primary Choice BB = Second Choice DD = Third Choice CC = Fourth Choice
03	M.DCH	Drop Chanel	0003 00XX 0000 0000 0000
04	M.PPTIME	Assign Time	0004 0000 0000 0000 0000
05	M.STEP	STEP	0005 0000 0000 0000 000
06			
07			
10	M.RSTOR	Request Storage	0010 FFFF 0000 00XX 0000 XX=10 request RBT Storage(CP8) (BYTE 2 = Low Unit)
11	M.CDF	Complete Dayfile	0011 0000 0000 0000 0000
12	M.DPP	Drop PP	0012 0000 0000 0000 0000
13	M.ABORT	Abort	0013 0000 0000 0000 0000
14	M.NTIME	Time Limit	0014 TTTT 0000 0000 0000
15	M.RCP	Request CP	0015 0000 0000 0000 0000
16	M.DCP	Drop CP	0016 0000 0000 0000 0000
17	M.PAUSE	Pause	0017 0000 0000 0000 0000
20	M.RPP	Request PP	0020 0000 0000 0000 0000
21	M.RCLCP	Recall CP	0021 0000 0000 0000 0000
22	M.REQP	Request EQP	0022 EEEE 0000 0000 0000

23	M.DEQP	Drop EQP	0023 00EE 0000 0000 0000
24	M.RPRI	Priority	0024 00PP 0000 0000 0000
25	M.REM	Exit Mode	0025 000M 0000 0000 0000
26			
27			
30	M.OPDROP	Operator Drop	0030 00CP 00EE 0000 0000
31	M.RTAPE	ON	0031 00EE 0000 0000 0000
32	M.DTAPE	OFF	0032 00EE 0000 0000 0000
33	M.AEQP	Assign EQPT	0033 00EE 000P 0000 0000
34	M.EREQS	Enter Request	0034 000R 00AA SSSS 0000
35	M.CCPA	Change CP	0035 0000 0000 0000 00CP
36	M.CPUST	Change CPY	0036 0000 0000 0000 0000
37	M.RPJ	Delay PP Call	0037 MMMM MMMM bbbb 0000

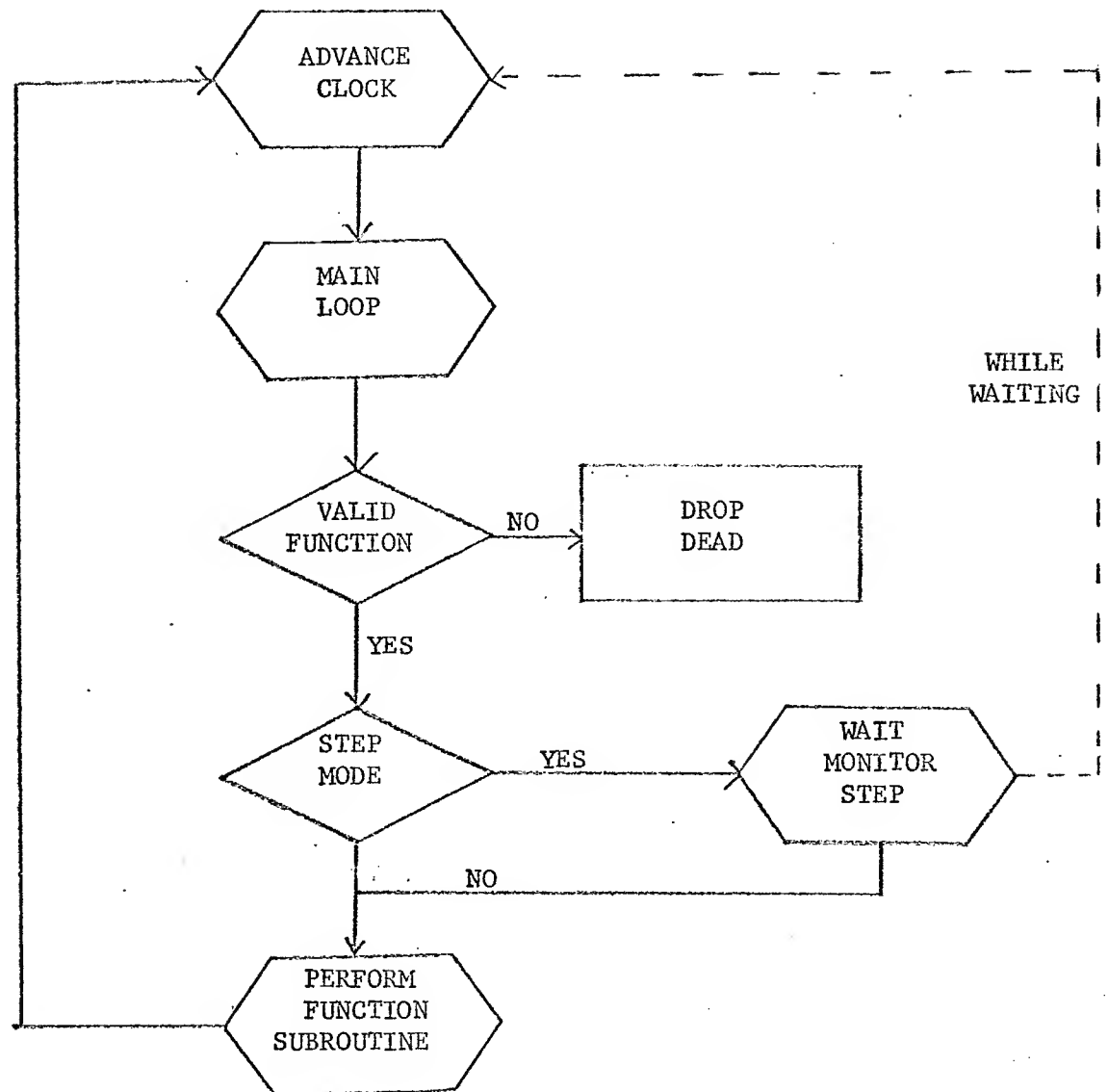
- MONITOR -

5-1

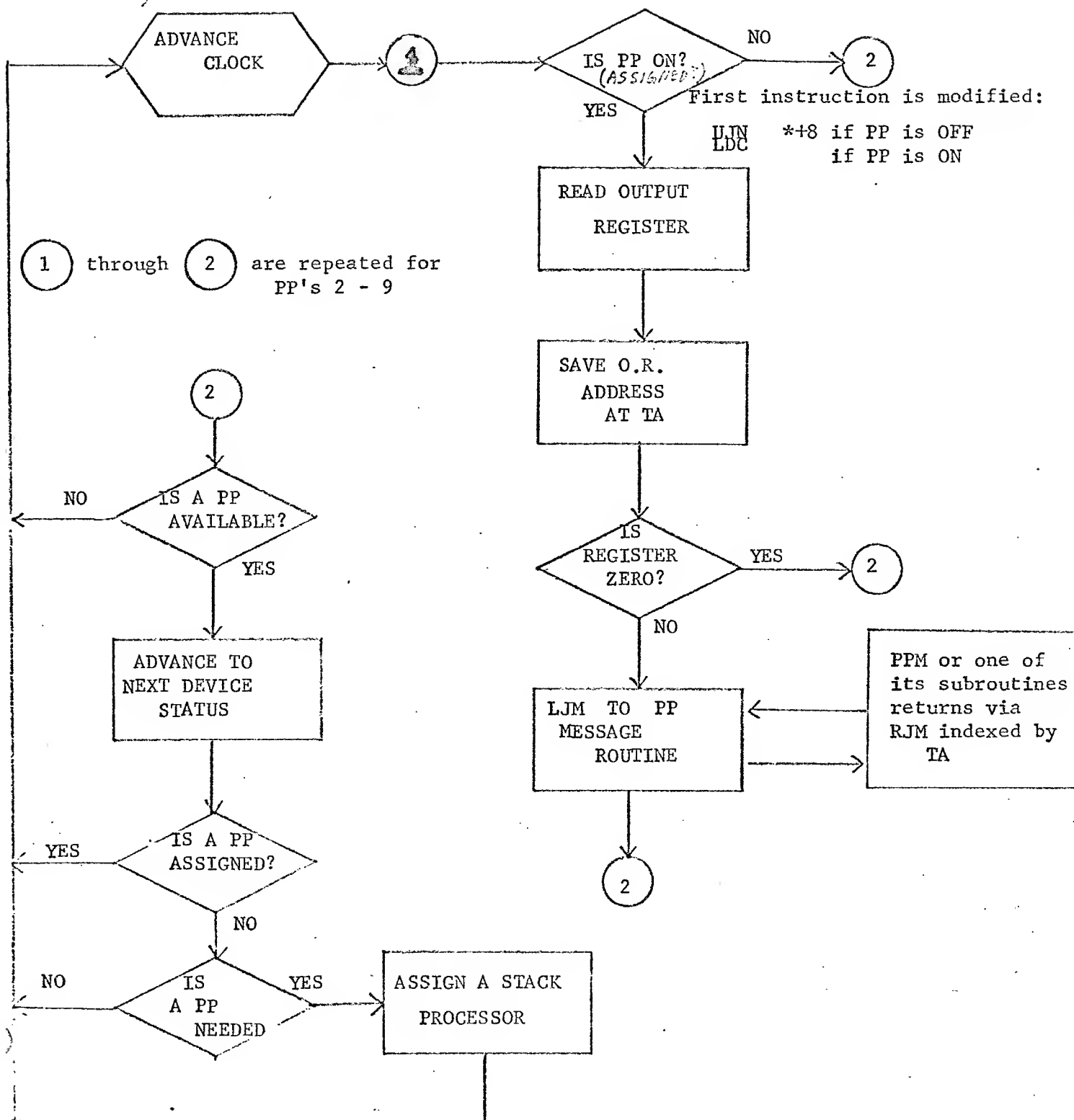


PERIPHERAL PROCESSOR MONITOR FUNCTION

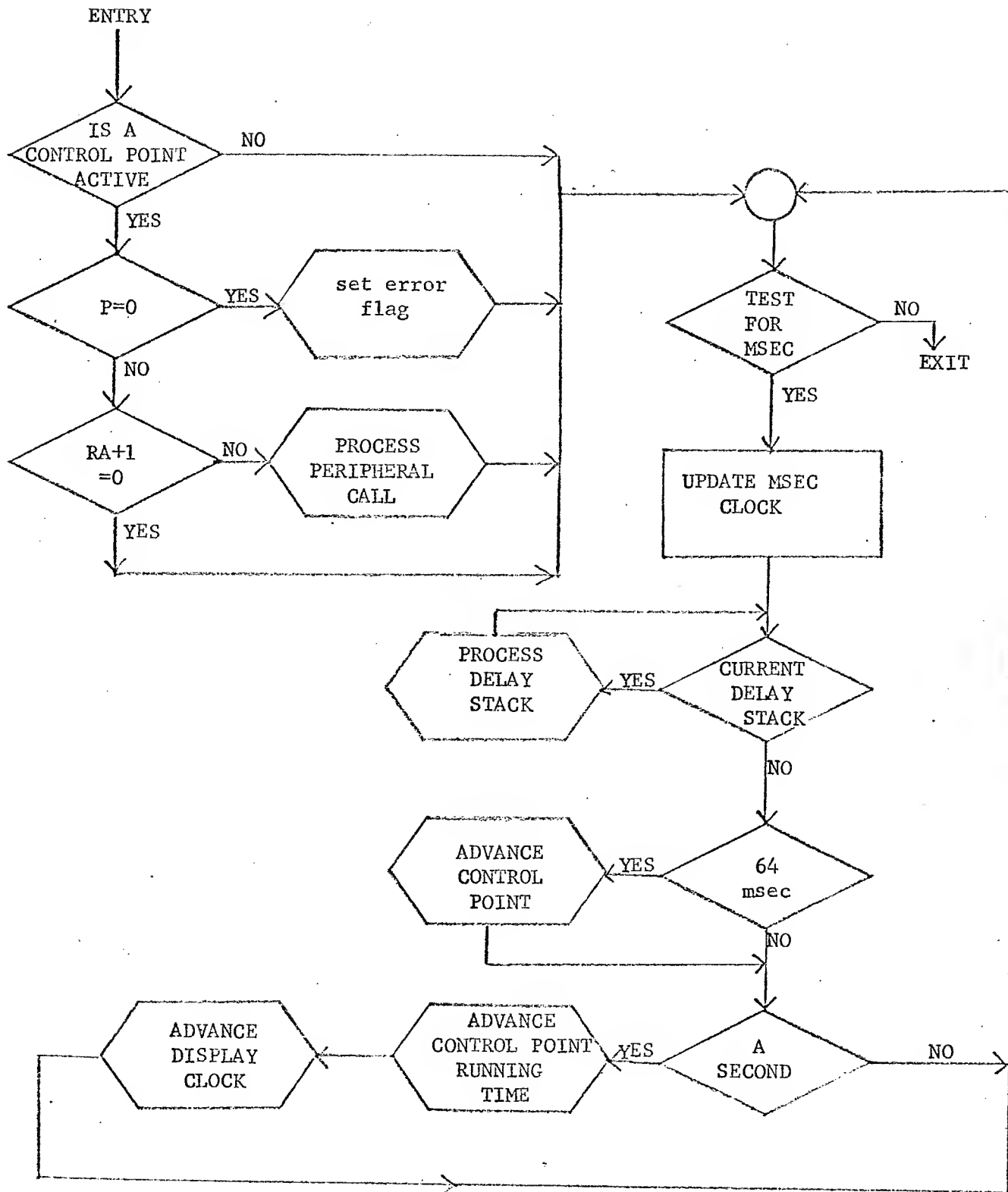
5-2



MTR Main Loop



ADVANCE CLOCK



CONCORDANCE OF 2.0/3.0 MONITOR FUNCTIONS

<u>Monitor Function</u>	<u>Mnemonic</u>	<u>Changed</u>	<u>Comments</u>
01	M.DFM	yes	(0 normal request (1 dayfile only ("A" display) - same as 2.0 byte 1 = (2 control point 0 (system) message (4 no "A" display flags may be combined - the upper 6 bits are used by DSD
02	M.RCH	yes	byte 4 = alternate command (may be 0) if channel is assigned byte 4 is set to 7777B hang till assignment achieved by setting byte 4 = 2
03	M.DCH	no	
04	M.PPTIME	no	
05	M.STEP	no	
06	assign track	gone	these functions are no-operations
07	release track	gone	
10	M.RSTOR	yes	also used to request ECS, RBT storage (0 request CM only (1 request ECS only byte 3 = (2 request both (10 request RBT storage (CP 8) (other request CM only for byte 3 = 1 or 2, byte 2 = ECS size in 512 word units for byte 3 = 10, byte 2 = requested low limit for RBT
11	M.CDF	no	now handled by DSD

12	M.DPP	no	MTR will check for auto recall ready
13	M.ABORT	no	
14	M.NTIME	no	
15	M.RCP	no	does not override auto recall
16	M.DCP	no	does not clear auto recall
17	M.PAUSE	no	
20	M.RPP	no	not satisfied until after all RPJs (see 37)
21	M.RCLCP	no	does not override auto recall
22	M.REQP	no	however, response byte moved in Control Point area
23	M.DEQP	no	
24	M.RPRI	no	7777B altered to 7776B
25	M.REM	no	
26	Set dayfile pointers	gone	these functions are no-operations
27	Toggle simulator	gone	
30	M.OPDROP	yes	byte 2 = value to set error flag (6 for drop)
31	M.SEF M.RTAPE	no	
32	M.DTAPE	no	

33	M.AEQP	no	
34	M.EREQS	new	enter words 0 and 1 of message buffer in request stack
35	M.CCPA	new	change control point assignment
			decrement stack count of given Control Point
36	unused		
37	M.RPJ	new	start peripheral job in word 0 of message buffer after given msec delay (may be 0) reply optional

Subject: SCOPE 3.1 Priority System

The priority of a job in the SCOPE 3.1 system is a twelve bit number. The twelve bits are divided into two fields. The high-order two bits are designated the priority level, the low-order ten bits are the sub-level.

The priority level is obtained from the priority field (P) of the JOB card. The four possible levels and their meaning are:

Level

- 0 Normal Priority
- 1 Normal Expedite (Requires signature of 101 or 595 supervision)
- 2 Special Expedite (Requires signature of Manager)
- 3 Panic Expedite (Requires signature of a Director or higher)

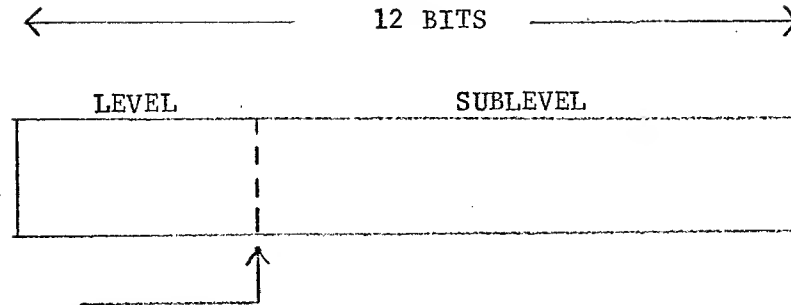
The ten-bit sub-level is initially determined from the time (T) and central memory (CM) estimates. Jobs in the input queue will periodically have their sub-levels increased. Thus the priority increases as a function of time until the job is assigned to a control point. Once a job is at a control point, its sub-level will be re-computed at regular intervals in an attempt to balance the utilization of the CPU between compute-bound and I/O-bound jobs. I/O jobs will tend to have their priorities increased; while the priorities of compute-bound jobs will tend to decrease. The priority sub-levels of jobs in the output queue (PRINT) will be inversely proportional to the amount of output to be printed.

It should be noted that the priority level will never change. That is, a priority 0 will never become a 1 or vice versa.

The following examples demonstrate the initial priority determination. All numbers are expressed in octal.

<u>Job Card Parameters</u>			<u>Priority</u>
<u>P</u>	<u>T</u>	<u>CM</u>	
0	100	00000	1400
0	600	00000	1000 (0000 + 1000)
1	600	00000	3000 (2000 + 1000)
2	600	00000	5000 (4000 + 1000)
0	100	12000	1100
0	600	15000	0400
0	100	40000	1700

PRIORITY



DETERMINED BY:

IP.MPR - THE MAXIMUM PRIORITY LEVEL ASSIGNABLE BY USER.

INSTALLATION OPTIONS AFFECTING PRIORITY (RE) EVALUATION

- IP.LV^FG - LEVEL ABOVE WHICH SUBLEVEL WILL REMAIN FIXED.
- IP.IQD - DETERMINES DELAY BETWEEN INCREMENTING SUB-
LEVEL OF A JOB IN THE INPUT QUEUE.
- IP.OQD - DETERMINES DELAY BETWEEN INCREMENTING SUB-
LEVEL OF A JOB IN THE OUTPUT QUEUE.
- IP.CPD - DETERMINES DELAY BETWEEN REEVALUATIONS OF
SUBLEVEL OF JOBS IN CP.
- IP.OSW - PARAMETER IN DYNAMIC REEVALUATION ALGORITHM
FOR WEIGHTING THE FORMER SUBLEVEL.

JOB CARD MACROS

WEIGHT	FIELD, RELATION, VALUE, ADDITIVE
OCTAL	FIELD
DECIMAL	FIELD

FIELD mnemonics:

CM = CENTRAL MEMORY REQUIREMENT

T = OVERALL JOB TIME LIMIT

EC = ECS MEMORY REQUIREMENT

RELATION mnemonics:

GE

LE

VALUE units:

VALUE_{CM} = 64 WORDS

VALUE_T = 8 SECONDS

VALUE_{EC} = 512 WORDS

INITIAL SUBLEVEL COMPUTATION EXAMPLES

INSTALLATION SPECIFIES:

DECIMAL T

WEIGHT T, LE, 8, 10B

WEIGHT CM, LE, 400B, 4

WEIGHT T, LE, 4, 20B

WEIGHT EC, GE, 100B, 100B

SAMPLE JOB CARDS:

JOB1, T40, CM30000.

SUBLEVEL = 10B + 4 = 14B

JOB2, T20, CM20000.

SUBLEVEL = 10B + 4 + 20B = 34B

JOB3, T8000, CM30000, EC200.

SUBLEVEL = 4 + 100B = 104B

DYNAMIC SUBLEVEL REEVALUATION FORMULA

$$S_{i+1} = \left[(2^{(IP.CPD+6)} - CPT) / ADJ + (2^{IP.OSW-1} - 1) S_i \right] / 2^{IP.OSW}$$

WHERE:

S = CURRENT PRIORITY SUBLEVEL

CPT = CENTRAL PROCESSOR TIME USED SINCE LAST EVALUATION

ADJ = ADJUSTMENT FACTOR TO NORMALIZE TIME UNITS

$$ADJ = 2^{(IP.CPD+6) - NBS}$$

WHERE:

NBS = THE NUMBER OF BITS IN THE SUBLEVEL.

PRIORITY SUBLEVEL REEVALUATION EXAMPLE

IF:

--IP.MPR = 63, THE PRIORITY LEVEL MAY RANGE FROM 1 to 77_8 AND
WILL OCCUPY THE HIGH ORDER 6 BITS OF THE PRIORITY
ITEM.

IP.LVF = 61, THE UPPER TWO PRIORITY LEVELS (76_8 AND 77_8) WILL
BE FIXED, AND NO RECOMPUTATION OF SUBLEVELS WILL TAKE
PLACE.

IP.CPD = 4, THE PRIORITY SUBLEVEL WILL BE RECOMPUTED EACH
 $2^{(IP.CPD+6)}$ MSEC, OR ABOUT ONCE A SECOND.

IP.OSW = 2, THE OLD SUBLEVEL WILL CARRY A WEIGHT OF
 $2^{(IP.OSW) - 1}$, OR 3.

ADJ WILL HAVE A VALUE OF $2^{(IP.CPD+6) - NBS} = 2^{(4+6-6)} = 16$.

THUS THE REEVALUATION COMPUTATION WOULD BE:

$$S_{i+1} = (1024 \text{ MS} - \text{CPT})/16 + 3S_i/4$$

Assume four jobs running at the same priority. If all four were demanding all of the CP time, each would receive 25% of the CP time. Job A is I/O bound and cannot use more than 10% of the CP time. That leaves 30% for each of the other three jobs. Job B is also I/O bound but can use 40% of the CP time. C and D can make use of all the CP time they can get. The following chart shows how these jobs would share the CP time. Within the chart the first number is the value of the priority sublevel and the second is the percentage of CP time actually used during the period. The example is in decimal for easier comprehension.

PRIORITY / % CP TIME USED

JOB	FIRST PERIOD	SECOND PERIOD	THIRD PERIOD	FOURTH PERIOD
A	50/10	60/0	70/10	75/10
B	50/40	52/0	64/40	63/40
C	50/50	50/0	62/50	59/50
D	50/0	62/100	46/0	59/0
	FIFTH PERIOD	SIXTH PERIOD	SEVENTH PERIOD	EIGHTH PERIOD
A	78/10	81/10	83/10	84/10
B	62/0	71/40	68/40	66/0
C	56/0	67/50	62/0	71/90
D	69/90	54/0	65/50	61/0
	NINTH PERIOD	TENTH PERIOD	ELEVENTH PERIOD	TWELFTH PERIOD
A	85/10	86/10	87/10	87/10
B	74/40	70/40	67/0	75/40
C	55/0	66/50	62/0	71/50
D	70/50	65/0	73/90	54/0
	THIRTEENTH PERIOD	FOURTEENTH PERIOD	FIFTEENTH PERIOD	TOTALS
A	87/10	87/10	87/10	9.3%
B	71/40	68/0	76/40	26.7%
C	65/50	61/0	70/50	32.7%
D	65/0	73/90	54/0	31.3%

5-16

STORAGE MOVE EXAMPLE - All figures /100_g

Control Point 5 Requests an FL of 300

<u>Control Point</u>	<u>Before</u>			<u>After</u>		
	<u>RA</u>	<u>FL</u>	<u>UAS</u>	<u>RA</u>	<u>FL</u>	<u>UAS</u>
0	0	142	0	0	142	0
1	142	33	0	142	33	0
2	175	31	0	175	31	0
3	226	0	500	226	0	0
4	726	20	130	226	20	0
5	1076	100	0	246	300	430
6	1176	150	0	1176	150	0
7	1346	0	430	1346	0	430

SYSTEM DISPLAY - DAYFILE PROCESSING⁶⁻¹¹⁻⁶⁻²

DSD OPERATOR ENTRY CHANGES

Deletions:

SIM.

DATE.

RECOVER.

n.EXPRESS.

Additions:

n.STEP. Places control point n in STEP mode. Only one control point may be in n.STEP at any one time. Former STEP entry still valid for system STEP mode.

n.LOADX. Loads jobs from non-3.0 tapes.

n.BLANK. Prepares degaused or blank tapes for use by SCOPE 3.0 by writing blank labels (double tape mark) at start of tape.

n.RECHECK. Rechecks tape labels following pause for operator action due to label error.

n.VRN,xxxxxx. Enters visual reel number up to 6 digits long following request for same by system.

Modifications:

n.DAYFILE,uu. Dumps system dayfile to equipment type specified by parameter uu. May be LP, CP, or MT.

Internal Changes:

n.DIS. May not be entered when system or a control point is in STEP mode.

The following type-ins will cause DSD to be temporarily assigned to the specified control point while the requested operation is being performed.

n.ENDx.

n.GO.

n.OFFSWx.

n.ONSWx.

n.RECHECK.

n.REPEATx.

n.SUPPRESSx.

Display Changes:

The format of the H display has been modified to display only selected files. The format may be altered by typing in H,xy. where x and/or y may be:

I -- INPUT

Ø -- OUTPUT

P -- PUNCH

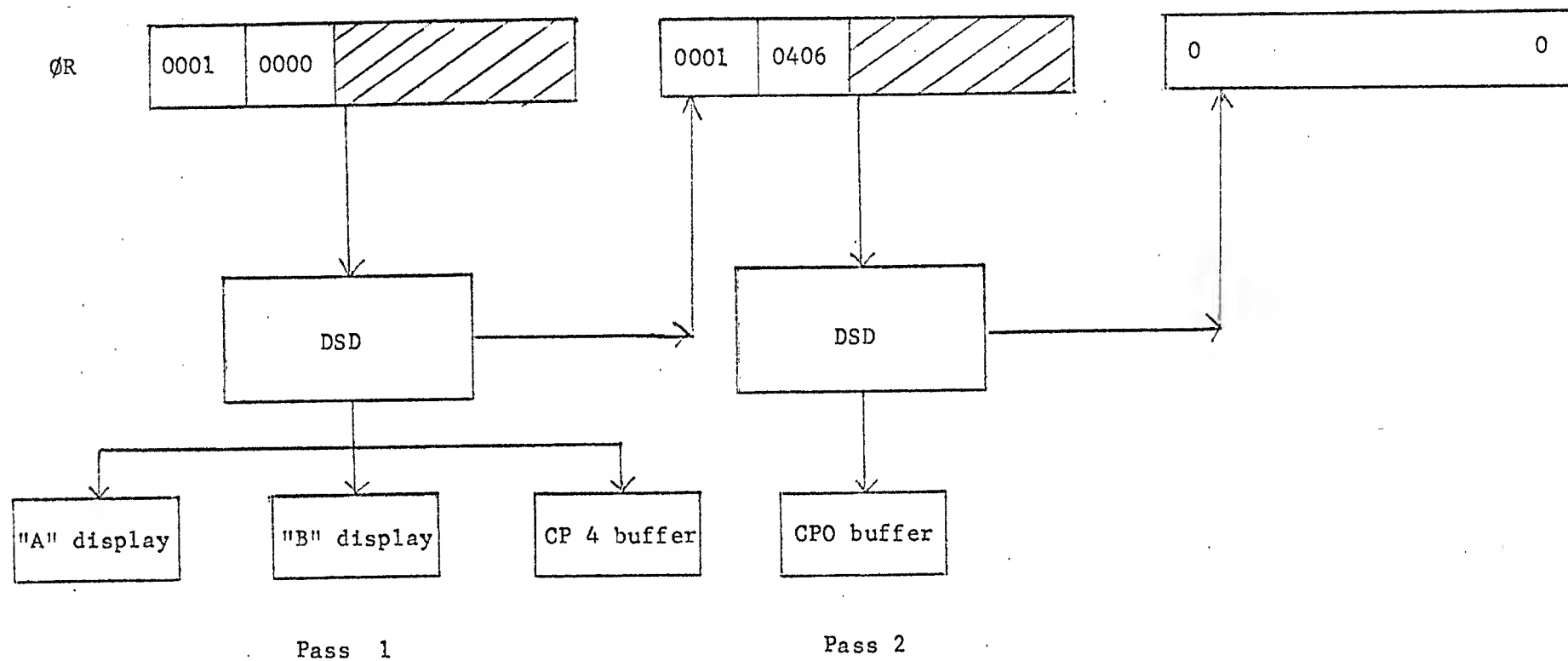
C -- COMMON

L -- LOCKED

FET AT CONTROL POINT 0

DAYFILE		STATUS/CODE
		FIRST
		IN
		OUT
FNT ADDRESS		LIMIT
		MAX

DAYFILE MESSAGE REQUEST PATH (MESSAGE FOR CP 4)



STACK PROCESSOR

7-10

Stack Processor Orders.

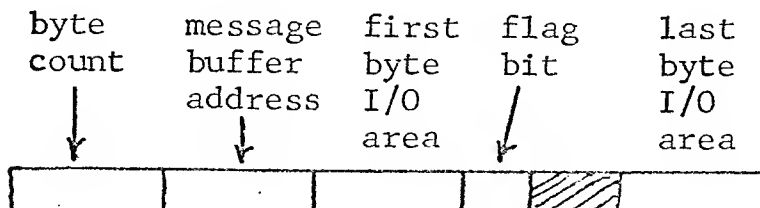
Note: These codes are communicated to the stack processor. They are not the codes used in the code/status word.

Order codes requiring a specific request format are indicated. In most cases the format is determined by flag bit settings rather than order code.

- O.READ (00): Read into central memory until a short PRU is encountered or the buffer is full(IN=OUT).
- O.RDSK (01): Read into central memory until a short PRU is encountered or until the buffer is full. Set the FST to reference the first PRU following the first end-of-record of level x or greater. The level is given in the high-order 6 bits of the order byte. (ie. 1401 would request a read, then would require positioning following an EOR of level 14 or greater)
- O.RGMPR (02): Read into central memory after dropping the first three CM words of the first PRU. This is used by STITCH for loading a program from the system library, eliminating the three word header added to system programs by EDITLIB.
- 03 : Undefined. Will result in an invalid stack entry message and may abort the control point.
- O.WRT (04): Write full PRU's from central memory.
- O.WRTR (05): Write from central memory, ending with a short PRU of the level specified in the high-order six bits of the order byte. If an EOF flag bit is found in this order, a zero length PRU of level 17 will be written following the short PRU terminating the record.

06,07 : Undefined.

- O.RDP (10) : Read into the requesting PP's memory until a short PRU is encountered, or until the input area is full. For this and other PP I/O orders the format of the second word of the request will be:



- O.RDPNP (11): Read into requesting PP after dropping first three cm words of the first PRU. This is used for all PP system program calls.
- O.SKP (12): Skip forward n records of level x or greater. The level is specified in the high six bits of the order byte; the number of records to be skipped is given in the third byte of the second word of the order. No data is transmitted.
- O.SKB (13): Skip backward n records of level x or greater. The level is specified in the high six bits of the order byte; the number of records to be skipped is given in the third byte of the second word of the order. No data is transmitted.
- O.WRP (14): Write from requesting PP, full PRU's only.
- O.WRPR (15): Write from requesting PP, ending with a short PRU of the level specified in the high order six bits of the order byte. If an EOF flag bit is set in this order, a zero length PRU of level 17 will be written following the short PRU terminating the record.
- UBPRU (16): Backspace n PRU's. The number of PRU's to be backspaced is given in the third byte of the second word of the order. Note that this requests repositioning defined by physical rather than logical units. No data is transmitted.
- O.RCHN (17): Release chain. All record blocks assigned to a file are released, and the RBT word pairs containing them are released. The FST is reset to an empty condition, if its address is supplied in the order.
Note: requests 16 and 17 require no communication with the device, and therefore are given the highest priority in the search for the next order to be executed. All other requests are assigned priority based on repositioning required.

REQUEST FOR PP TRANSACTION

WORD 1
IF FNT bit = 0

Address of FNT word				Order (specifies PP transaction)	Control Point Number	Physical Unit
RBT word	RBT ordinal	byte	PRU			
PP word count	Address of PP message word	PP FWA	R E C A L L	F N T	F E T	A V A I L A B L E
			PP LWA + 1			

WORD 1
IF FNT bit = 1
(NO FNT)

WORD 2

Release EOF
or
After EOR

REQUEST FOR CM Transaction

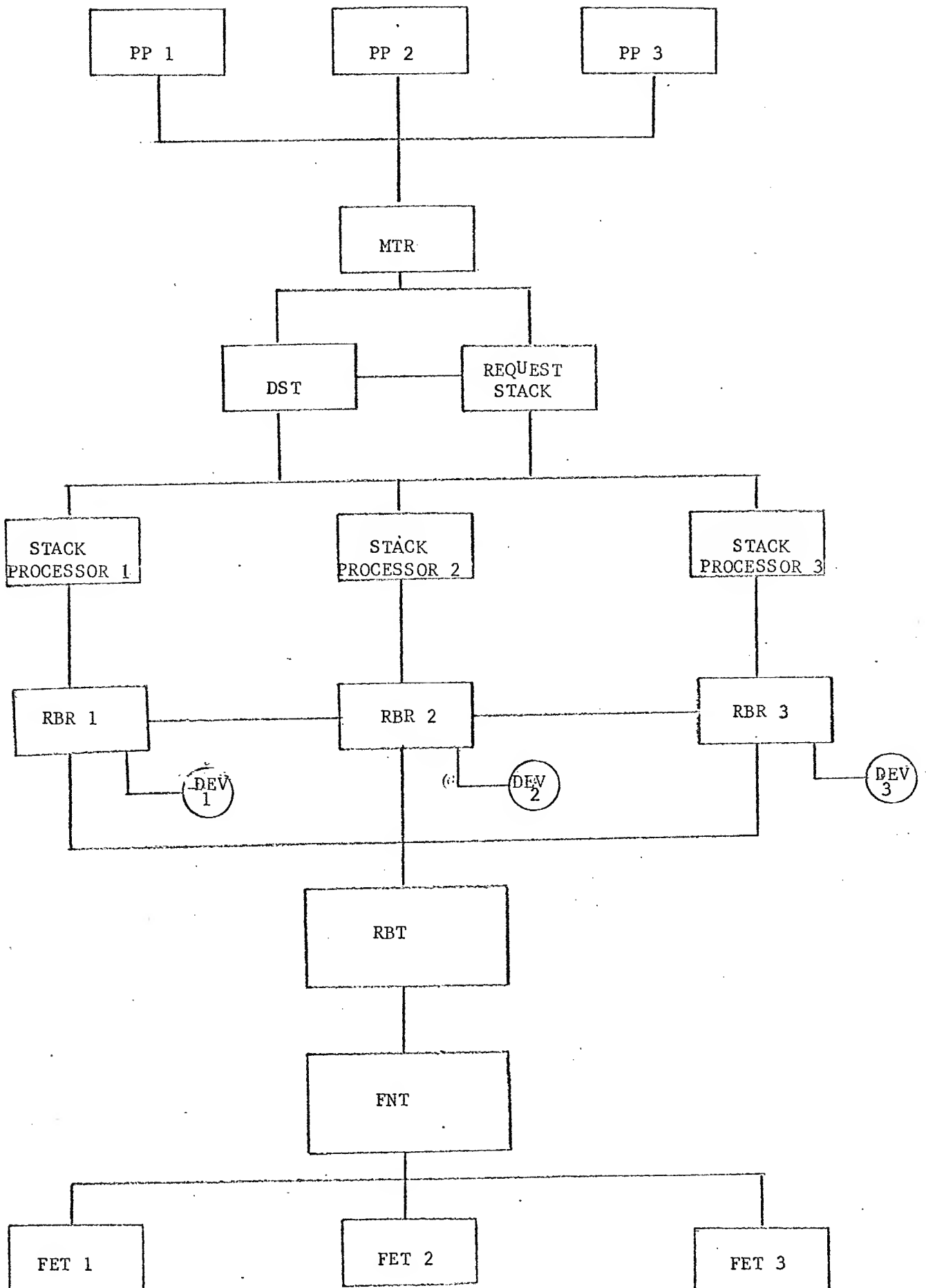
Word 1
if FNT bit=0

Word 1
if FNT bit=1
(no FNT)

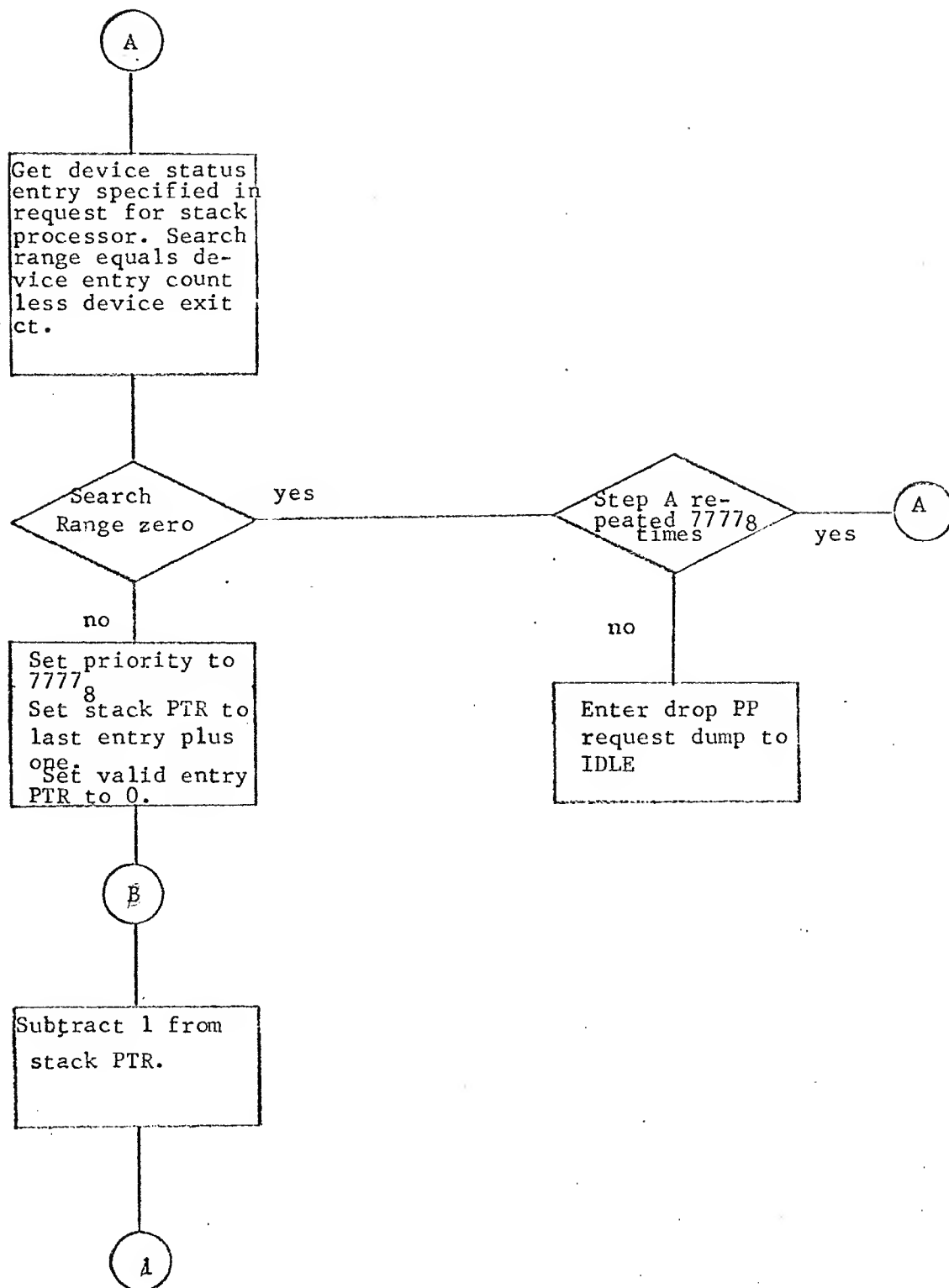
Word 2

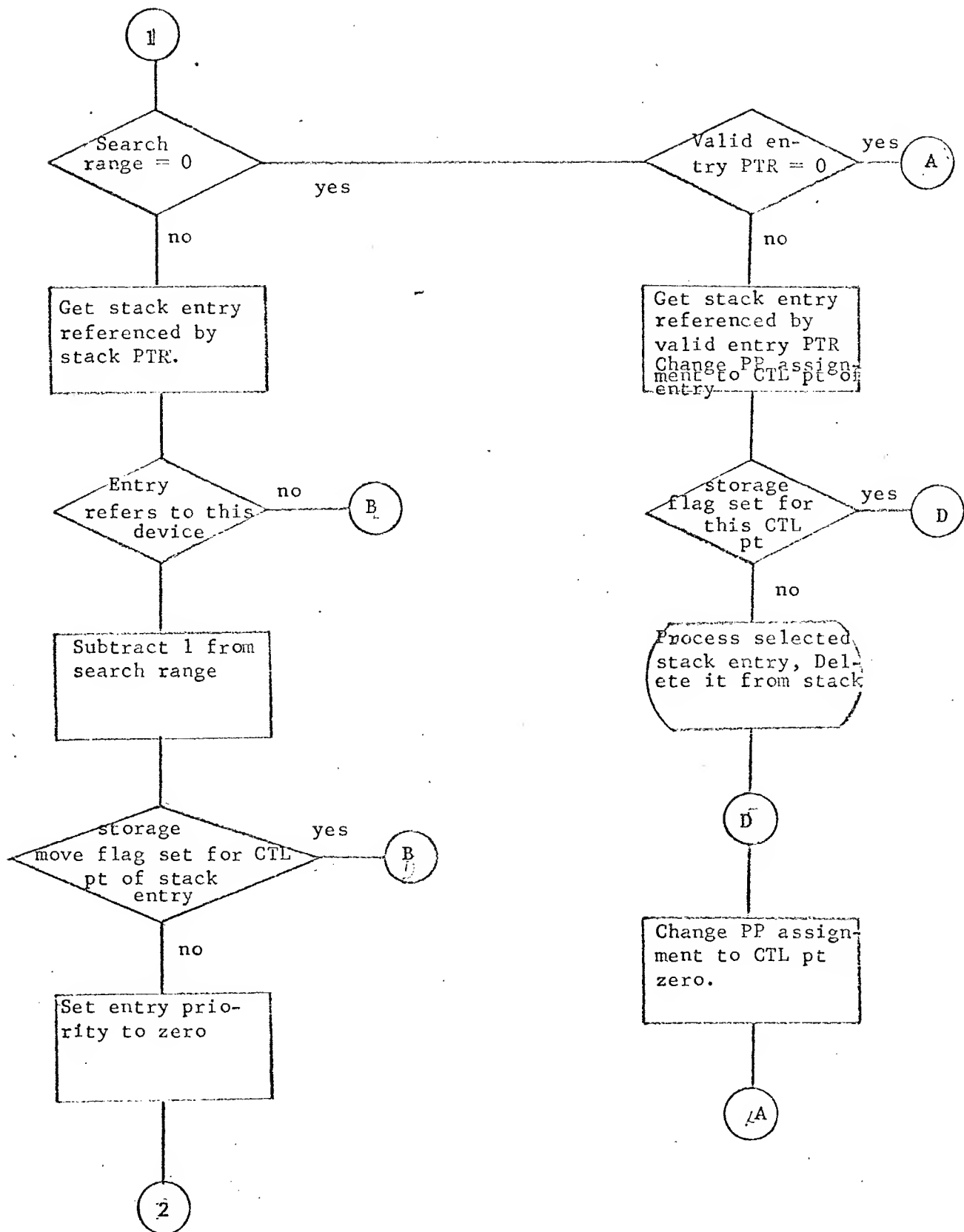
Address of FNT word				Order (specifies CM transaction)	Control Point Number	Physical Unit
RBT word	RBT ordinal	byte	PRU			
FET FWA	FIRST for transmission or RECORD (PRU) COUNT-n			R E C A L L	F F N E T T	A V A I L A B L E
				LAST for transmission		

Release EOF
or
After EOR

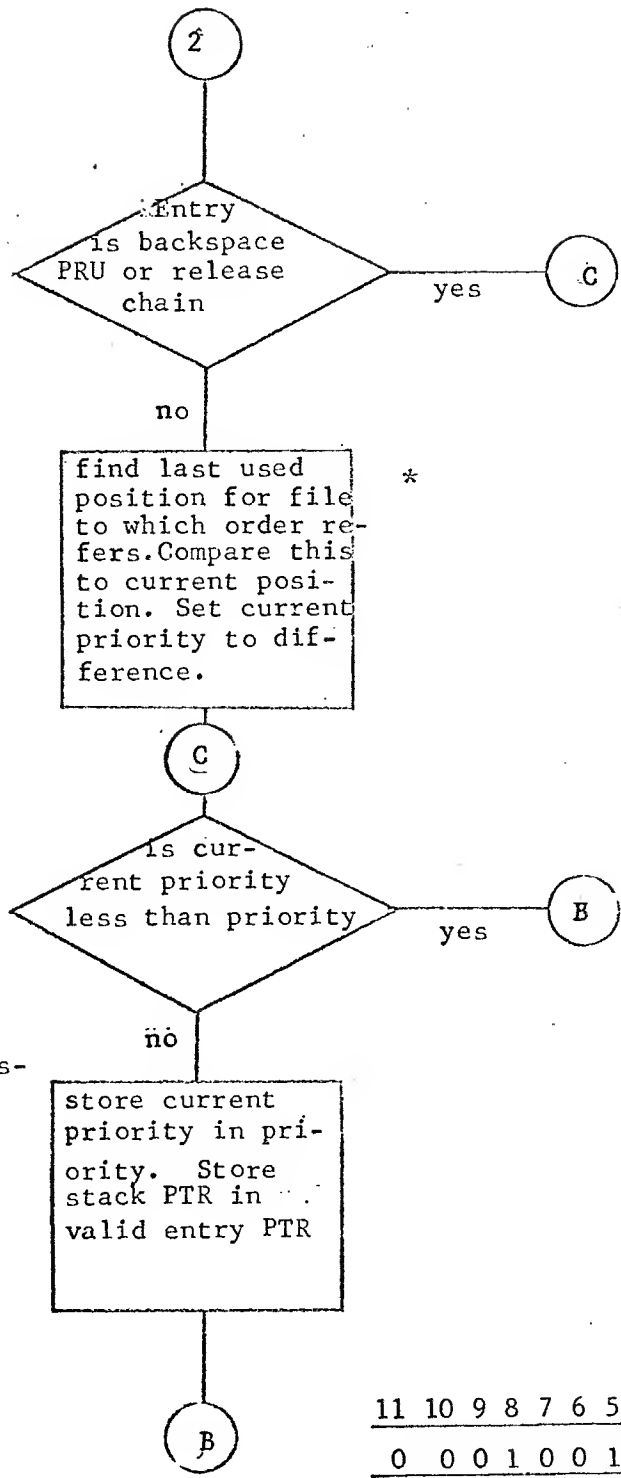


DISK STACK PROCESSING





Note: Backspace PRU and release chain orders do not reference disk and therefore require no positioning.



* Current position and last-used position are both used in the format shown below (Fig A). By use of logical minus, and subsequent editing of the result, repositioning is expressed as a binary number. (see example, Fig. B)

Track									Zone GP	Zone	Half track	RB Flag
11	10	9	8	7	6	5	4	3	2	1	0	

Fig. A

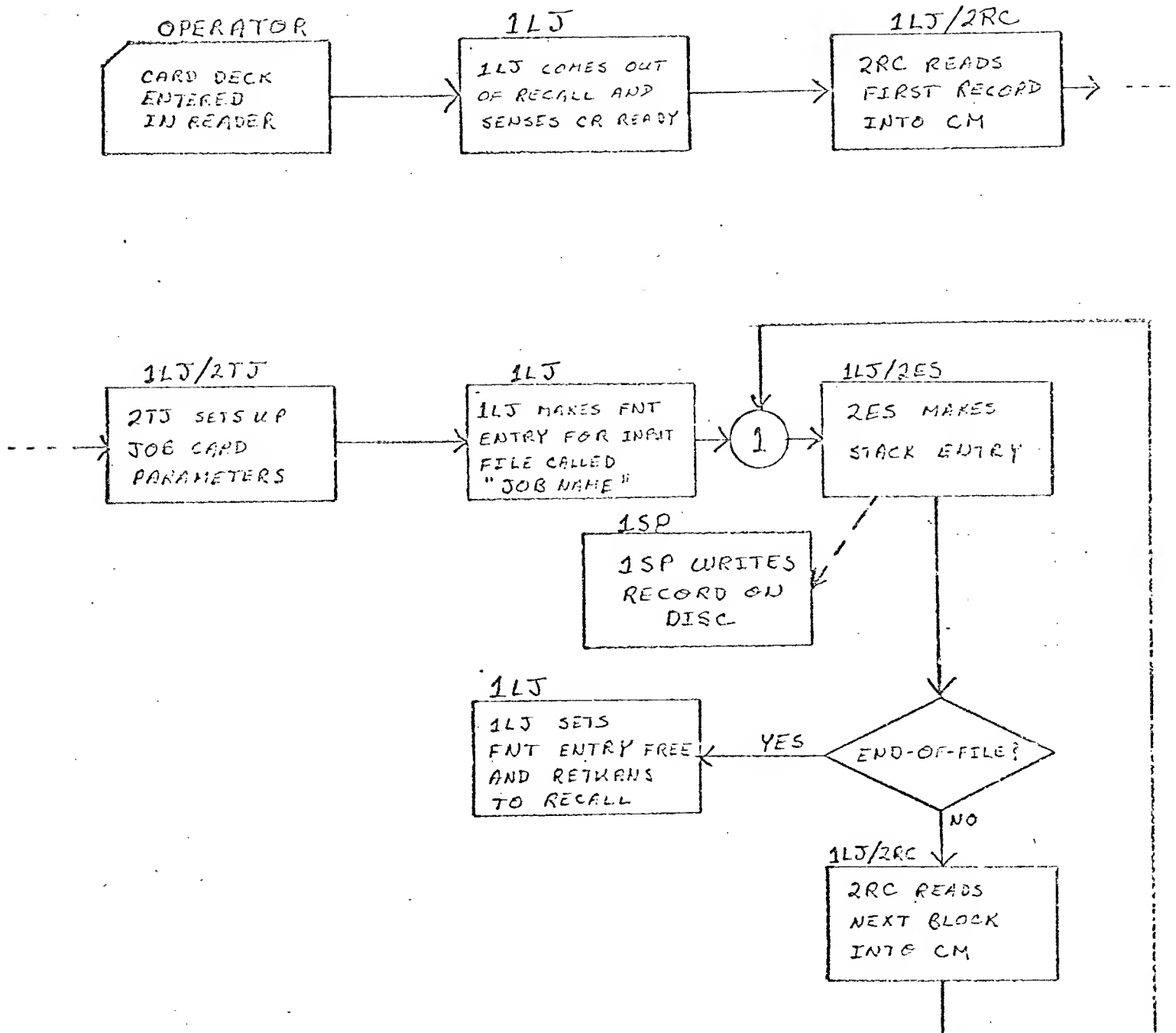
11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	1	0	0	1	0	1	1	0	1
0	0	0	0	0	1	0	1	1	1	0	1
0	0	0	1	1	0	0	1	0	1	1	0
0	0	0	1	1	0	0					

Fig. B

current position
last used file pos.
logical difference
edited current priority

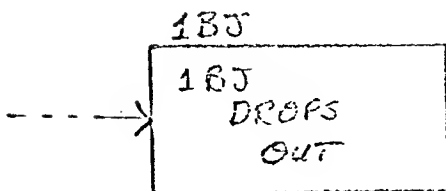
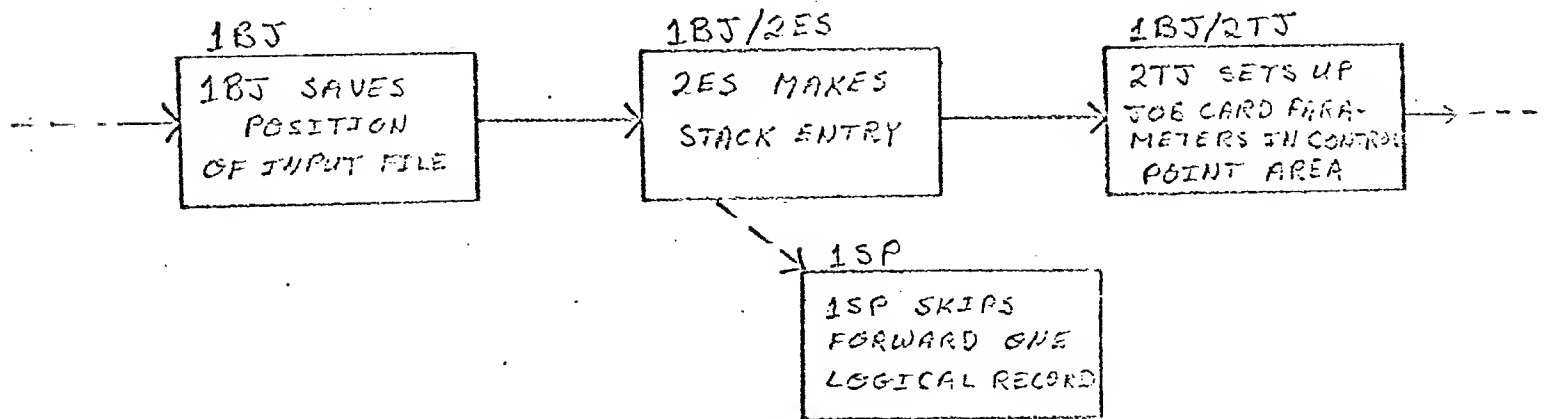
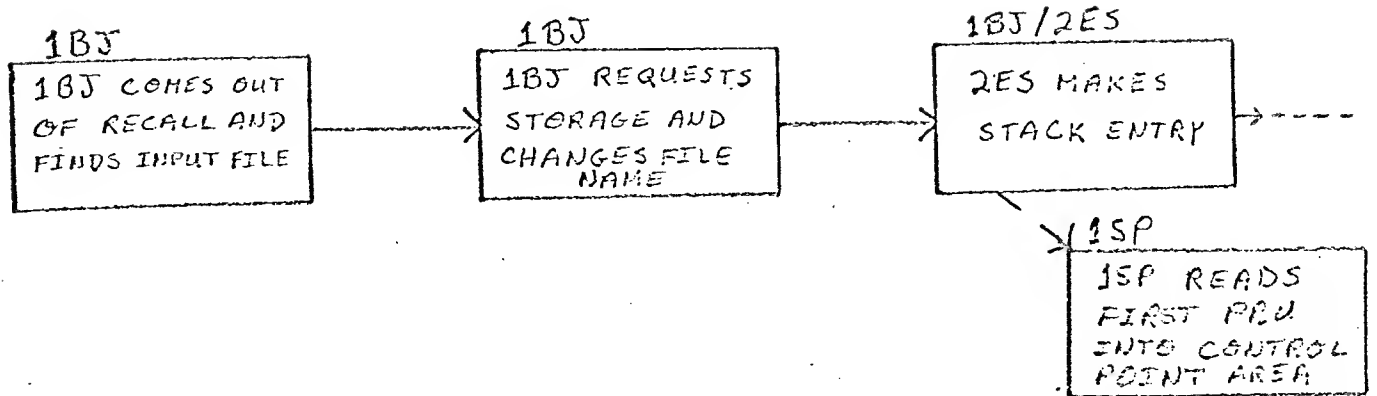
JOB FLOW

LOAD JOB



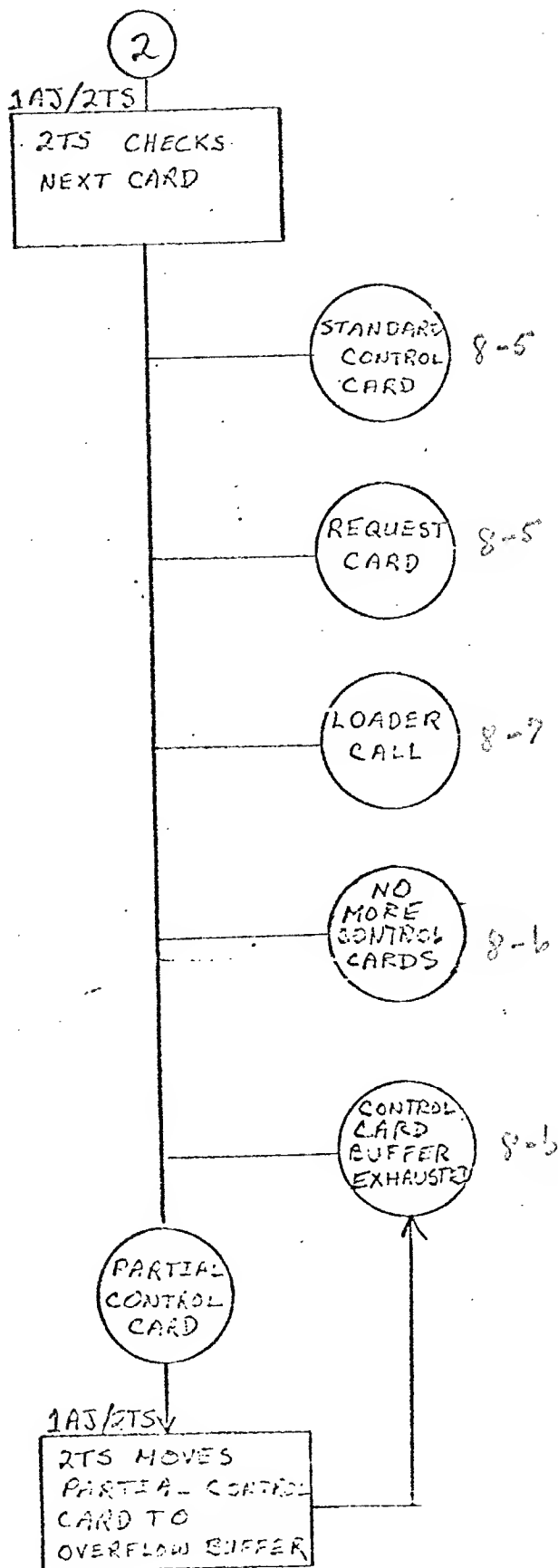
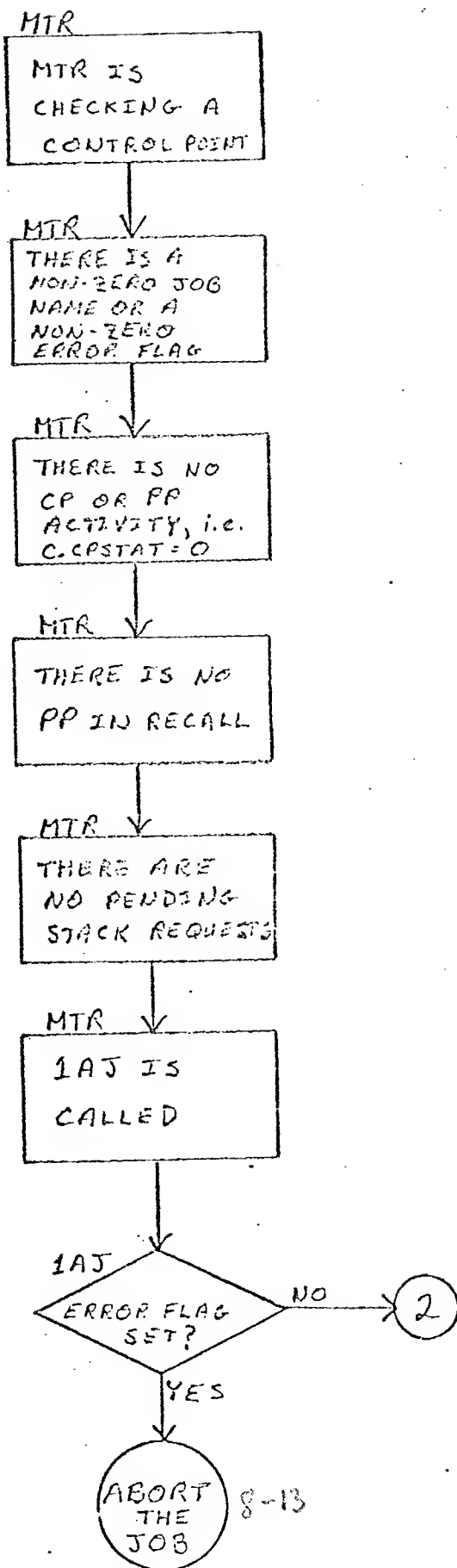
BEGINNING A JOB

8-3



PROCESSING A JOB (ADVANCE JOB)

8-4



The following paper will describe the major SCOPE 3.0 Peripheral Processor routines involved in getting jobs into, through, and out of a 6000 series computer.

Familiarity with the hardware and with the external characteristics of SCOPE 3.0 is assumed.

The discussions will begin with the system as it stands immediately after completion of a dead-start load. The system monitor (MTR) is in PP0, idling as it waits for a request to process. (It should be realised that MTR will be continually searching for requests while other processes are being performed elsewhere. The actual functions done by MTR will be introduced and described as needed). The system display package is in PP9, initially displaying on the left console the message :

ENTER DATE

MTR & DSD are permanently assigned to PPs 0, and 9, respectively.

Central memory contains the CM resident tables, flags and routines, collectively known as CMR. The major portions of CMR reside in low numbered central memory, although there are some dynamically variable disc tables located in high core. Again relevant tables and flags will be introduced and defined as needed.

The last component of an idling SCOPE 3.0 system is the Peripheral Processor resident program. A copy of PP resident is loaded into each of the pool PPs (PPs 1 through 8). PP resident contains a number of small service routines which are commonly needed by most PP overlays; additionally, PP resident contains an idle loop which is executed initially and whenever the pool PP is not assigned to a specific task. Each of the pool PPs is aware of its own number and each PP has assigned to it a unique 8 word area in CMR known as a PP communication area. Each communication area has a location called a PP input register, one called a PP output register, and six called the PP message buffer. The PP resident idle loop consists primarily of continual examinations of the PP input register; a zero value input register is the idling condition. A non-zero value indicates that this PP is being requested to perform a task and therefore must load an overlay.

PP overlays are listed in the system library directory which is a part of CMR; the location of the overlay is also given. An overlay may be located either on disc or in central memory; the overlays which are frequently needed are usually kept in central memory, since loading speed is maximized. There are two basic types of PP overlays: primary and secondary. Primary overlays are those programs which may be requested through a PP input register; they are loaded by PP resident at 773B (there is a non-executable 5 PP word header appended by COMPASS) and entered by a

LJM 1000B.

Secondary overlays are essentially subroutines of one or more primary overlays; they are loaded into higher PP core when the controlling primary overlay makes a request to a PP resident routine (R.ØVL) to do so. The majority of secondary overlays are coded to be loaded at 1773B and entered by a

RJM 2001B.

However, there are several which must be loaded at 2773B, 3773B, etc. Also, there are a few secondary overlays which utilize a set of macros (called RELØC) which have the effect of making the overlay self-relocating; such overlays may be loaded anywhere, providing of course that they do not wrap around the end of memory. This latter condition is not detected by the system; it will eventually be recognised by the operator since the destruction of a PP will (usually) destroy the system ultimately.

All PP overlays have three-character names; there are certain conventions which are followed with respect to the first character when assigning names. Primary overlay names either begin with an alphabetic character or the digit one (1). Alphabetic overlays may be called from a running program, a control card (in some cases), or from within the system; numeric overlays may only be called from within the system, i.e. from a PP. Secondary overlay names usually begin with the digit which is closest to the 1000_8 word block at which the overlay is to be loaded; e.g. 2TR is loaded at 1773B. No conventions have been established for the second and third characters, except that they should have some (be it ever so tenuous) mnemonic value.

Now, to return to the idling system, let us begin to do some processing. To set

the system going the operator must type in today's date in the correct format. Until this is done, DSD remains in an idle loop waiting for keyboard input. When the date is received, its legality is checked and DSD enters the value into a flag location in CMR for use by the remainder of the system. DSD then brings up the A (dayfile) display on the left console and the B (control point) display on the right console.

At this point, normally the operator would type in AUTO to prepare the system for full multiprogramming; for instructional purposes, we shall simplify the procedure

and only type in

1. READ.

DSD has been cycling through its main loop which consists primarily of refreshing the console displays and checking for keyboard input. When the above command is received, DSD branches to the routine which processes a request for the READ package. It is known to DSD that the name of the PP overlay comprising the READ package is 11J (load job) and that 11J must be assigned to a pool PP.

1. DSD posts in its output register a request to MTR (DSD also has a PP communication area. Its input register is not used since PP9 is not a pool PP and cannot be assigned tasks by MTR). This particular request will be a "request peripheral processor" (M.RPP). (It will be seen later that M.RPP is a subset of another monitor function and exists only for historical reasons.) DSD's output register will contain the code for M.RPP and the control point number to which the pool PP will be assigned, in this case, 1. As is usual, additional information is given in the PP message buffer; in the case of an M.RPP function, the first word of the message buffer is expected to contain the input register setting for the requested pool PP.

2. Part of MTR's loop consists of checking the output registers of the 8 pool PPs and DSD. A Zero value constitutes no request; a non-zero value indicates some function to be performed. In this case, MTR will pick up the first word of DSD's message buffer, insert the control point number (1), and place the word into the input register of an available PP. In this particular situation, the PP chosen will be PP1 since nothing

else is happening; however, in other cases, any of the other 7 pool PPs might have been selected.

3. The PP resident in PPI will now notice that its input register is non-zero, specifically that the high order 18 bits contain the display code name of the PP primary overlay which is required to perform the desired task, (1LJ), the next 3 bits are zero, and the next 3 bits contain the control point to which the task is assigned (1). PP resident looks up the name 1LJ in the system library directory, reads the overlay into its own memory and transfers control to 1LJ at 1000_8 .

4. 1LJ first performs the initialization which is common to nearly all primary overlays. The contents of the PP input register are read into PP low core; the control point number is used to compute the address of the control point area. (CMR contains a control point area for each of the seven control points. This area contains information about the state of the job (if any) currently assigned to the control point. Control point area contents will be described as needed during the course of this discussion.) The control point reference address (RA) and field length (FL) are obtained from the control point area. These quantities are retained in PP low core for use by the primary overlay and also any secondary overlays called during the performance of the assigned task.

5. 1LJ enters the job name READ into control point area 1 in order that DSD may display, for the operator's information, the current job. Prior to this time, control point 1 (and all other control points) had a zero job name indicating that the control point was vacant.

6. 1LJ then posts a request (M.RSTOR) in its output register for central memory storage for buffer space. Again, MTR will sense the request and in this case will immediately give control point 1 the necessary storage. Granting this request consists only of setting the desired FL into control point area 1 and resetting the RAs of control points 2 - 7 equal to RA+FL of control point 1.

7. At this point, we shall introduce a table in CMR called the Equipment Status Table.

(EST). The EST contains an entry for each I/O device which is present in the configuration; one word is required for each device. Each entry contains a 2-character display code device type, a flag field to indicate assigned or available, and hardware information (channel number, equipment number etc.) Entries may be ordered whimsically, but the relative position (EST ordinal) with respect to the first word of the EST must be known by the operator in order to assign equipment.

Returning to the job processing, 1LJ now posts a request (M.REQP) in its output register for equipment of type CR (card reader). MTR searches the EST for such a device and will select the first unassigned card reader (lowest numbered ordinal) that it finds. Hence, the order of EST entries must also be known to determine which card reader goes to which READ package. (Note that if a second READ package were assigned to another control point, it would handle the card reader with the next highest EST ordinal, and so on.) Each copy of 1LJ can handle only one card reader. MTR marks the selected card reader EST entry as assigned and returns the EST Ordinal to 1LJ via 1LJ's message buffer.

8. 1LJ then sets up the appropriate hardware instructions to sense the status of its assigned card reader and let's assume first that it is not ready, i.e. the operator has not yet put any job decks in the hopper. The system, in general, tries to keep as many pool PPs available as possible and since it may be a while before the card reader is ready, we introduce a new concept, PP recall. 1LJ uses the old-fashioned form of PP recall for historical reasons only. Old-fashioned PP recall required the use of a location in the control point area, known as the PP recall register. 1LJ writes the contents of its input register into the PP recall register and then issues a monitor request (M. DPP) to return the PP to the pool. MTR examines control point PP recall registers during its main loop and if one contains a non-zero quantity, the register contents are placed in the input register of an available PP and the recall register is reset to zero.

Modern PP recall requires the use of a monitor function (M.RPJ - request peripheral

job) followed by an M.DPP. M.RPJ allows the specification of a time interval which is to elapse before a pool PP is assigned (the time interval may be zero, meaning that assignment should be made as soon as possible). It should be noted that:

- (a) M.RPJ can (and should) replace M.RPP, and
- (b) M.RPJ would allow two or more PPs at a single control point to enter recall simultaneously; use of the PP recall register does not permit this ability.

MTR maintains a list called the PP delay stack in PPO memory of PP programs to be called after a specified interval; the delay stack is examined each time the clock is updated. It can be seen that 1LJ could make effective use of the delay stack feature since the event it is awaiting (card reader ready) is measurable in human time; a pool PP would be required less frequently using M.RPJ than by using the PP recall register. But I digress. 1LJ is now said to be in recall; when it is reloaded, it starts anew, i.e., returns to step 4. This is necessary since it cannot rely on being loaded into a properly initialized PP. However, 1LJ need not repeat all the steps since it can determine that control point 1 already has CM storage and a card reader. Again it tests the status of the card reader, returning to a recall state if it is not ready. This time, we shall assume a ready condition.

9. Other than the test for card reader ready, 1LJ has no I/O driver capability; it must load a secondary overlay. The card reader driver overlay is named 2RC. 2RC is loaded (as are all other secondary overlays) by presetting 2 locations in PP low core with the display code overlay name, entering the load address in the A-register, and executing a RJN to the PP resident routine, R./VL. 1LJ then enters 2RC which reads cards until either the central memory buffer is filled or until an end-of-record or end-of-file card is encountered, whichever occurs first. When 2RC has finished processing, it returns to 1LJ.

10. 1LJ then loads 2TJ (translate job card). 2TJ will locate the first card image in the buffer and check its validity as a job card. (You will recall that the first card of the first logical record of a job is required to be a job card.) Additionally, 2TJ will save in PP low core the job name (appending a job count to keep the name unique),

the required field length and priority for later use. Control is returned to 1LJ.

11. Assuming that the job card was indeed valid, 1LJ now loads 2BP (check buffer parameters). 2BP is responsible for checking that various file tables (which are described elsewhere) contain legal values and for setting up PP low core locations in the form required by hardware drivers to be loaded subsequently. (Note: logically, 2BP should have been loaded before 2RC in step 9 above; the reason for not doing so are not known to the author) 2BP is generally called immediately prior to the loading of any I/O driver. It is more important that it be loaded when user I/O is involved, however, since one assumes that system routines such as 1LJ are not going to produce erroneous file tables.

12. 1LJ next loads 2ES (enter stack). 2ES is a pseudo I/O driver, in a sense. It is loaded and treated as though it actually could read and write an allocatable device (i.e. disc), but it can't. What it does do is to gather together information about the file to be read or written and prepare an entry for a CMR table called the request stack. The necessary information about physical device, file position and so forth is formatted within PP memory and 2ES enters a PP resident subroutine, R.EREQS. R.EREQS in turn searches the request stack for an available location and formats a request to MTR (M.EREQS) to make the entry. (PP resident performs the stack search in order to minimize the monitor loop time; MTR, however, must actually make the entry in order to provide an interlock on the request stack. If two copies of PP resident have simultaneously located the same empty entry, MTR refuses the second M.EREQS function and the rejected PP resident must re-search the stack). When the stack entry has been made, R.EREQS returns control to 2ES which then exits to 1LJ. 1LJ remains in a tight loop, waiting for the central memory buffer to be written to the disc. (The detailed mechanics of disc reading and writing will be covered elsewhere).

13. When the buffer is empty (i.e. the disc operation is complete), 1LJ loads 2BP and then 2RC to refill the buffer from the card reader. Steps 11 through 13 are now repeated.

until an end-of-file card has been read and written. It can be seen that there is now an exact logical record-by-logical record copy of an input file on the disc. The only thing that now remains is to dissociate this file from the READ package and make it available for running. This is done by making an entry in another CMR table called the File Name Table or FNT. (The FNT was used during the course of constructing the input file, but in an uninteresting and difficult-to-describe manner, so we shall introduce it here). The FNT contains a 3-word entry for every file known to the system; an entry contains the file name and type, the type of device on which the file resides, location information peculiar to the particular device, current state (busy or not busy), the priority and the control point to which the file belongs. In the case of the input file which we have just constructed, 1LJ will make an FNT entry containing the following items:

- file name; this will be the job name extracted from the job card; this name was made unique by virtue of the fact that 2TJ appended a job count;

- file type: this will be type input. (Other file types will be introduced as needed).

All files in the FNT of type input constitute the input queue.

- device type: this will be disc, either 6603 or 6638.

- location information: during the course of constructing the input file, various disc tables were built in central memory; the FNT entry contains pointers to these tables.

- current state: this will be not busy

- priority: this will be the priority extracted from the job card as the high order part (called priority level) and zero as the low order part (called sublevel). The use of priority bits is described elsewhere.

- control point assignment: this will be zero, indicating that the file is not assigned to any control point.

- additionally, in the case of a file of type input, the FNT contains the central memory and ECS field lengths requested on the job card.

Having entered the file in the input queue (i.e. made the FNT entry), 1LJ returns to check the status of the card reader (step 8), entering PP recall if there are no more cards in the hopper or constructing another input file if there are.

Let's assume there are 5 more jobs, so 1LJ reiterates the above process 3 more times ultimately reaching the state where it is going in and out of recall. We now have an input queue containing 4 jobs (files) it will be necessary to bring up a package to select one or more jobs for execution. Before we do that, however, to keep life simple and for instructional purposes, let us drop the READ package. This is sometimes done during normal running when there are no more job decks to be entered in order to free up the buffer space that 1LJ has.

POPPING
IE
EAD
PACKAGE

1. The operator must initiate the dropping of any control point; this is done by typing x.DRØP where x is the control point number. In this case, he types:

1.DRØP

Again, DSD will detect keyboard input during its main loop and will branch to process the DRØP command. Processing consists of sending a request (M.ØPDRØP) via DSD's output register to MTR. The M.ØPDRØP also contains a control point number, in this case, 1. The effect of this function is to cause the monitor to set an error flag within the control point area of control point one.

2. Now you will recall that 1LJ is still popping in and out, testing card reader ready. In addition to this already mentioned test, 1LJ also checks the control point 1 error flag. (It also checks the error flag during processing, so that READ may be dropped in the middle of card reading, but this is not too relevant to the current discussion). If at any time 1LJ discovers that the error flag is non-zero, it will drop its PP (via the M.DPP function) without setting the PP recall register. This leaves control point 1 still occupied but inactive, since there is now no way for any more activity to be performed.

3. Meanwhile, MTR has been cycling through its main loop, part of which consists of an Advance Control Point routine. The Advance Control Point routine is entered once every 64 milliseconds and each time it checks the condition of one control point, beginning with 1 cycling up to 7 and returning to 1. A control point is said to require advancing if all of the following conditions are true:

- There is a job name. (This condition is met since the job name READ has not been cleared.
- There is no central processor activity. (The READ package has never initiated any CP activity)
- There is no peripheral processor activity. (1LJ has dropped itself without recall and it was the only PP which was attached to control point 1).
- There are no PP recall requests (either in the delay stack or in the recall register). (1LJ did not set the recall register before it dropped).
- There are no pending entries in the request stack. (The only request stack entries made by 1LJ are during the course of preparing a job for entry into the input queue and there are no more jobs to be prepared).

It can be seen, therefore, that control point 1 must be advanced.

Advancing a control point, from MIR's point of view, consists of calling 1AJ (advance job) into a pool PP and attaching the PP to control point 1.

4. As will be seen later, 1AJ can follow several paths varying in complexity, depending on the state of the control point which is to be advanced. The case currently under investigation, namely that of dropping a system control point (READ) is probably the simplest and involves the following steps.

1AJ checks the error flag which is set. 1AJ then checks the priority of the control point which is given in the control point area; now, since READ is a system control, it has zero priority. This fact greatly simplifies the process as will be seen later. 1AJ requests the loading of the overlay 2EF (error flag). 2EF checks the actual value of the error flag to determine which message to enter into the dayfile. In this case, the value is 6 (F. ERØD) which directs 2EF to write the following message:

hh.mm.ss	READ	ØPERATØR DRØP
(Time of Day)	(job name)	(message)

2EF returns to 1AJ.

5. 1AJ begins a search of the FNT, looking for (in this case) files which are assigned to control point 1. Each time such a file is located, 1AJ will load the overlay 2DF (drop files). 2DF "removes" the file from the FNT; what constitutes removal of a file is

dependent upon the file type.

It should be specially noted that those files which were entered into the input queue were assigned to control point 0 and are therefore not dependent on the existence of their creating READ package). When all files have been removed from the FNT, 1AJ then searches the EST for equipment assigned to control point 1; in this case it will find exactly one equipment, the card reader. Again, 2DF is called to return the card reader to the pool of available equipment.

1AJ then posts an M.RSTOR function to MTR to request that all central memory storage associated with control point 1 be dropped. Control point area 1 is then cleared, i.e. the job name, etc, are set to zero.

We are now faced with a system in an essentially idle condition except that there are four jobs in the input queue. We must now initiate processing. To do so, we will bring up one copy of the next package at, say, control point 3.

1. The procedure for bringing up NEXT is nearly identical with that for bringing up READ and so we shall abbreviate the description. The operator initiates the process by typing 3.NEXT. It is known to DSD that the name of the PP overlay comprising the NEXT package is 1BJ (begin job). DSD requests, again with the M.RPP function, that 1BJ be assigned to a pool PP at control point 3. MTR will select a pool PP and PP resident will load the overlay 1BJ.

2. 1BJ will perform the usual initialization of PP low core (See READ PACKAGE - item 4) and then enters the job name NEXT into the control point area for display purposes.

3. MTR maintains in CMR a location (T.UAS) which contains the current amount of available central memory and ECS. 1BJ reads these quantities into PP low core and then requests the FNT channel. (A slight digression: MTR maintains a channel status table which includes entries for the 12 hardware channels, and several pseudo-channels.

Before a PP can use one of the hardware channels or alter a table represented by a pseudo-channel, it must request the use of the channel from MTR. This is another method of providing a table interlock.) In the case under consideration, the FNT

channel really need not be obtained since there is no possible conflict; but a more normal situation would include several copies of LBJ making identical searches and it is considered desirable to execute each job only once.

4. LBJ performs a search of the File Name Table, looking first for a file of type input. Having located such a file, it compares the requested CM and ECS field lengths with the currently available storage. If there is adequate storage for the job, then the job priority is compared with the priorities of all other jobs which meet the storage requirements. The job with the highest priority will be selected for execution. Also, during the course of every nth FNT search by LBJ (where n is an installation parameter), the priority sublevel of all files in the input queue will be incremented by one; the sublevel will not be allowed to overflow into the level, however. Hence, it can be seen (assuming that enough storage is always available) that priority level as declared on the job card is the major determinant of the order of execution and that, within a set of jobs of the same priority level, the length of time the job has been in the job queue will determine the order of execution.
- If LBJ had failed to find a candidate for execution, it would now enter PP recall in the old-fashioned manner, i.e. by setting the PP recall register and dropping out. However 4 jobs are known to be in the input queue and so we carry on.

5. The selected FNT entry must now be altered. The file name is changed from the unique job name to INPUT; the type is changed from input to local; and the control point assignment is set to 3. Also, the fields which contained the ECS and CM field length requirements are altered to the more usual disc position information. The position of the input file, at this point, is logically rewound, i.e. the next record to be read is the first (control card) record.

All this means that the user's input file will be available to him alone (since it's assigned to what will soon be his control point) and that it may be referenced simply by calling for the file named INPUT.

6. LBJ then sets up control point area 3 for the selected job. The job name is

changed from NEXT to the name of the job to be executed. The priority (with the sublevel field reset to zero) is entered into the control point area with the aid of a monitor function (M.RPRI). The current position of the input file is also saved in the control point area for reasons which may become clear later. 1BJ then calls 2ES to format a disc request to skip forward one logical record on the input file.

(You will recall that this first record contains only control cards and is presumably not of interest to the user's program). 2ES, through PP resident and the monitor, enters the request stack, leaves the INPUT FNT entry busy and then returns to 1BJ. 1BJ loops until the FNT entry becomes not busy (request completed). The INPUT FNT entry is now positioned at the user's first data record.

7. 1BJ now indulges in a game of fake-out-the-system, the object of which is to read the first sector of control cards into a small buffer in the control point area. The method used is, briefly, as follows. Temporarily, the FNT entry for INPUT is set to reflect a rewind condition (this can be done since the initial file position was saved in the control point area). 1BJ then sets up its own PP to look as if it belonged to control point 0 rather than control point 3. (Control point 0 is a convenient fiction invented for just such cases as this. The area occupied by CMR is said to belong to control point 0). This is necessary because it is illegal for a PP attached to one control point to request that data be read into the area belonging to another control point. 1BJ then loads 2ES to format a request, in the usual manner, to read data into the sector-sized buffer in the control point area. 1BJ then awaits completion of the operation and restores the proper values in its PP and in the FNT. However, instead of returning the rewind position of the INPUT file to the control point area, it sets the position of the next sector of control cards; this allows 1BJ to perform the above sneaky operation each time another sector of control cards is required.

8. 1BJ next calls 2TJ to reprocess and skip over the job card which is the first card image in the control card buffer. 2TJ will again check the validity of the job card. (Because 1BJ has no facilities for writing an error message for the user, a job with an erroneous job card will be passed along into the input queue so that 1BJ/2TJ can sense the error and inform the user). On this pass, 2TJ also saves the job time limit in PP low core. 2TJ returns control to 1BJ which then completes the setting up of the control point area, i.e. the job time limit is entered (via the monitor function M.NTIME); the exit mode is set (via the monitor function M.REM); and the required CM and ECS field lengths are requested (M.RSTOR). If the necessary field length cannot be granted, 1BJ goes into recall. There are two noteworthy items which follow from this.

(a) All during the course of the NEXT package discussion, 1BJ was making tests to determine if the next task had already been accomplished; such tests were not mentioned for the sake of clarity.

(b) At first glance, it might seem that 1BJ would always obtain storage since it did first check to see that there was enough storage for the job. However, there is no interlock on the value contained in T.UAS and 2 copies of 1BJ might have decided there was adequate storage for their respective jobs; if the sum of the storage requirements exceed available core, then the second 1BJ to request storage will have to wait. Also, a job at another control point might have requested, and been granted, an expanded field length after 1BJ has read T.UAS and before 1BJ's request for storage.

However, in this case, storage will be granted and 1BJ will drop out (M.DPP).

9. Referring back to DROPPING THE READ PACKAGE, item 3, it can be seen that the identical condition now exists at control point 3, that is, MTR must advance control point 3. (There is a job name; there is no CP or PP activity; there are no PPs in recall or in the delay stack; and there are no pending requests). IAJ is now loaded into a pool PP and attached to control point 3.

IAJ first checks the error flag which is now not set, so the overlay 2TS (translate statement) is called. As far as 2TS is concerned, there are several major groups of control card types, each of which will be briefly described below. 2TS reads the next control statement from the control card buffer and proceeds in one of the following ways:

(a) SWITCH, MODE, COMMENT, COMMON, RELEASE.

All of these cards require only simple processing and are handled completely within 2TS. When the requested action has been completed and the control card buffer pointers updated, 2TS drops the PP (M.DPP) without returning to IAJ. This leaves control point 3 ready to be advanced again.

(b) REQUEST

Only a very little initial parameter scanning is done by 2TS. A few parameters are stored in PP low core and 2TS calls REQ (request card-processor) by altering the PP input register and returning directly to the PP resident idle loop. (REQ was made into a separate program to avoid duplications since REQ is also callable from a running program) REQ finishes the parameter processing and posts a request for operator action on the console display. It then tests for operator response which naturally will not happen for a period of time measurable in human terms. REQ then simply drops the PP without updating the

control card buffer pointer. (The console message remains, however). This has the effect of leaving the control point ready for advancing but insures that the same REQUEST card will be completely reprocessed. This procedure is reiterated until an operator response is received. REQ then completes processing, updates the control card buffer pointers and then drops the PP; this leaves control point 3 ready to be advanced again.

(c) LOAD, EXECUTE, NOGO and program call card

Minimal parameter scanning is performed and information is left in PP low core and in locations RA + 2 to RA + 100 within central memory. 2TS then calls LOAD into the PP by using the PP resident function R.OVL. (It might be worth noting that this overlay load unconventionally causes a program - LOAD - to be loaded at a lower numbered location than its caller, i.e. LOAD is loaded over LAJ. The tacit assumption is made that LOAD is not sufficiently long to overwrite the spot in 2TS from which the call is made; this is a reasonable assumption in this case, but it should not be made generally).

It is not the purpose of this discussion to detail the operation of the relocatable loader; adequate documentation appears elsewhere. A few general remarks will be made, however. The loading process is performed by one or more of the following routines: LOAD, LDR, 2LA, 2LB, 2LE, and in the central processor, LOADER, & OVERLOAD. The number of routines involved depends upon the type and complexity of the load operation. At the end of loader processing whichever of the 5 PP loader routines is in control drops the PP. If neither a PP or CP execution was initiated during the course of processing (e.g. NOGO card), then the control point is again ready for advancing.

Otherwise, the control point is considered to be profitably occupied and cannot be advanced until PP and CP activity has ceased.

10. Assuming that the job we have selected for execution at control point 3 contains more than 1 sector of control cards, 1AJ/2TS will eventually be faced with the problem of obtaining the second (or n^{th}) sector. And it is a problem; the method used defies detailed description. In brief, 2TS will detect that it has either a partial control card or none at all to process and so it will return control to 1AJ.

(This and the job termination case are the only instances when 1AJ regains control from 2TS.) 1AJ then calls 1BJ into the same PP by the underhanded method of adding one to the second character of the PP input register and returning directly to the PP resident idle loop. (Note well, anyone who is thinking of whimsically changing the name of either 1AJ or 1BJ!). Now referring back to the NEXT PACKAGE, item 7 and accepting on faith that 1BJ knows how to pick its way through the maze, it can be seen that the next sector of control cards will be read into the control point area. After this operation is completed, 1BJ then drops out, leaving it to MTR to advance the control point.

11. Ultimately, 2TS will detect either an EXIT CARD or the end of the control statements and will return to 1AJ for job termination.

MINATION

The following section will occasionally refer back to DROPPING THE READ PACKAGE because some portions of the two procedures are identical. In the current case, however, we are dealing with dropping a user job.

1. Control has just been returned to 1AJ from 2TS. The first decision that is made is that this is the end of the job and not just a request for another control card buffer load. Then the control point priority is tested and it is determined that this is a user job (i.e. it has a

nonzero priority). Next, the control point error flag is checked. Assuming that the job at control point 3 has terminated abnormally for some reason, steps 2, 3 and 4 will be performed; otherwise, 1AJ resumes processing at step 5.

2. 1AJ calls 2CA (checkpoint abort). 2CA checks to see whether or not the job took checkpoint dumps during the course of its execution. If so, all files associated with the job are locked, i.e. left in a state suitable for restarting the job at some later time; if not 2CA simply returns to 1AJ.
3. 2EF is called. In addition to the writing of the error message (see DROPPING THE READ PACKAGE, item 4) 2EF will locate the file named OUTPUT and will insure that there is an end-of-record mark. Also, 2EF will supply a partial core dump of the control point field length. The latter is accomplished through use of the M.RPJ function (without a delay specified) to load DMP (central memory dump program) into another PP. 2EF waits for the dump to be completed by waiting until the input register of the other PP no longer contains the request for DMP. All this additional processing is done because the control point priority is nonzero.
4. 1AJ then uses 2TS to cycle through the remaining control cards to the next EXIT card or to the end of the statements. If an EXIT card is encountered, processing reverts to normal, i.e. the following cards are processed as though no error had occurred. In the latter case, 1AJ ultimately regains control of the JOB TERMINATION point; this time, however, the error flag will not be set and steps 2, 3 and 4 are skipped.

5. At this point, 1AJ records the running time used at control point 3; all further PP time incurred in terminating the job is considered to be system overhead.
6. 1AJ calls 2CJ (complete job day file) which appends to the file OUTPUT all dayfile messages relevant to the job at control point 3. During the course of execution, these messages were being collected in two places: chronologically within the system dayfile for display purposes and as a permanent record of system activity; and in the job dayfile. Messages are collected in this manner to eliminate the problem of extracting messages belonging to one job when the job terminates.
7. 2CF (close files) is called to search the FNT for files associated with this control point. If a file of name OUTPUT, PUNCH, or PUNCHB is located, the file is assigned a disposition code of print, punch Hollerith or punch binary, respectively. Then, all local files associated with control point 3 which have a nonzero disposition code (including those just set to nonzero) are entered into the output queue. A file is considered to be in the output queue when the following changes have been made to the FNT entry:
 - (a) the file name is changed to the job name (as recorded in the control point area)
 - (b) the control point assignment is cleared, i.e. set to zero
 - (c) the priority level (as recorded in the control point area) with a sublevel of 1 has been entered.

With one exception, files in the output queue are of type local (local to control point zero). The exception is the file which was named OUTPUT during execution; it is set to type output. The reasons for this are obscure.

8. LAJ now goes through the process described in item 5 of DROPPING THE READ PACKAGE, i.e. in conjunction with 2DF, it releases files and equipment associated with control point 3. Again it should be noted that the files entered in the output queue were dissociated from the control point and hence, are not affected. Storage is released and the control point area is cleared.
9. In the case of dropping a system (zero priority) job, LAJ then dropped out leaving the control point vacant. Now, since it would be a great nuisance to require the operator to bring up NEXT for each job to be run, LAJ will automatically bring NEXT back to control point 3. This is done on the theory that once a NEXT package resides at a control point, it ought to remain until the input queue is empty. LBJ is summoned into the same PP by means of the "adding one" method described in item 10 of NEXT PACKAGE; LBJ begins processing at item 2.

PUT
KAGE

We will now leave control point 3 to the process of selecting and executing jobs and consider the problem of what to do with the output queue.

1. The procedure for bringing up the OUTPUT package is nearly identical with that for bringing up READ and NEXT, so again we shall abbreviate. The operator initiates the process by typing, say 2oOUTPUT. The name of the PP overlay comprising the OUTPUT package is 10T. 10T will be assigned to a pool PP at control point 2.

2. *10T performs the usual initialization of PP low core (see READ PACKAGE 4), and then enters the job name OUTPUT into the control*
10T performs the usual initialization of PP low core (see READ PACKAGE 4), and then enters the job name OUTPUT into the control

point area for display purposes.

5. lØT is capable of simultaneously supervising the printing or punching of as many as six files from the output queue, but since each requires the permanent allocation of a 1000_8 word CM buffer (in addition to the base 100_8), lØT counts the number of printers and punches to minimize buffer space. The EST is searched (directly by lØT) to count the number of equipments of type LP or CP. (Note that these equipments are not reserved at this point). An M.RSTOR function is then issued to MTR for $N*1000B + 100B$ words (maximum $6100B$), where N is the number of output devices. This storage will remain assigned for as long as OUTPUT occupies the control point. The base 100_8 words of storage are used to control the ownership of each of the N buffers.

4. lØT then checks for a free buffer; since this is the first time through, there will naturally be at least one. The FNT channel is reserved via a monitor function and the FNT is searched for files in the output queue. During the single pass through the FNT, several things occur.

- (a) The print file, if any, with the highest priority is selected for printing. A buffer is assigned to the file and an equipment of type LP is requested from MTR. If either a buffer or a printer cannot be obtained (as can happen at later stages) the file is simply returned to the output queue.

- (b) The punch file, if any, with the highest priority is selected for punching. A buffer and a card punch are assigned as in (a) above. The PP occupied by 1ØT is capable of processing only print files, so another PP must be brought into play for a punch file. 1ØT requests (M.RPP) that MTR assign a pool PP to control point 2 for the overlay called 1PØ (punch output).
 - (c) At an interval selected by the installation, the priority sublevels of all files remaining in the output queue are incremented by one. This insures that for jobs with the same priority level, the output of the eldest in the output queue will be processed first.
 - (d) Any files selected for processing during this pass must be removed from the output queue; this is accomplished by assigning the PNT entry to, in this case, control point 2.
5. If there were no files selected from the output queue, 1ØT would set the PP recall register and drop out. However, let us assume a print file has been found. 1ØT then calls 3ØT.
6. 3ØT has a few peculiarities which are worth pointing out. 3ØT is called only by 1ØT and in fact, operates as though it were part of 1ØT. It is merely an accident of history that it is not physically a part of 1ØT. Before 3ØT is entered, 1ØT sets some addresses of 1ØT subroutines into PP low core in order that 3ØT may use them. The most noteworthy of these subroutines is the one described in item 4. This subroutine is entered frequently by 3ØT during its processing to insure that printers

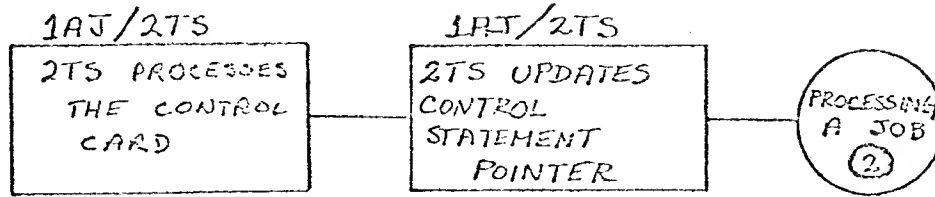
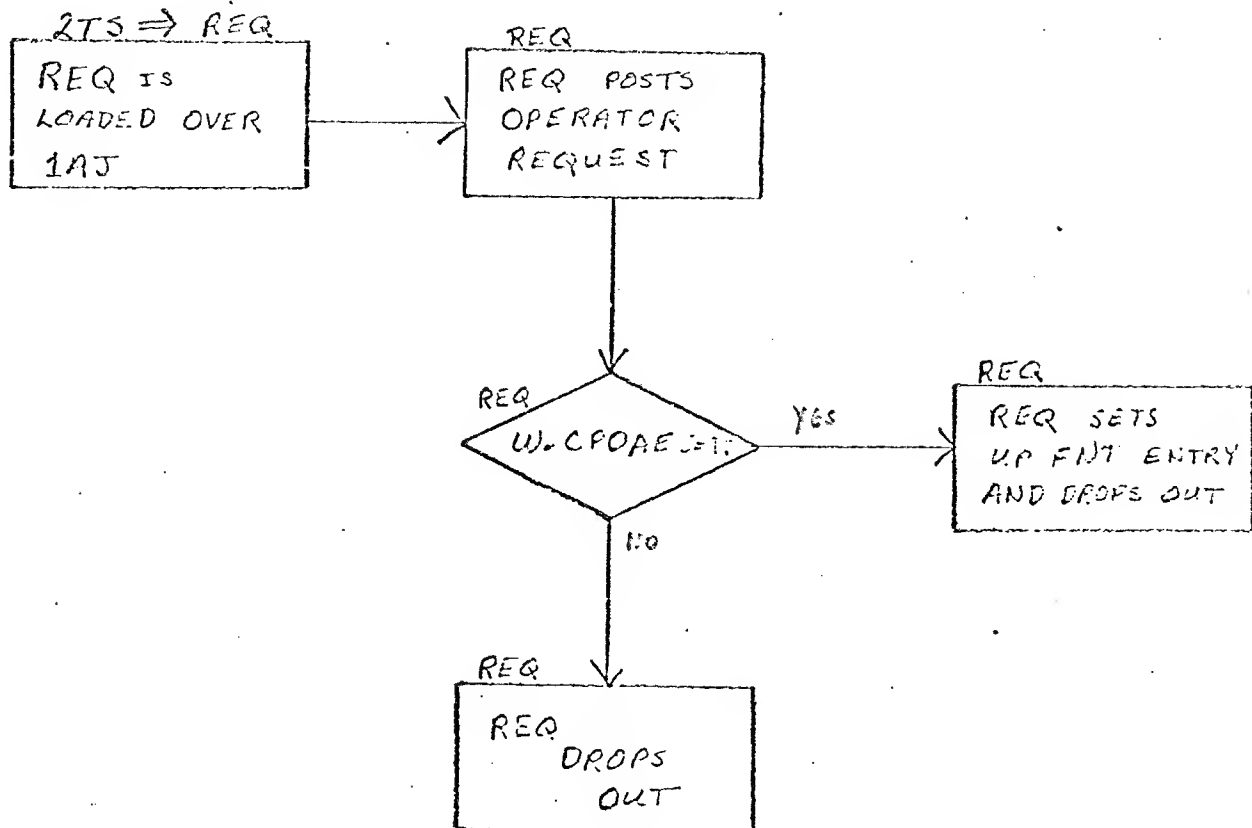
and punches are kept as busy as possible. The punch files are farmed out, as mentioned before, and 3ØT keeps track of the progress in printing the several files for which it may be responsible at any one time.

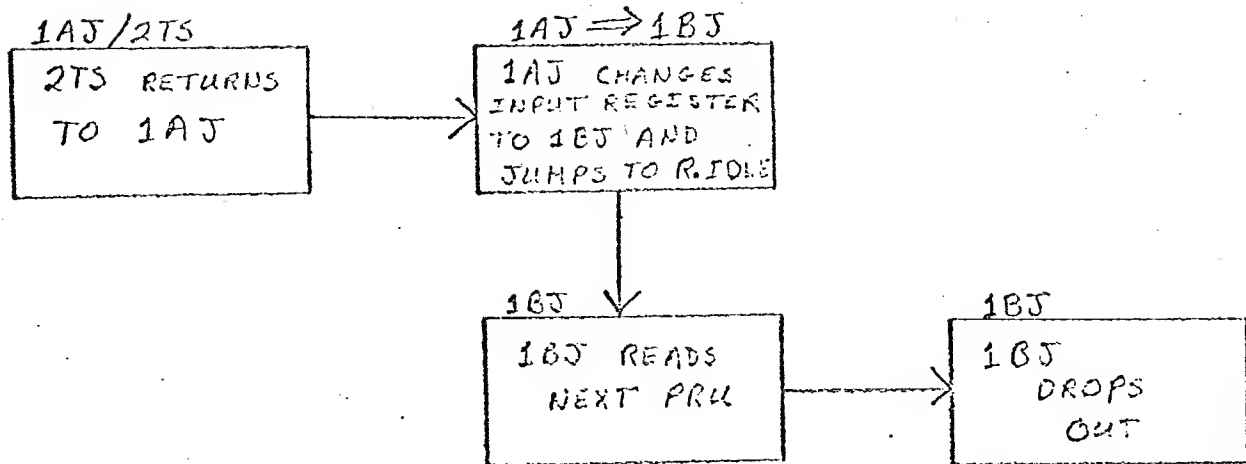
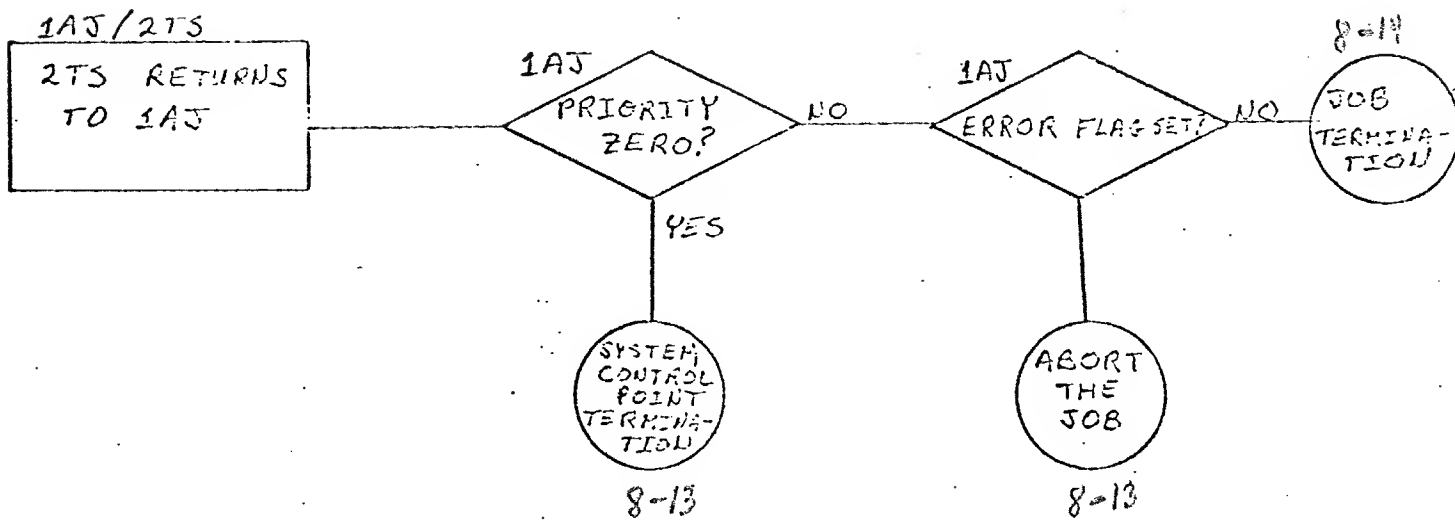
7. We shall now largely ignore the fact that the printing of other files may be initiated and concentrate on the processing of a single file. However, because several files may be printed at once, the technique used by 11J (i.e. that of alternating 2BP/2ES and 2BP/2RC) is not possible for 1ØT/3ØT. The reading of the disc file must be done elsewhere.

The PP overlay CIO was designed primarily to deal with user I/Q; however, it may also be called from another PP program provided that user-type parameters are supplied. 1ØT requests from monitor (M.RPP) the assignment of a PP for CIO, supplying for the input register an address within the base 100_8 word area as well as the name, CIO. (As mentioned previously, space is reserved within this area for information about each print file; the first five words for each file are known as a File Environment Table.(FET). The FET format used by 1ØT is identical to that specified for user I/Q). CIO independently begins to fill the proper buffer with data from the disc output file and 3ØT loops waiting for completion of the read operation. (During this loop it may also be printing data from other buffers or checking for new files to print or punch). When the buffer is full (or when an-end-of record or end-of-file is read), 3ØT begins to print the data on the printer selected for this buffer, interspersing this process with operations on other buffers and

and printers. The printing of the entire file proceeds in this manner. When the end-of-file is reached, 3PT calls 1C0 (complete output) into another PP.

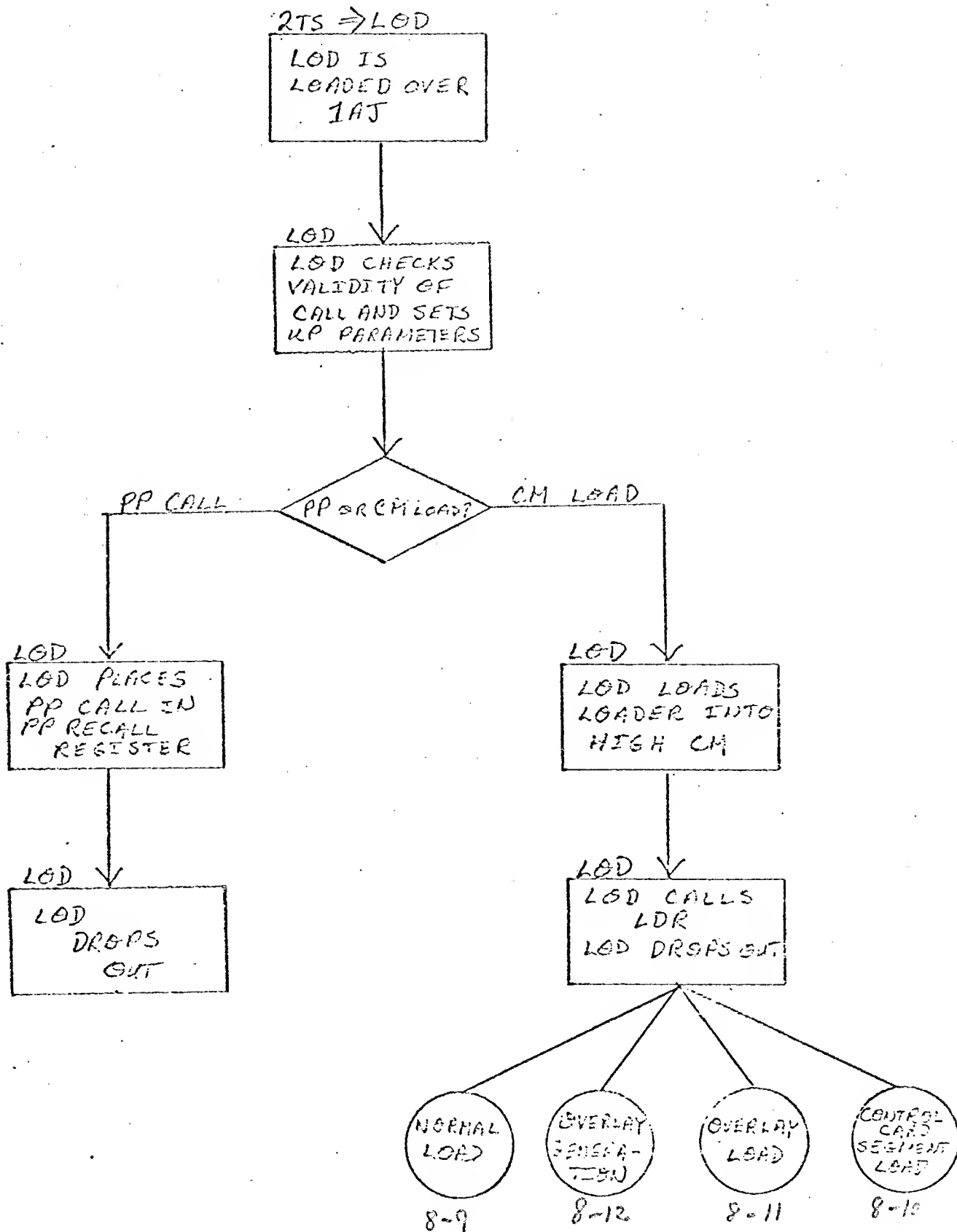
8. 1C0 is given the buffer number and the EST ordinal of the printer used. The printer is released (via a monitor function) back to the pool of available equipment. The buffer number is used to locate the relevant PET, and hence the file name; the buffer is then made available for another print or punch file. 1C0 then loads 2DF to remove the now defunct disc output file. 1C0 then drops out.
(SEE # 6 ON PAGE 19 FOR DAYFILE HANDLING INFORMATION)
9. Meanwhile, (assuming that a punch file was found) 1P0 has been punching cards by alternately calling 2BP/2ES to read the disc file into its CM buffer and 2BP/2PC (punch cards) to punch data from the buffer. CIO is not required since 1P0 only handles one punch file at a time; another PP and another copy of 1P0 would be required if two punches were to operate simultaneously. When the end-of-file is reached, 1P0 itself performs the actions described in (8) above, calling 2DF to remove the disc punch file. 1P0 then drops out.

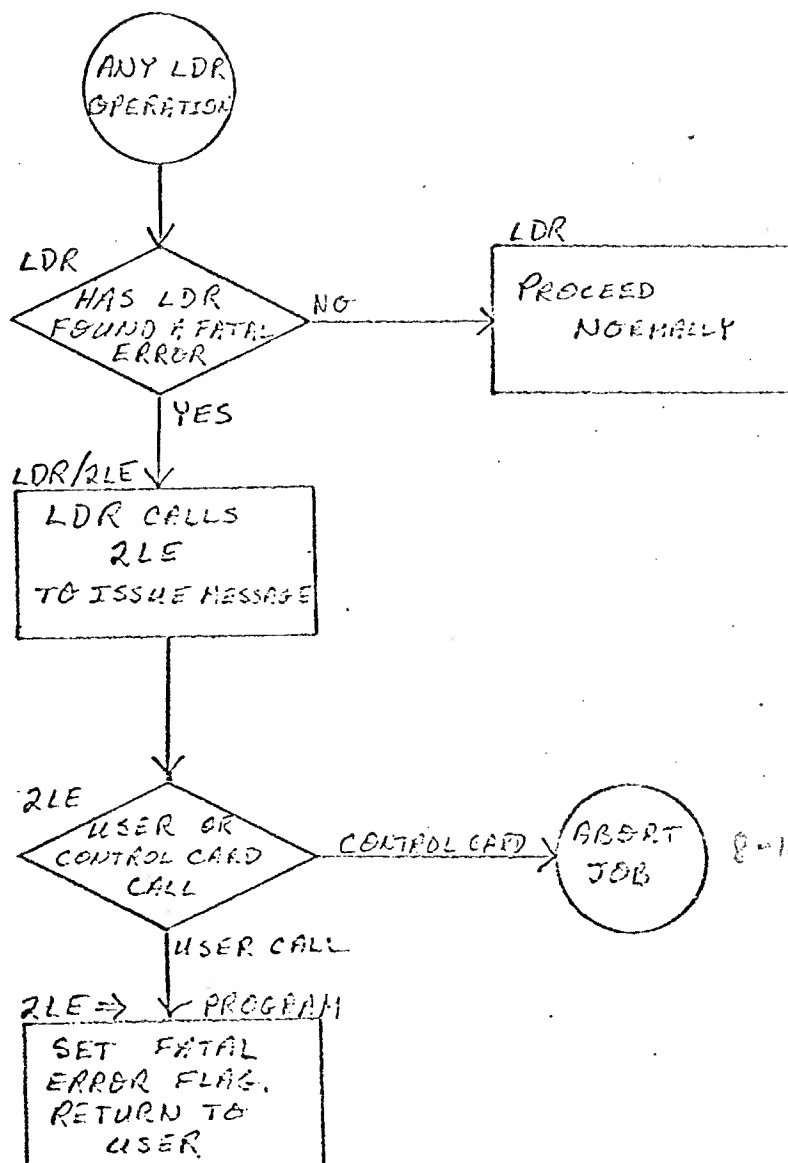
STANDARD CONTROL CARDREQUEST CARD

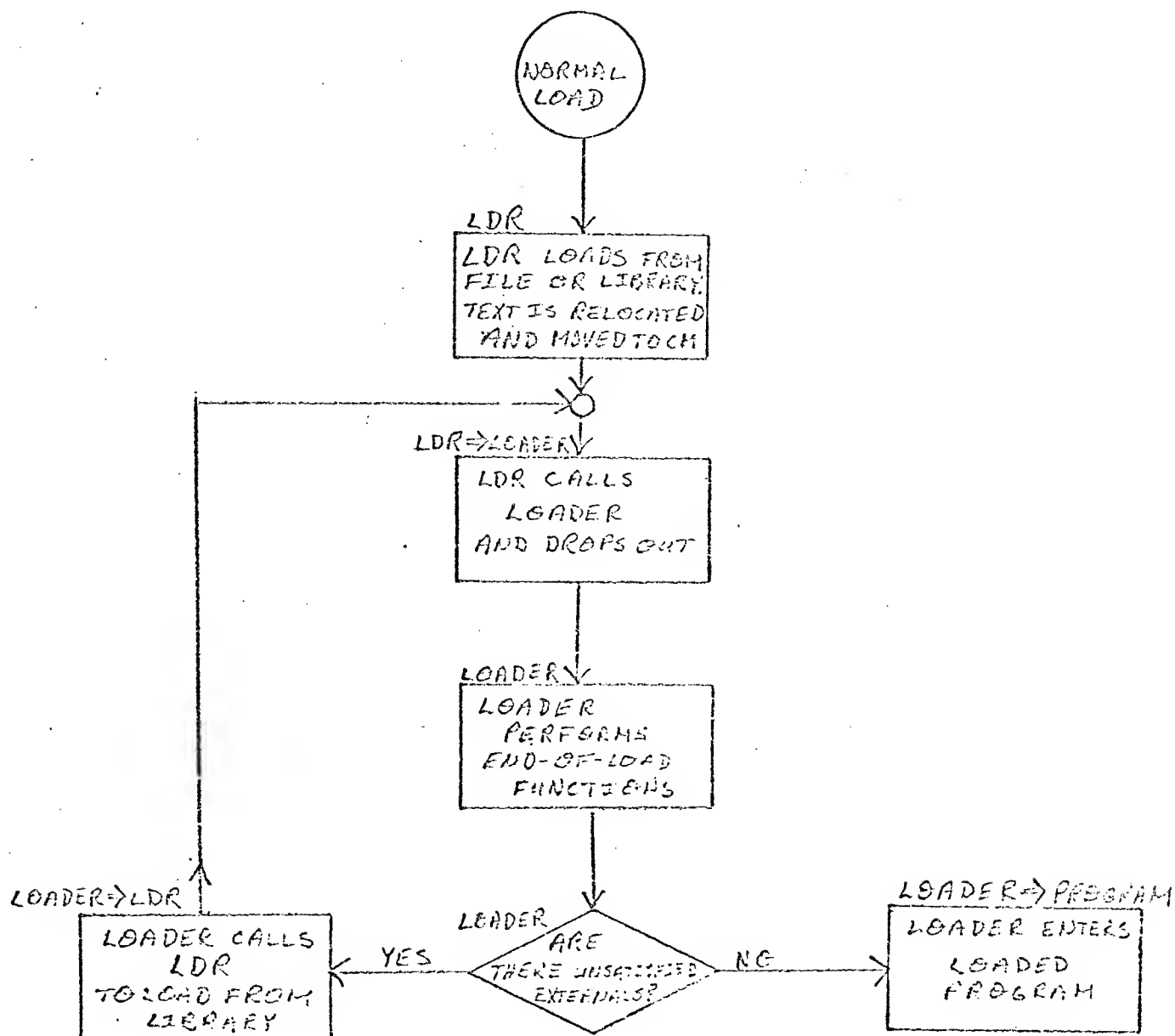
CONTROL CARD BUFFER EXHAUSTEDNO MORE CONTROL CARDS

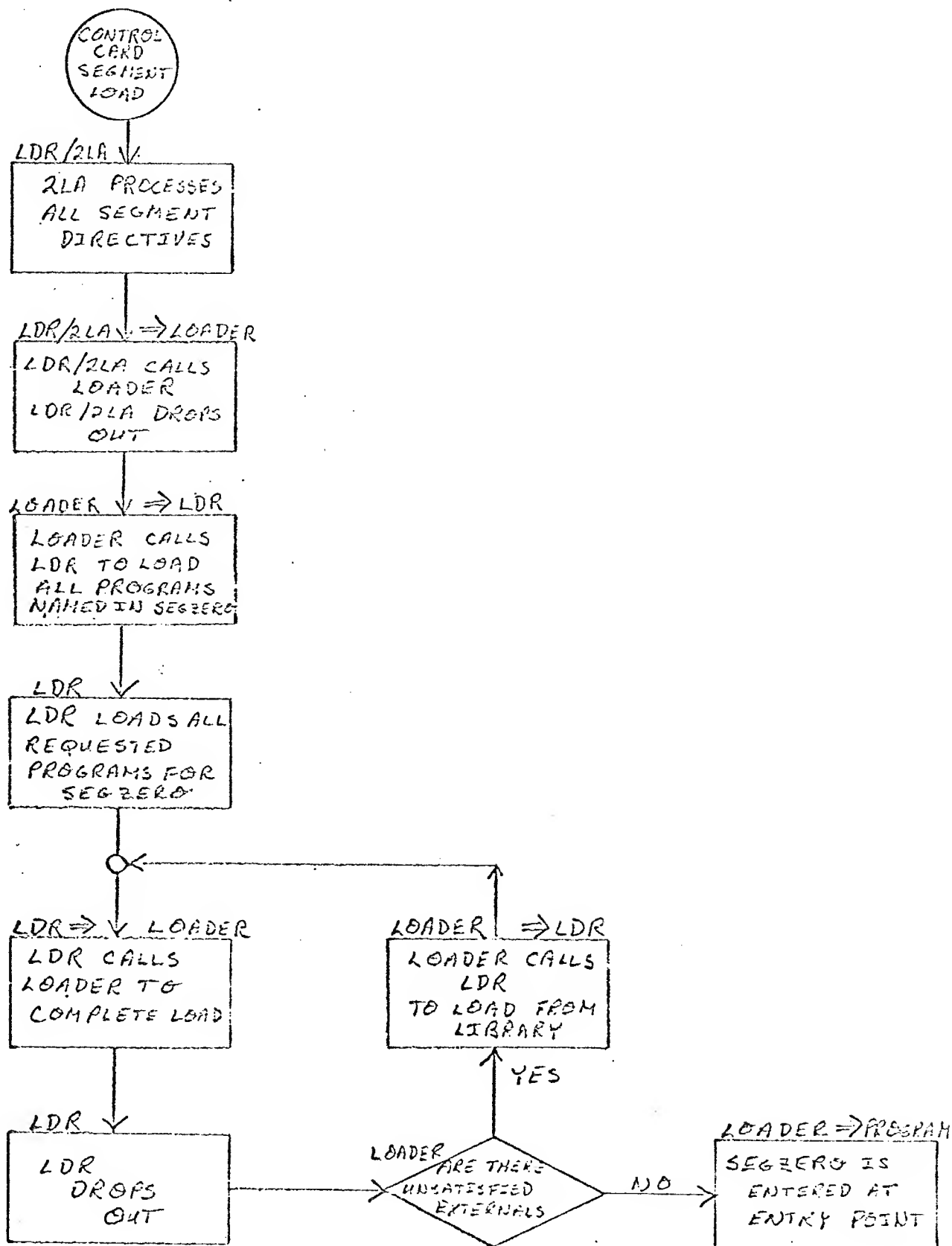
LOADER CALL

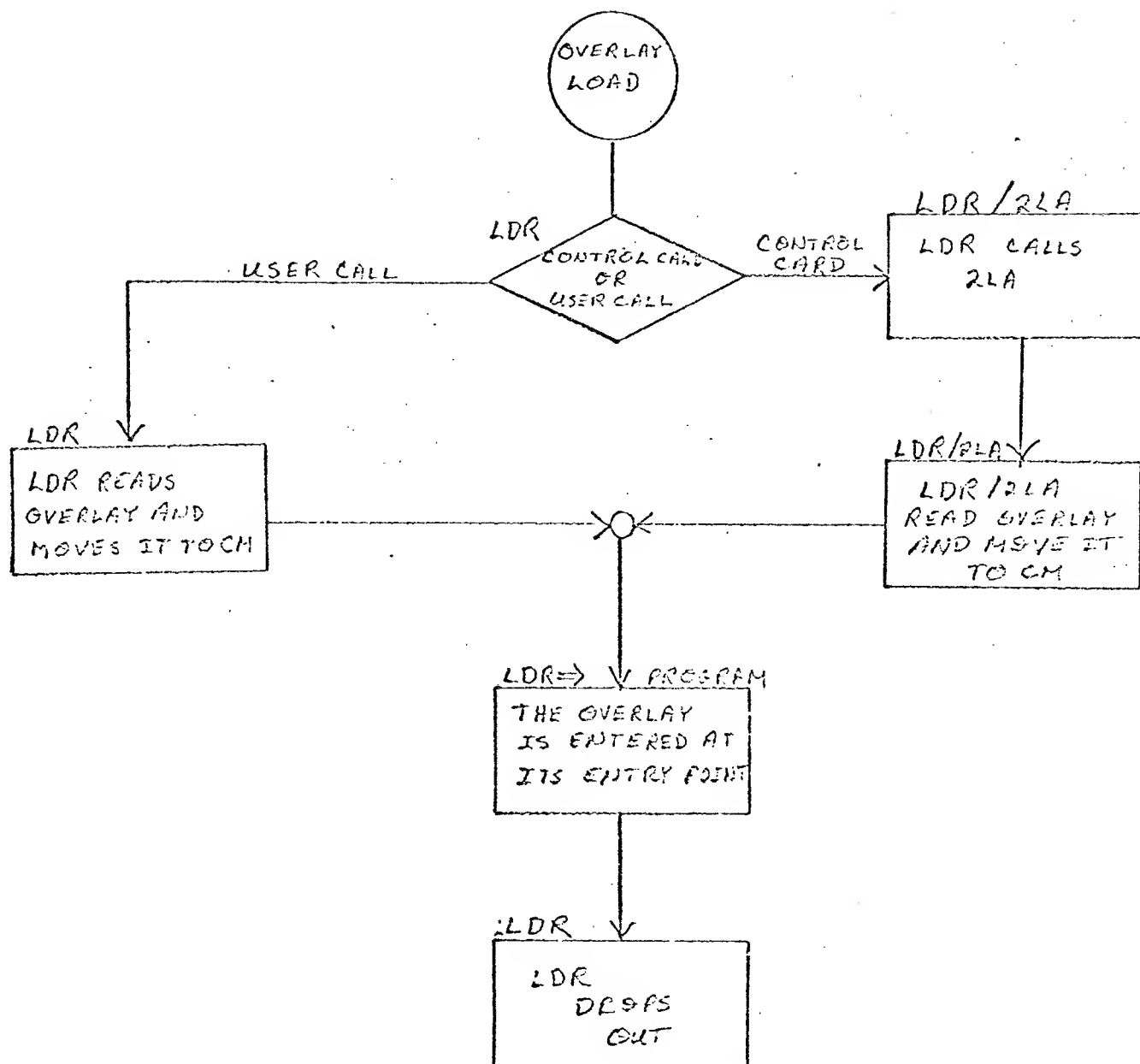
8-7



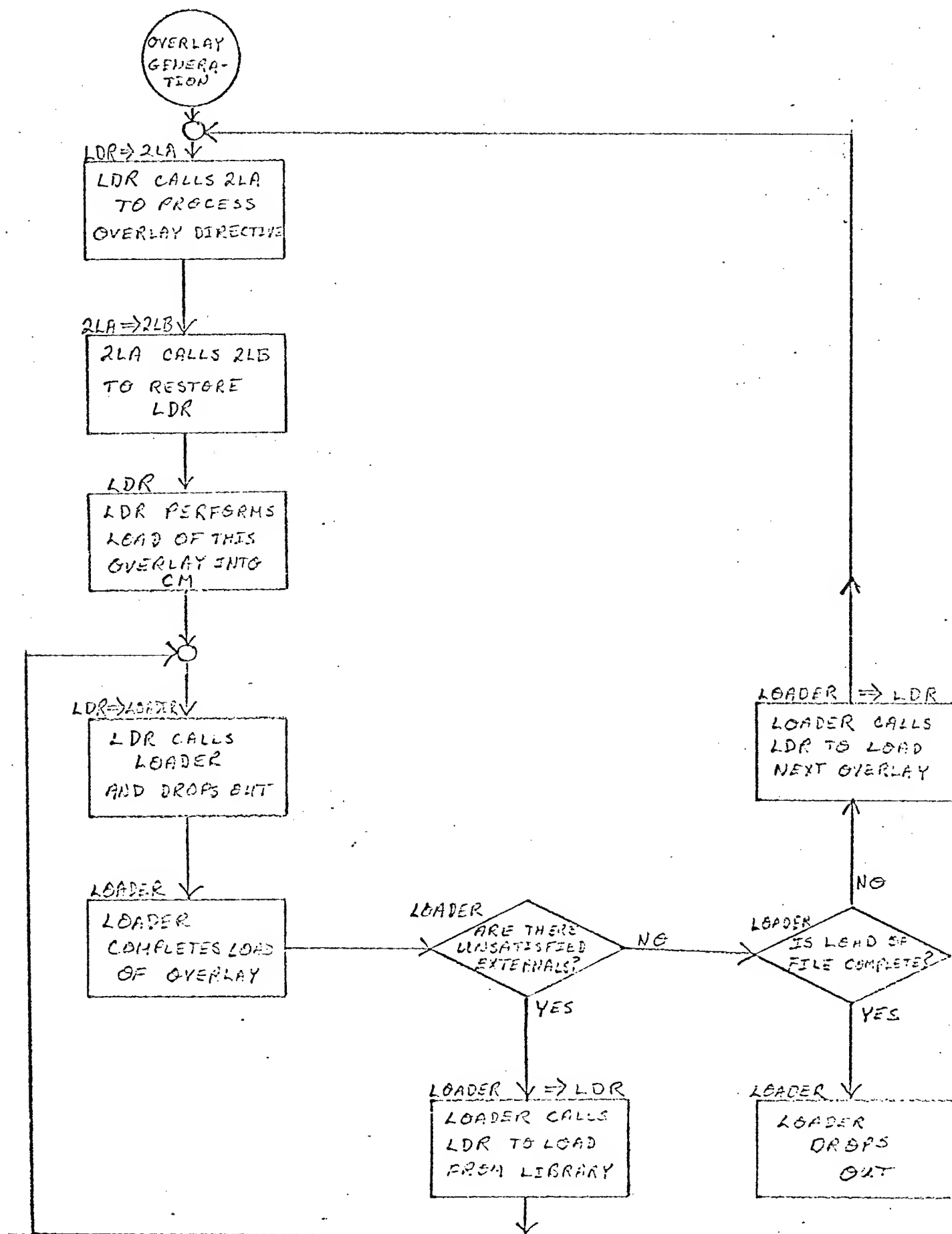


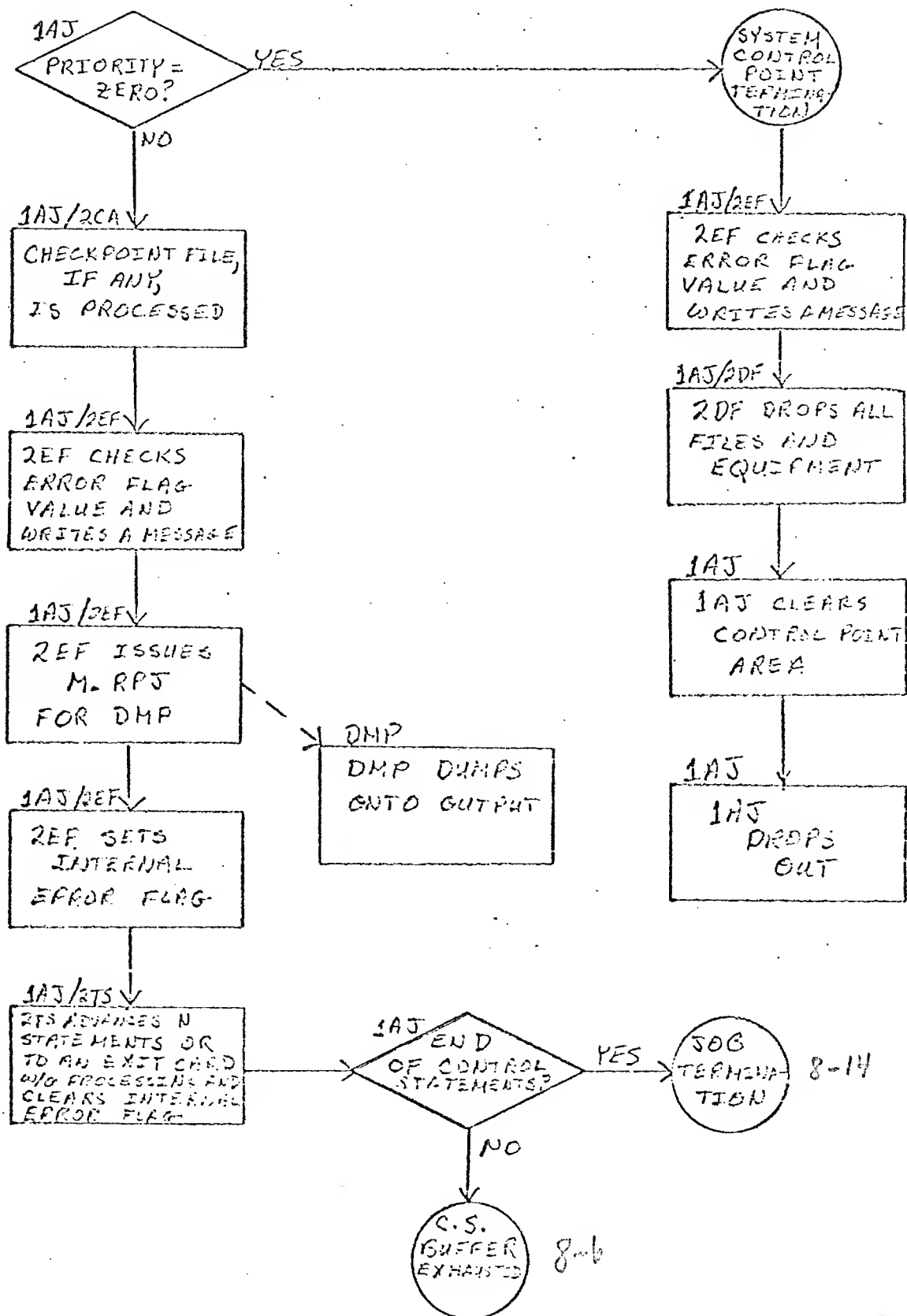
LOADER CALL (CONTINUED)

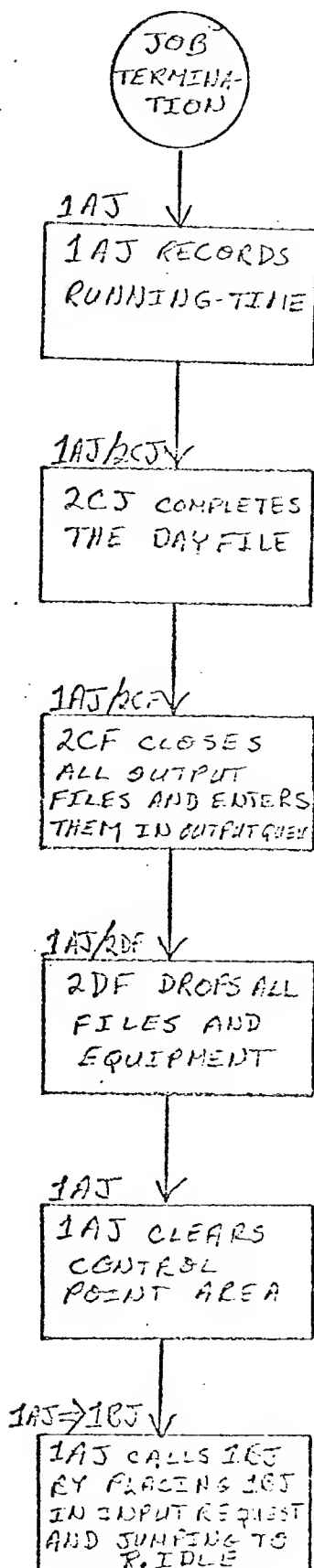
LOADER CALL (CONTINUED)

LOADER CALL (CONTINUED)

LOADER CALL (CONTINUED)

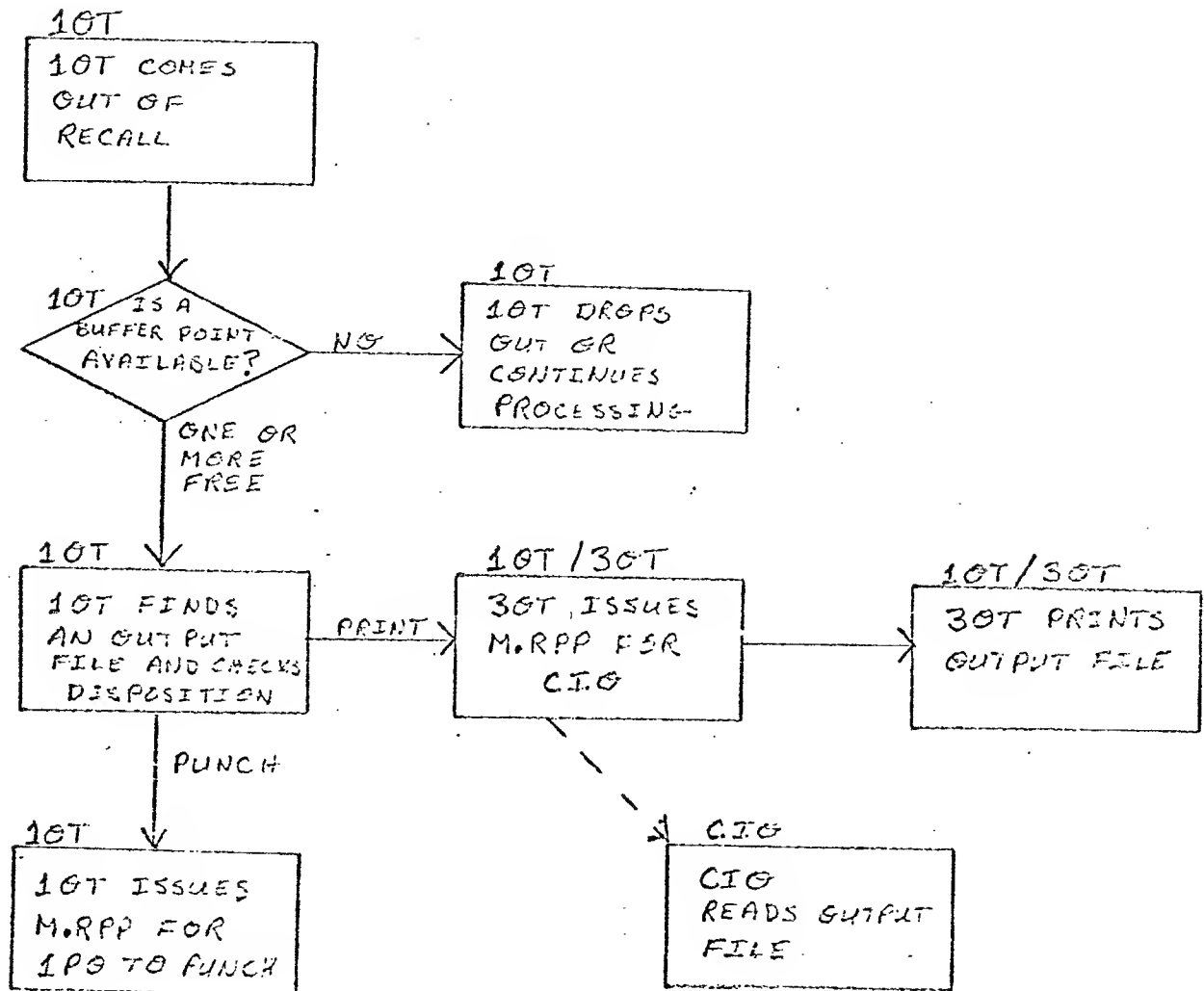


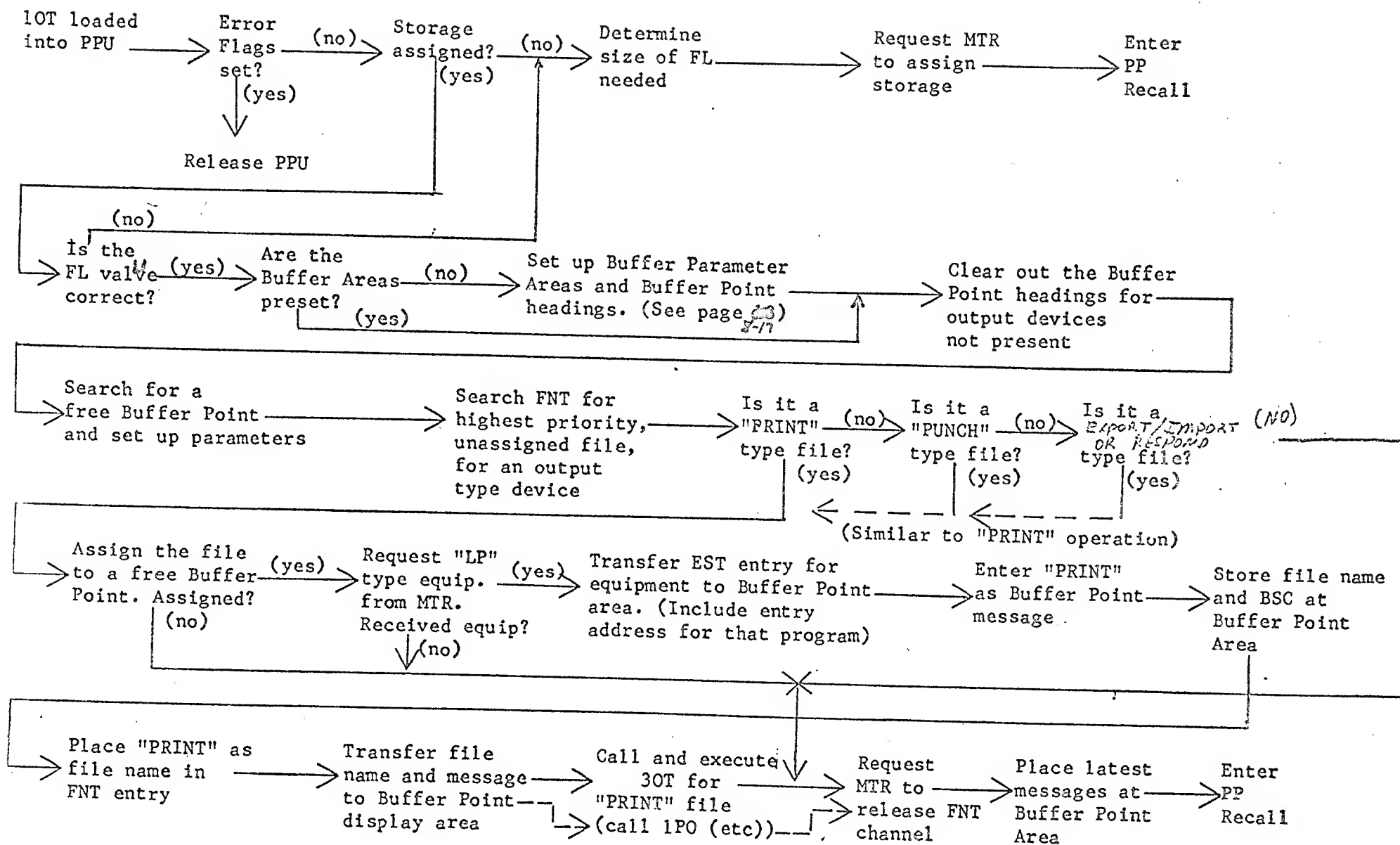
ABORTING A JOB

ABORTING A JOB (CONTINUED)

PROCESSING OUTPUT FROM A JOB
(DUMP JOB)

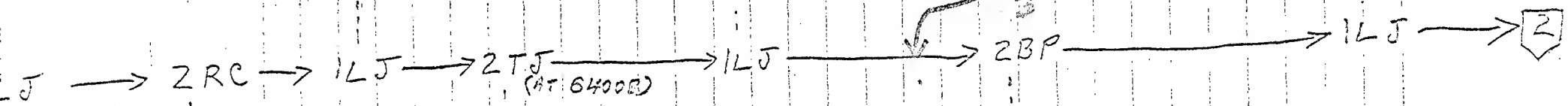
8-15





"OUTPUT" SIMPLIFIED FLOW CHART

LOAD JOB



CHECK TO SEE
IF SOMETHING IS
IN BUFFER AREA

- CODE + STATUS TO FET
- JOB NAME TO FET
- JOB NAME TO CP

- ZERO OUT ENTRY AREA
- LOAD RES

- READ 1 CARD
IN HOLLERITH

- CONVERT HOL -> EXT
BCD -> DPC, THEN
PACKED INTO BYTES
WRITE TO CM

- LOOP UNTIL EOR/EOF/BUFFER FULL

- TRANSLATE JOB CARD
- BUILDING FNT ENTRY
- BUILD PRIORITY SCHEDULE

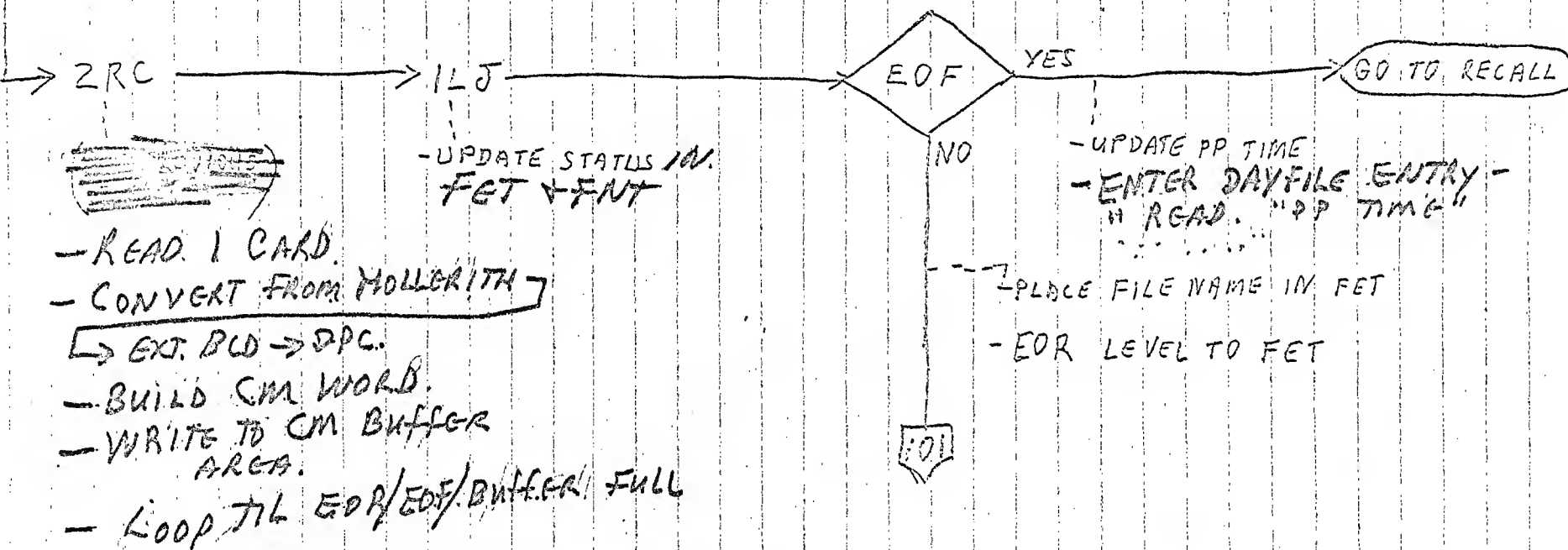
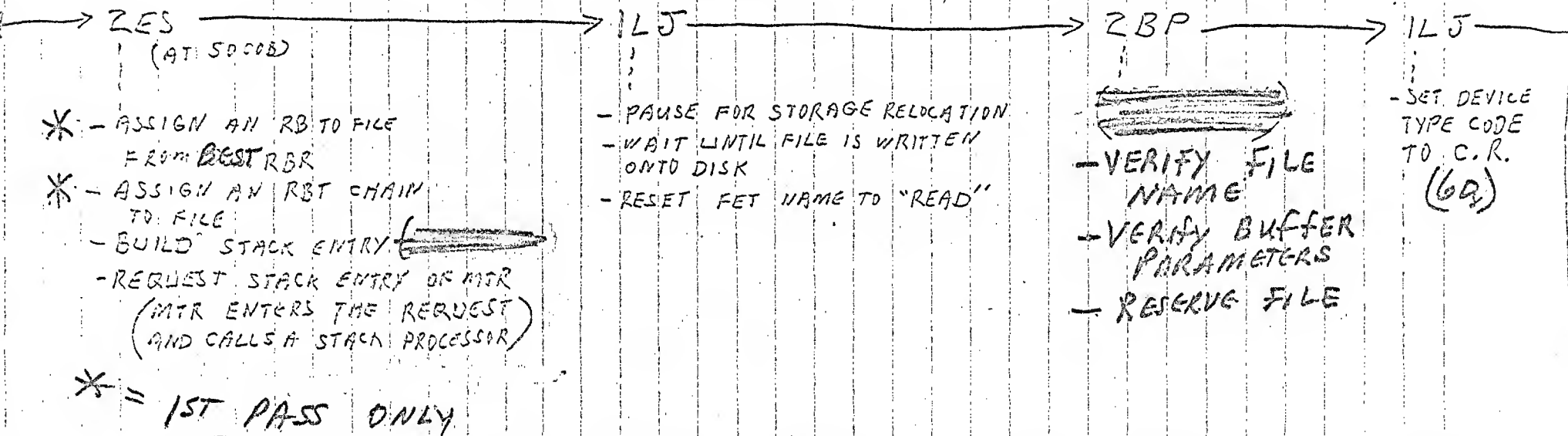
- VERIFY FILE NAME
- VERIFY BUFFER PARAMETERS
- * - ENTER FILE NAME AND CODE
+ STATUS INTO A FREE FNT ENTRY (000024)
- * - ENTER FILE AS A "LOCAL",
ASSIGNING TO CPI
- * - CLEAR FNT (2)
- * - SET FNT (3) = 0 - 01
- * - ENTER FET ADDRESS IN FNT (3)
- SET UP DEVICE TYPE IN FET

* = 1ST PASS ONLY

11 = 000 004 01 0100 = FIRST
12 = 0 205 = IN
13 = 0 00100 = OUT
14 = 0 2100 = LIMIT
15 = UNCHANGED
20 =
21 3100
22 3100
23 3100
24 2300
25 3100

RA+100 = 1ST CARD (10 CHAR)
101 = ABC, P7, T100, CM 400 000.
102 = 0 - 0

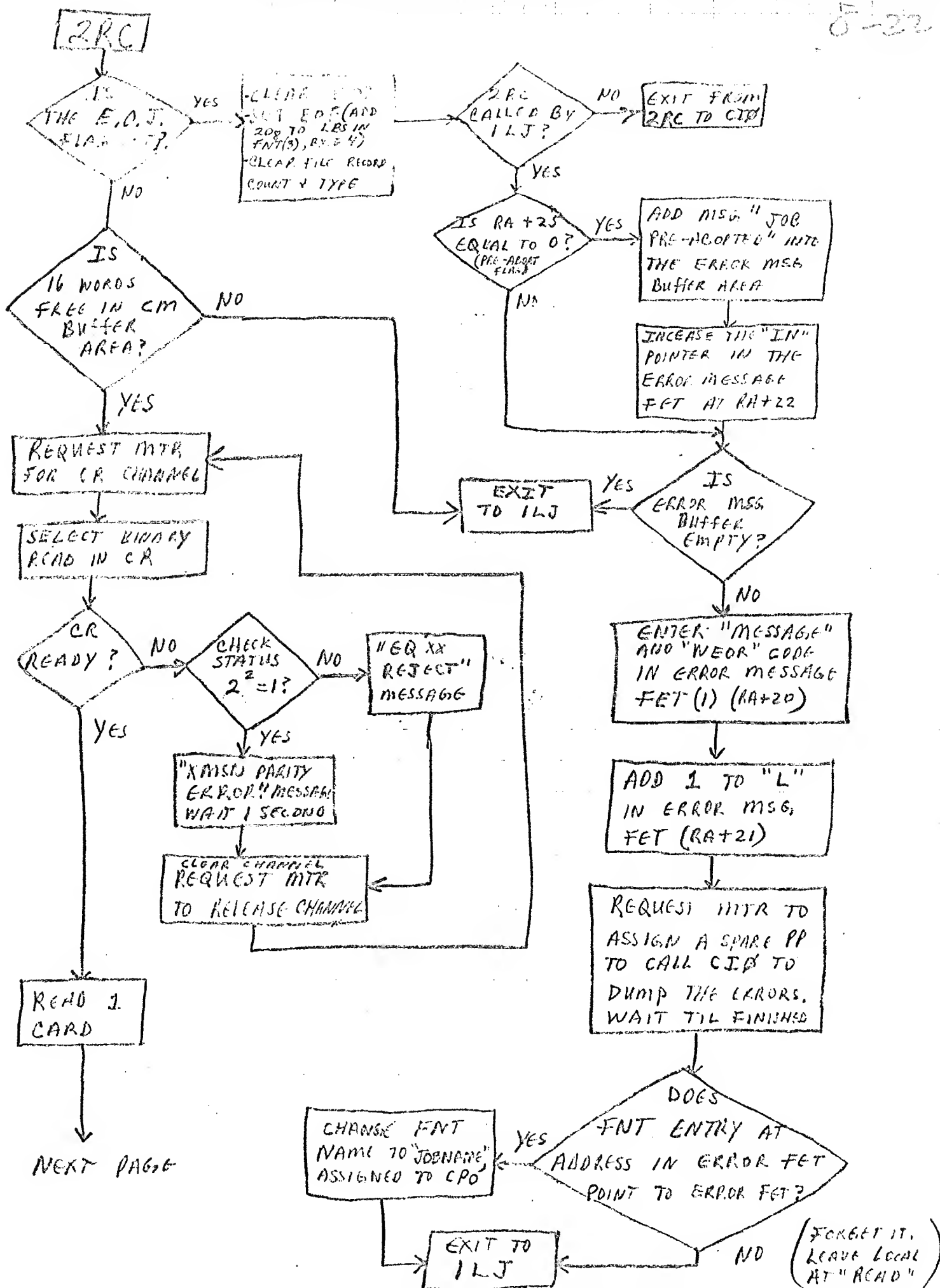
8-17-8
8-18-8



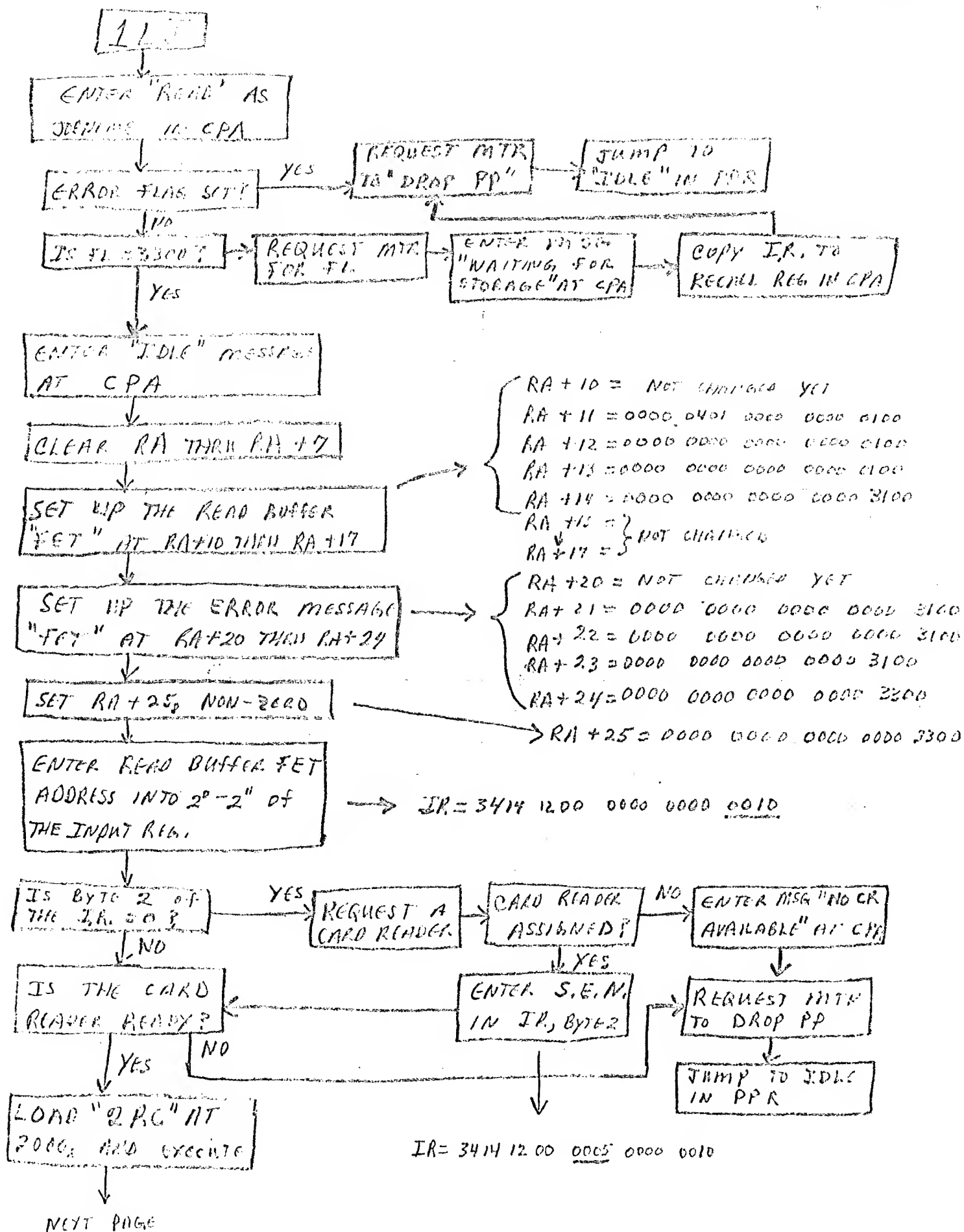
8-20

2ES Parameters

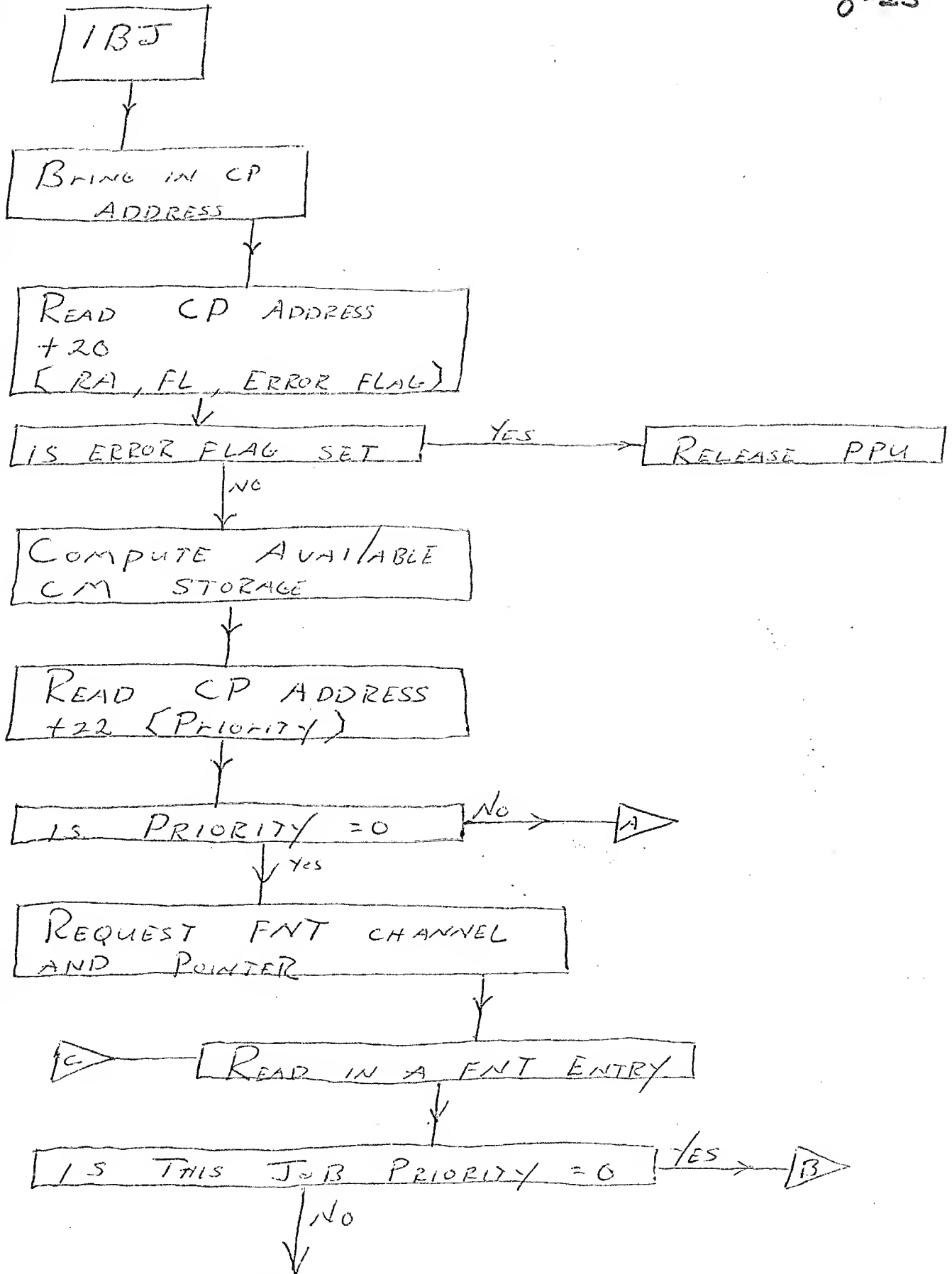
<u>PP Call</u>		
D.FNT	20	
	21	
	22	FNT(2)
	23	
	24	
	25	FNT(3)
	26	
	27	
D.EST	30	← level number in low 6 bits
	31	← OP Code Also to FNT(3)
	32	Flags 4=NOFET, 1=PP AVAIL
	33	CP RECALL 1=RECALL, 0=NO RECALL (Ignored if D.CPAD=0)
D.BA	40	First word of FET, 2nd word of FET If NO FET BIT If open function Is not set (for random bit)
	41	
	42	
	43	
	44	
D.FR5	45	Last buffer status
D.PPIRB	50	
	51	
	52	Number of records FET address
	53	
	54	
D.FA	57	
D.FIRST	60	First
	61	
D.LIMIT	66	Limit
	67	
D.CPAD	74	Control point address



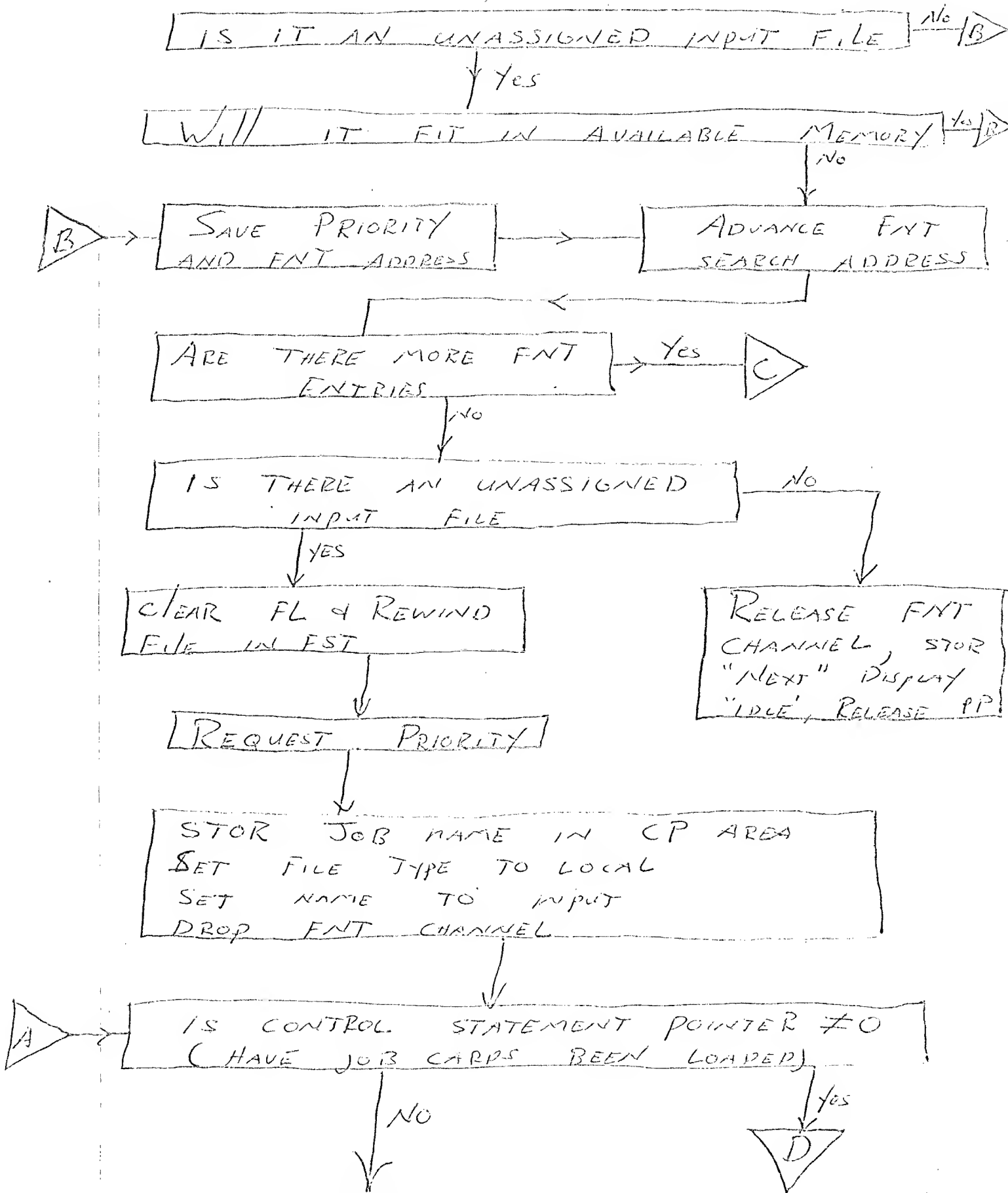
5-27



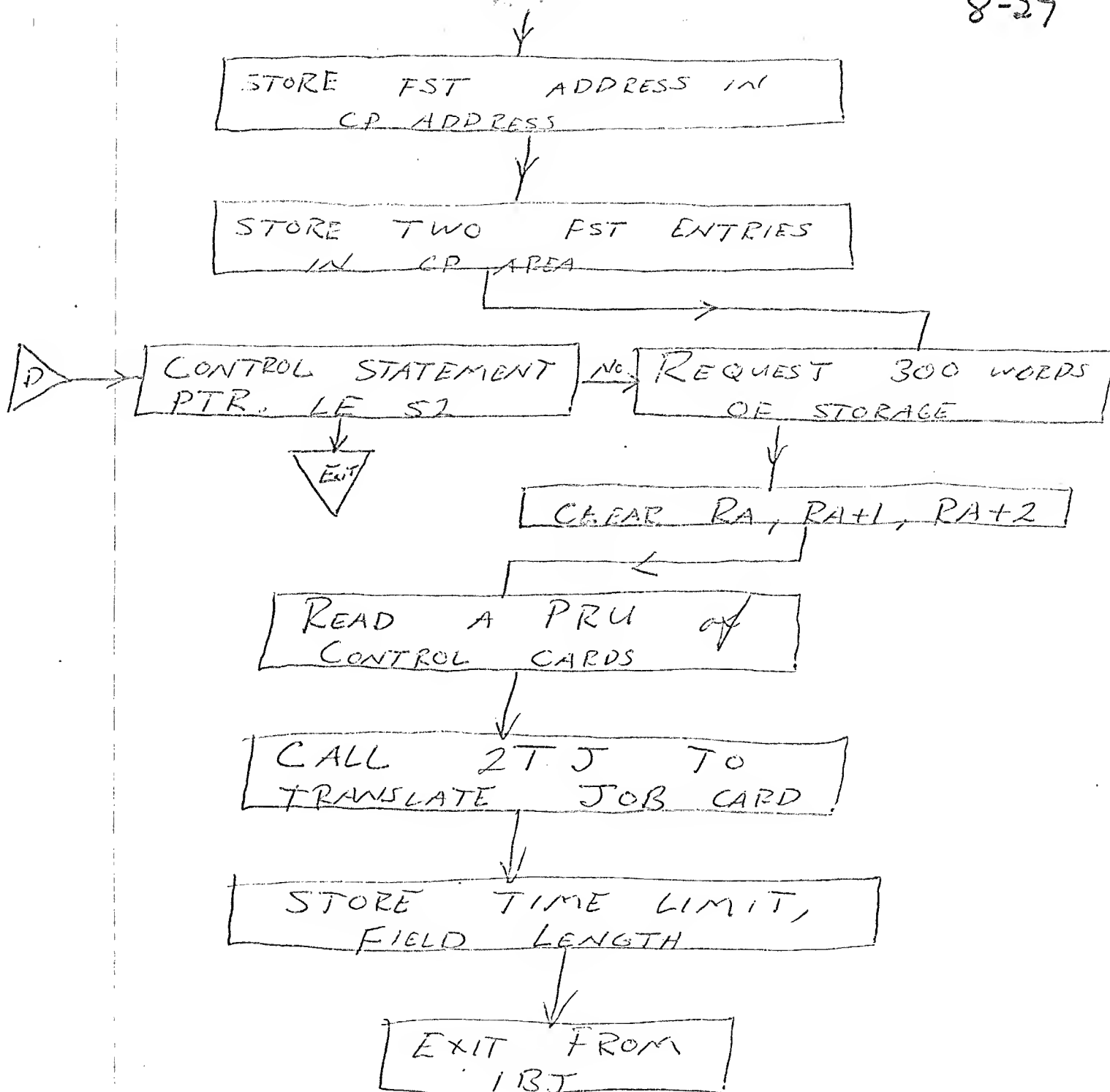
8-25

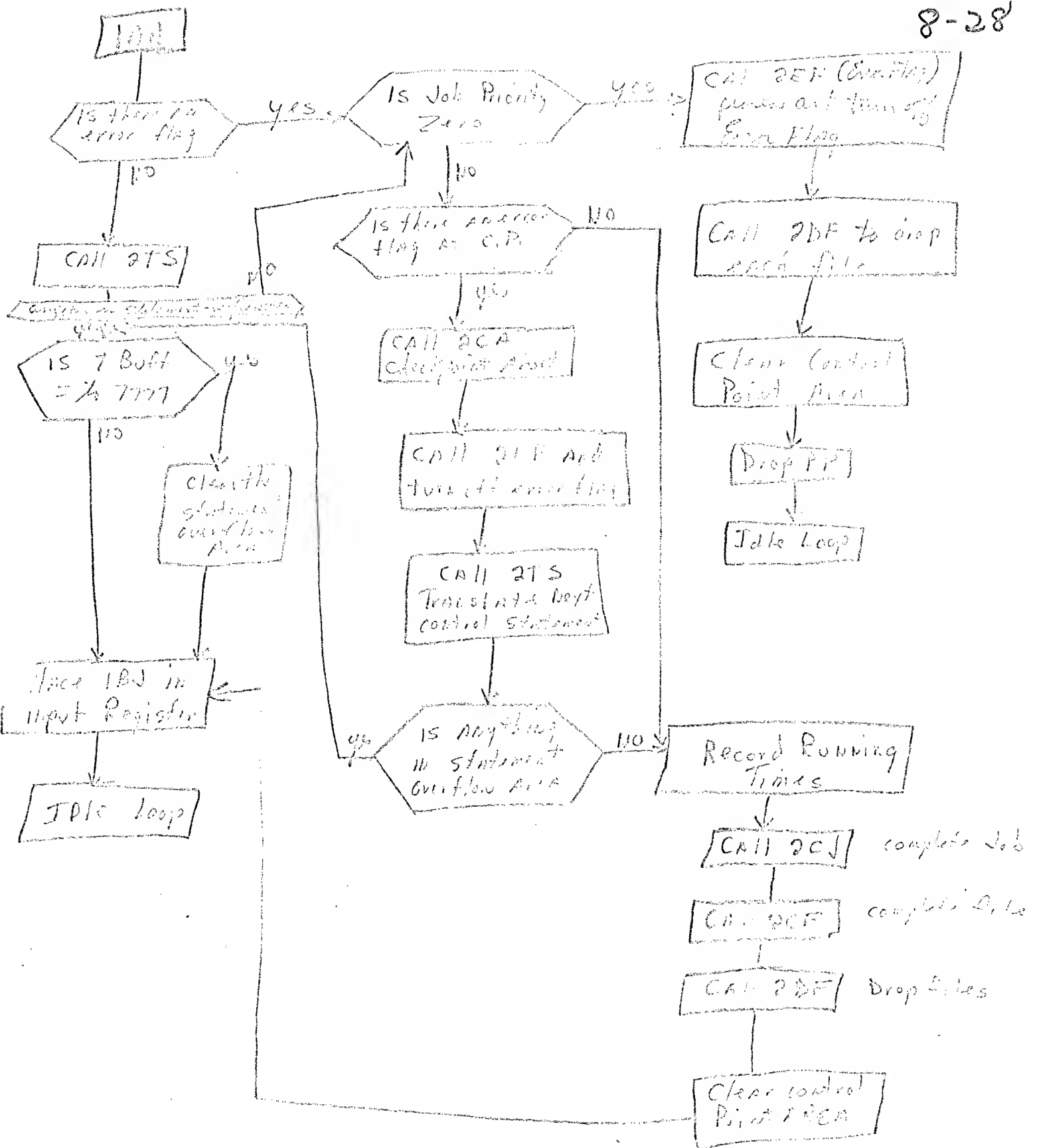


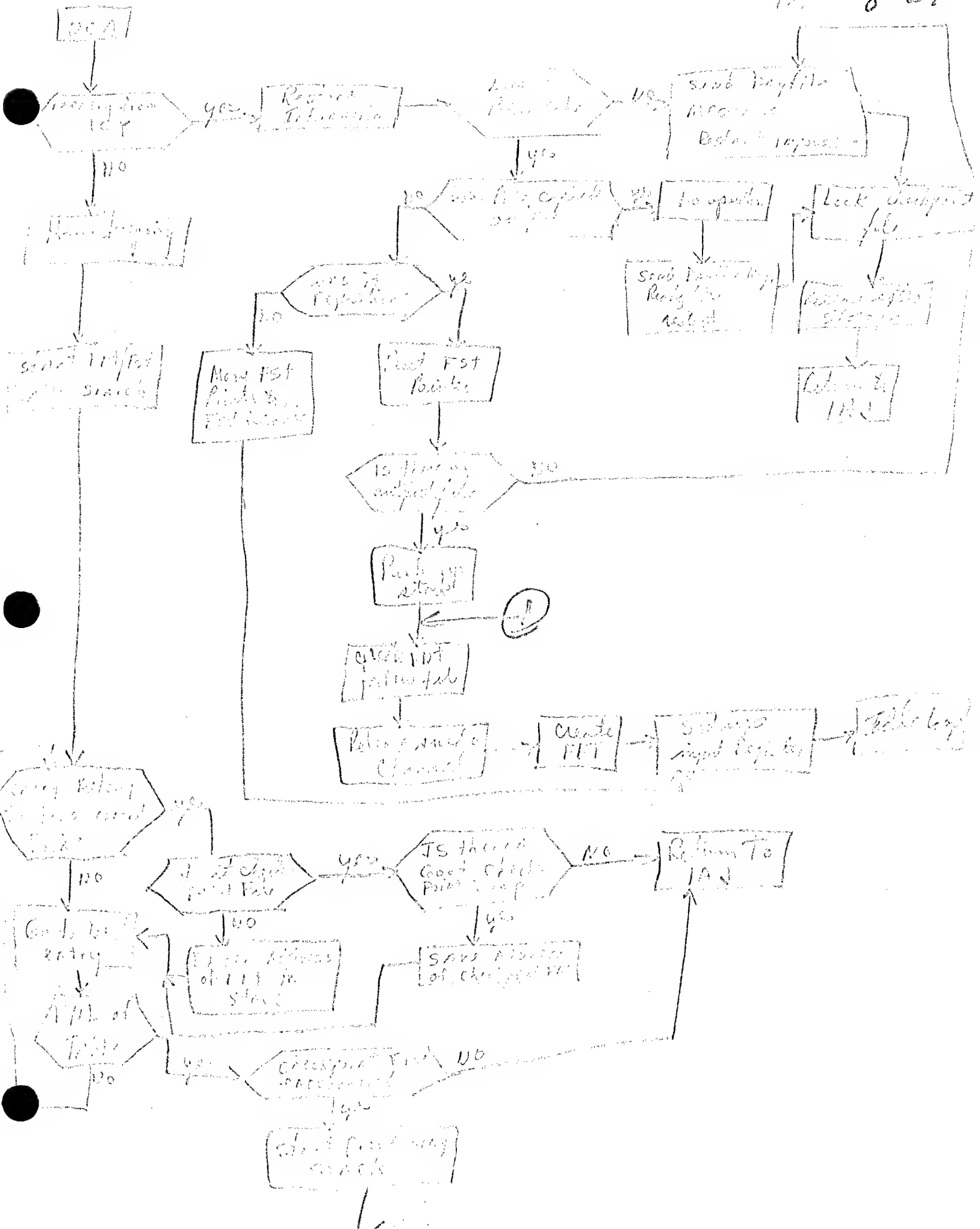
8-26

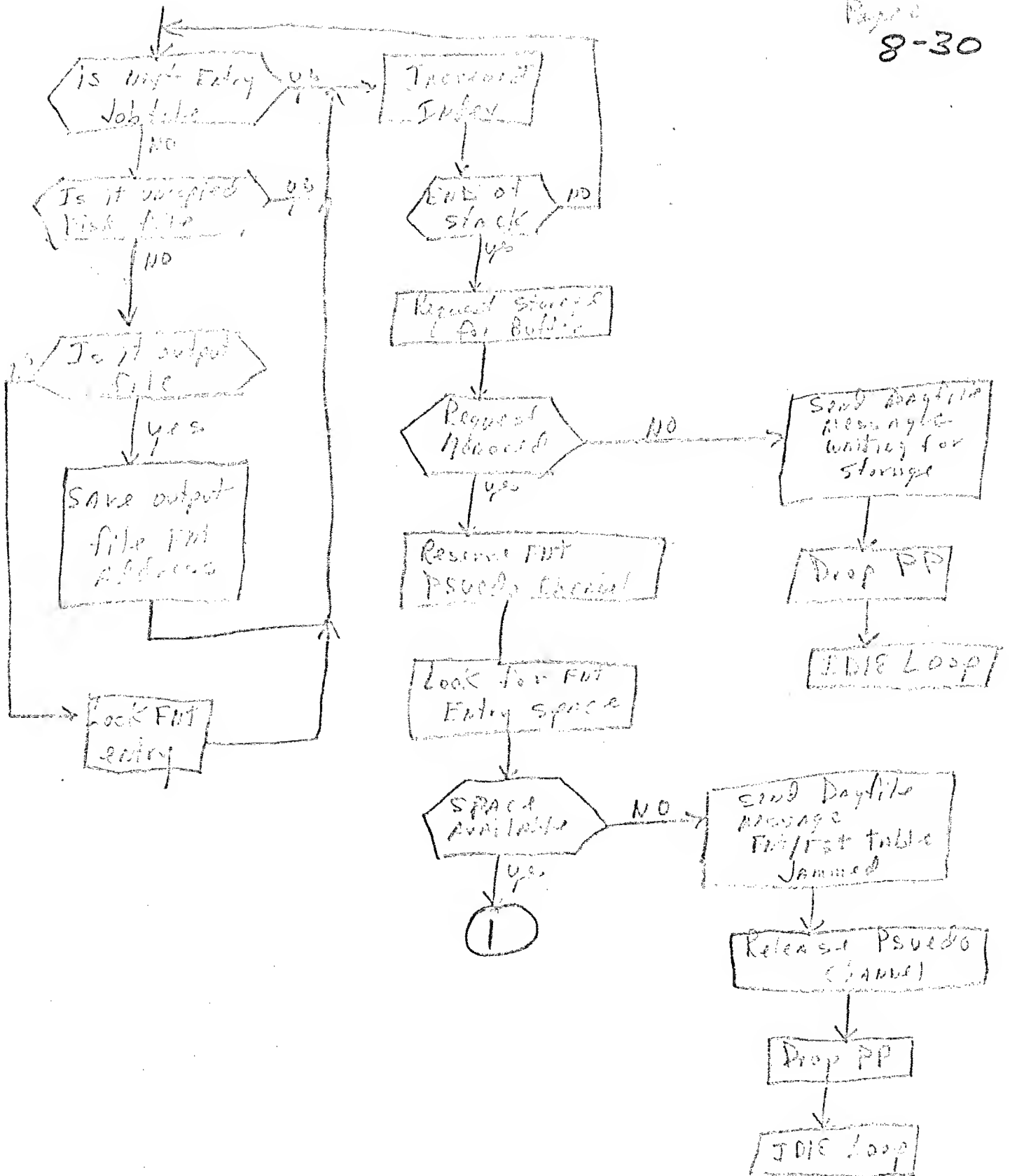


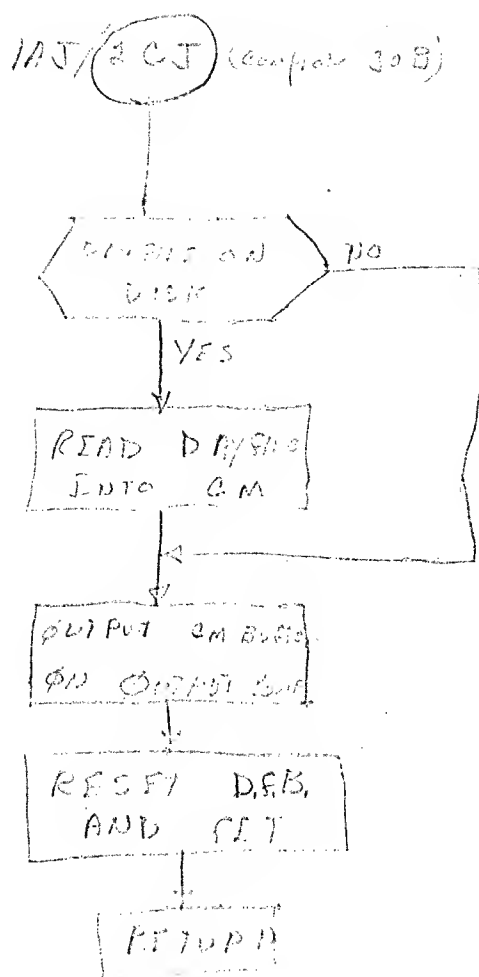
8-27



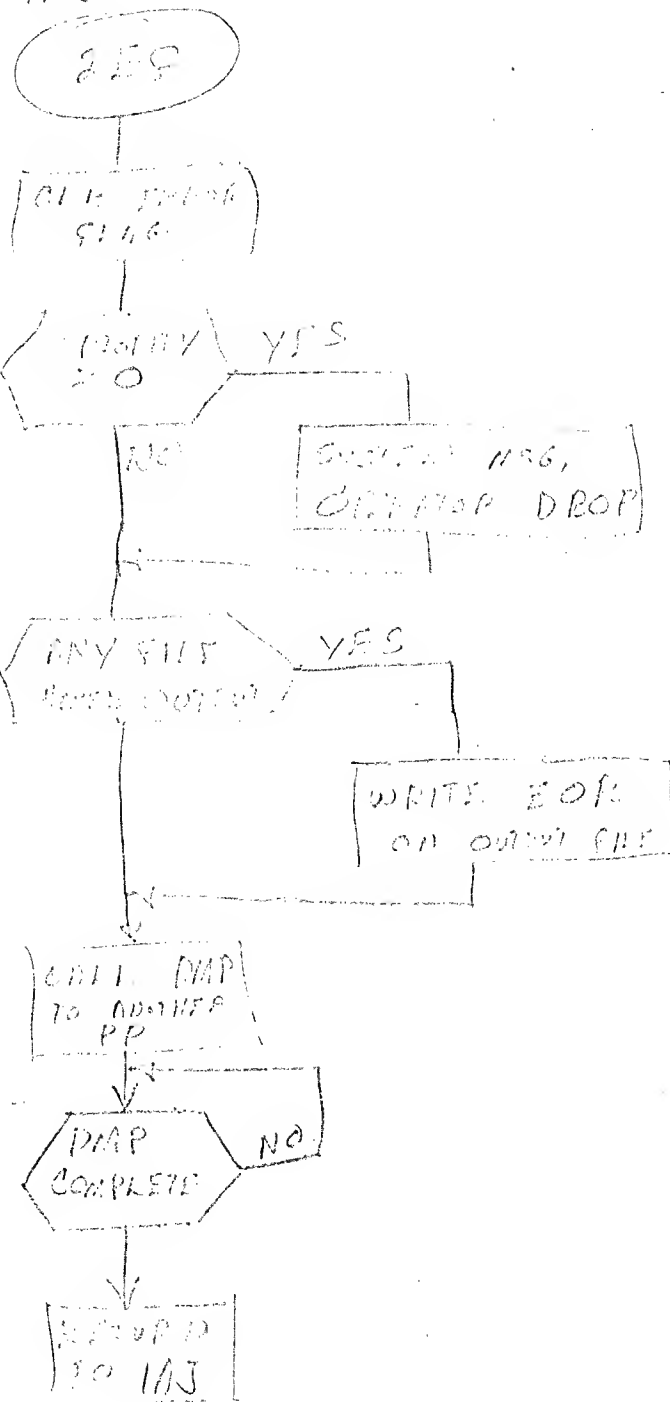








103



IAJ

2CF

SETUP MODEL
IN D.BA -> D.BOL

YES

D.FNT = 0

NO

GET DEATH
Pointer ENTRY

PER
IDF

YES

DISP.
Code = 0

NO

RJM.
RLS

SET V.FNT
To 10000000

EXIT

AXI FNT ENTRY
FOR THE G.P.

NO

YES

PER
IDF

DISP. Code
= 0

YES

NO

PRINT
TYPE

NO

YES

MODEL
IN USE

YES

NO

SAVE ENTRY
ADDRESS
IN ADDRESS

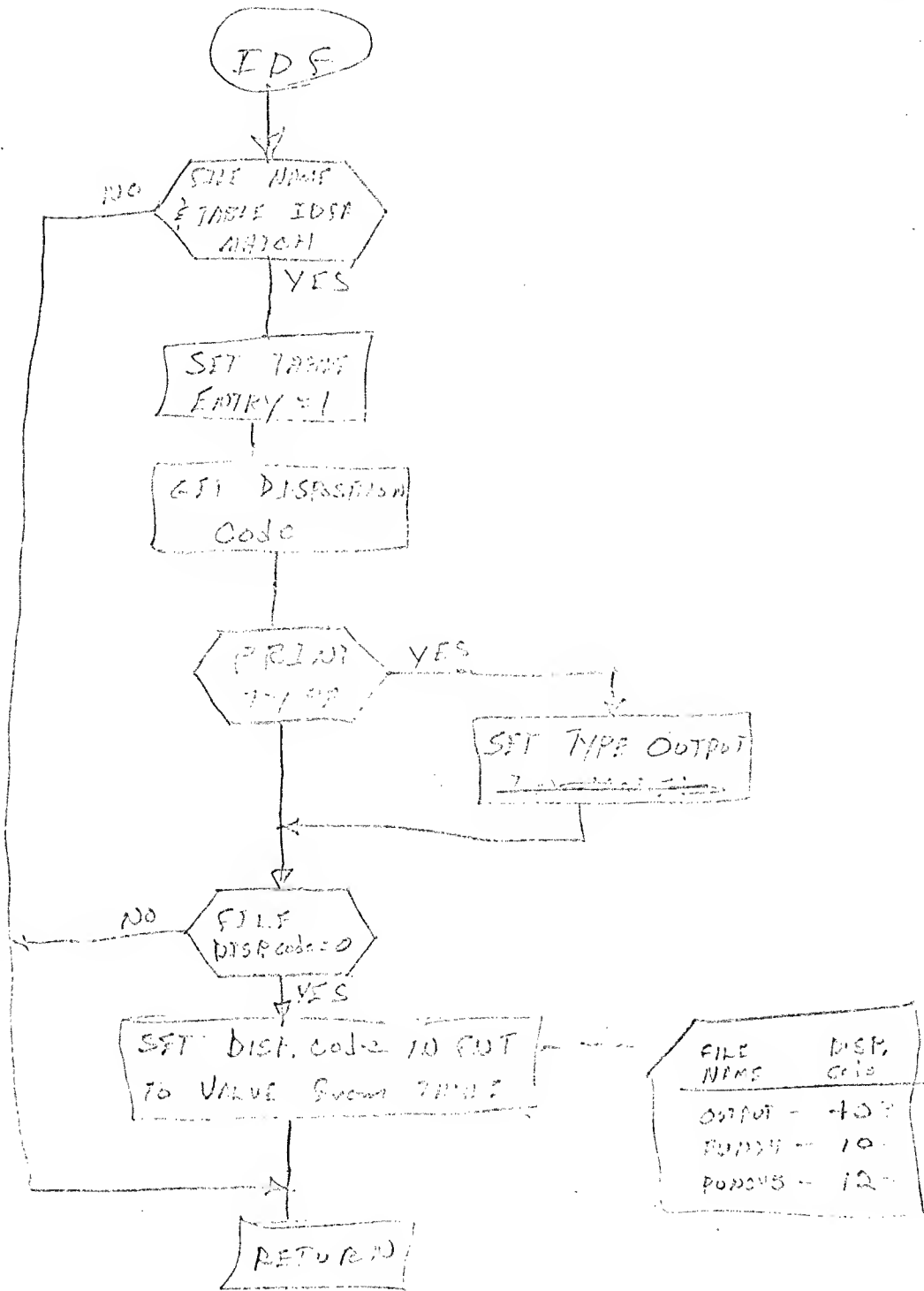
END of
FNT

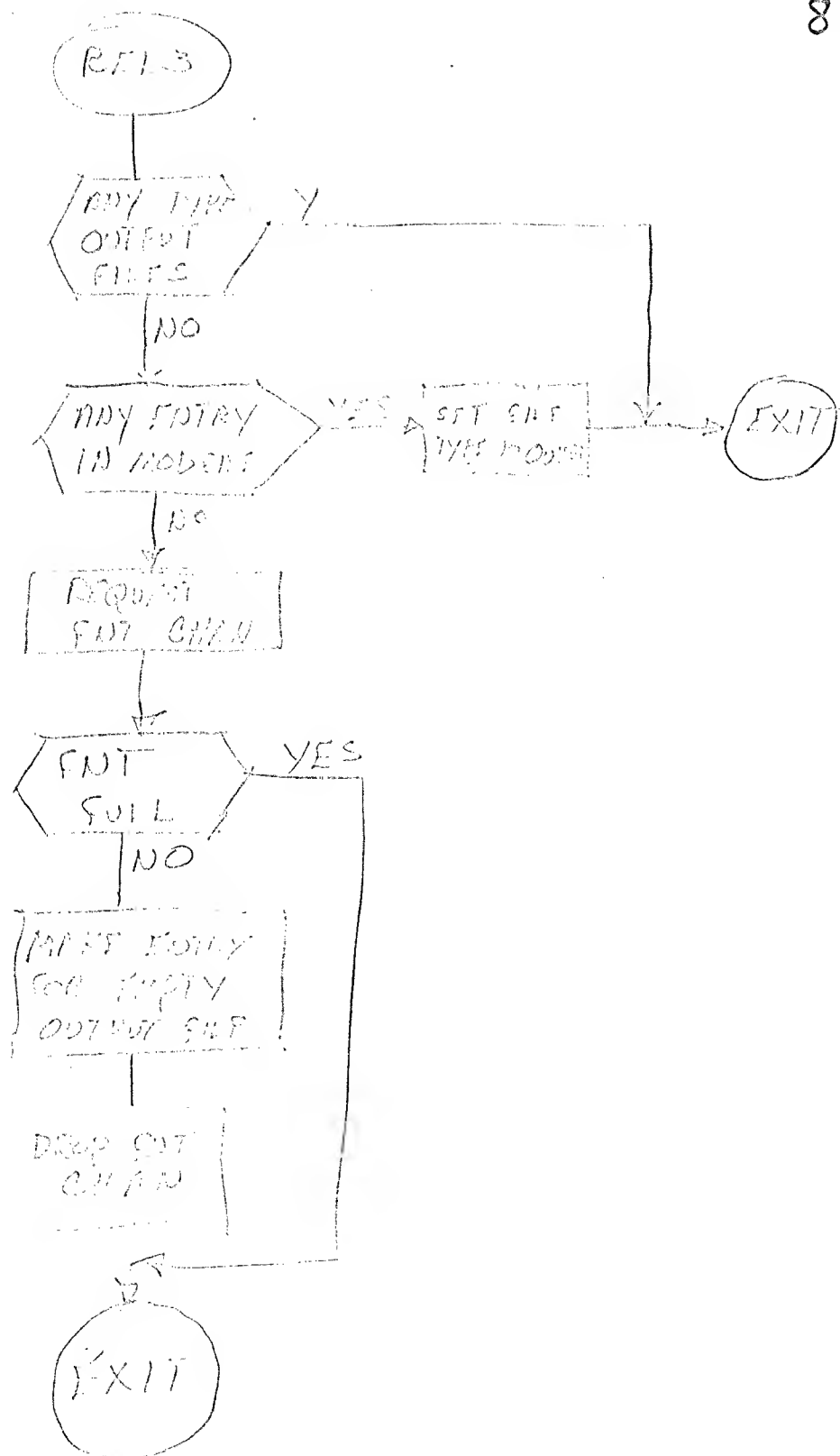
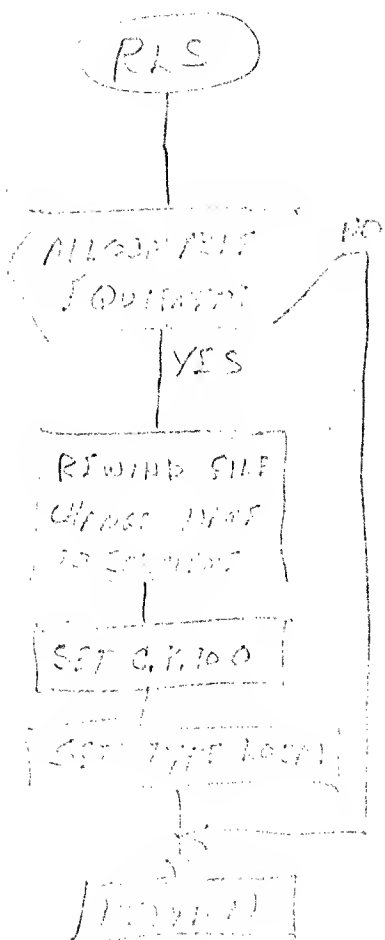
NO

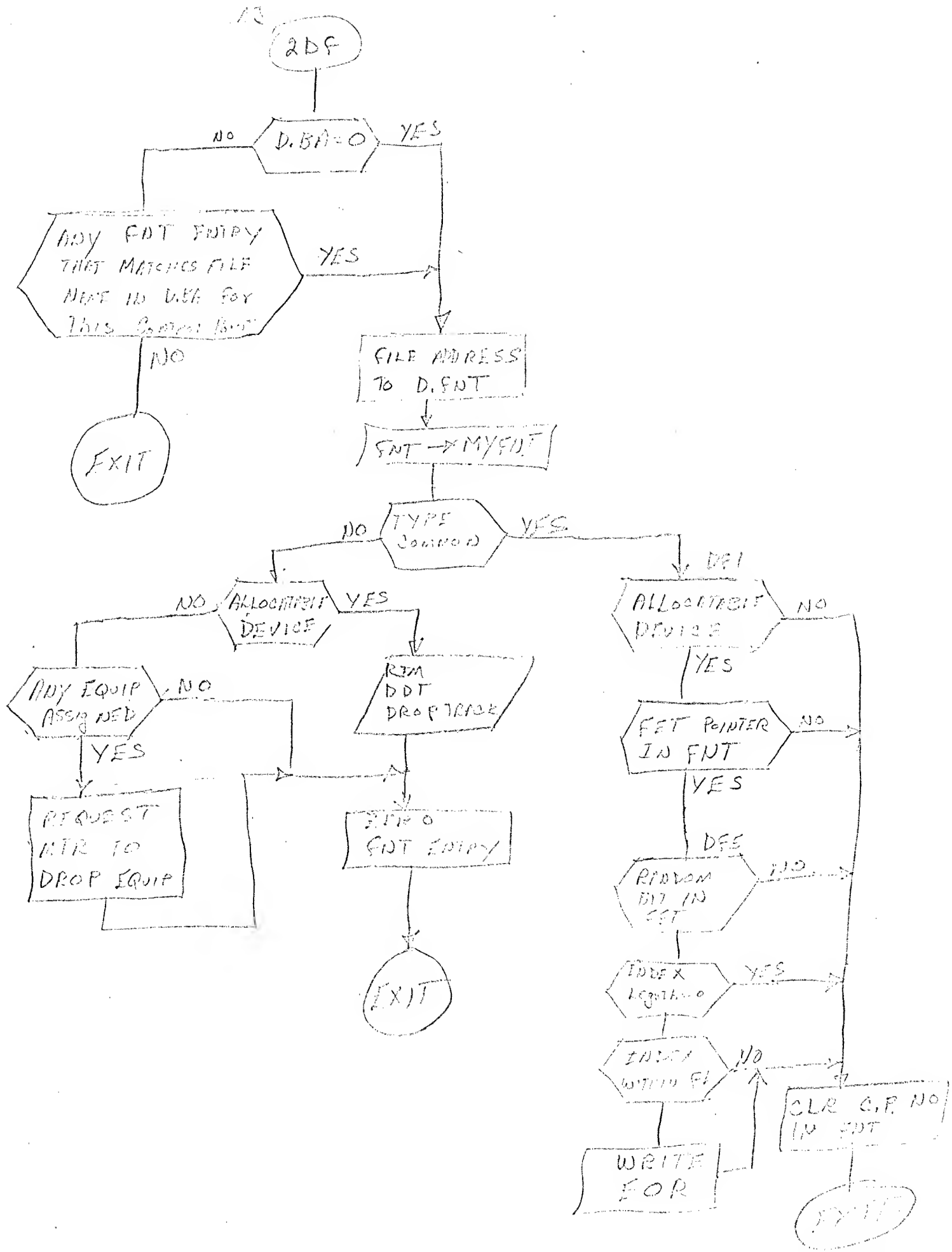
YES

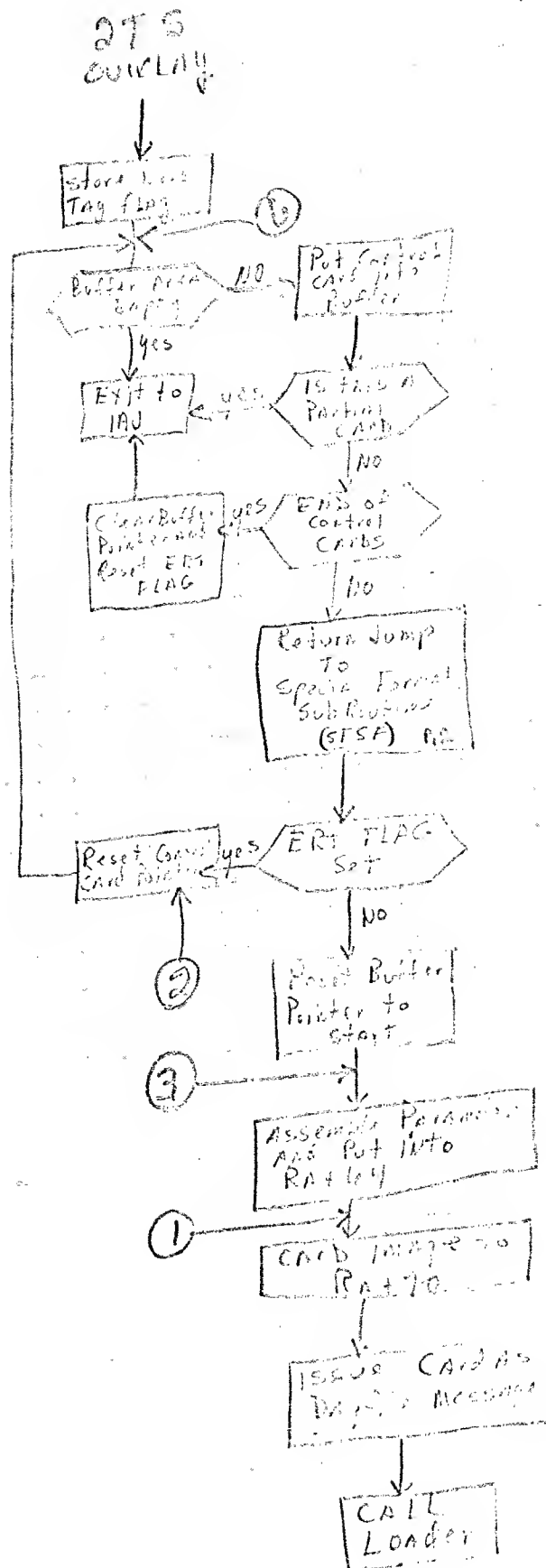
REL 3

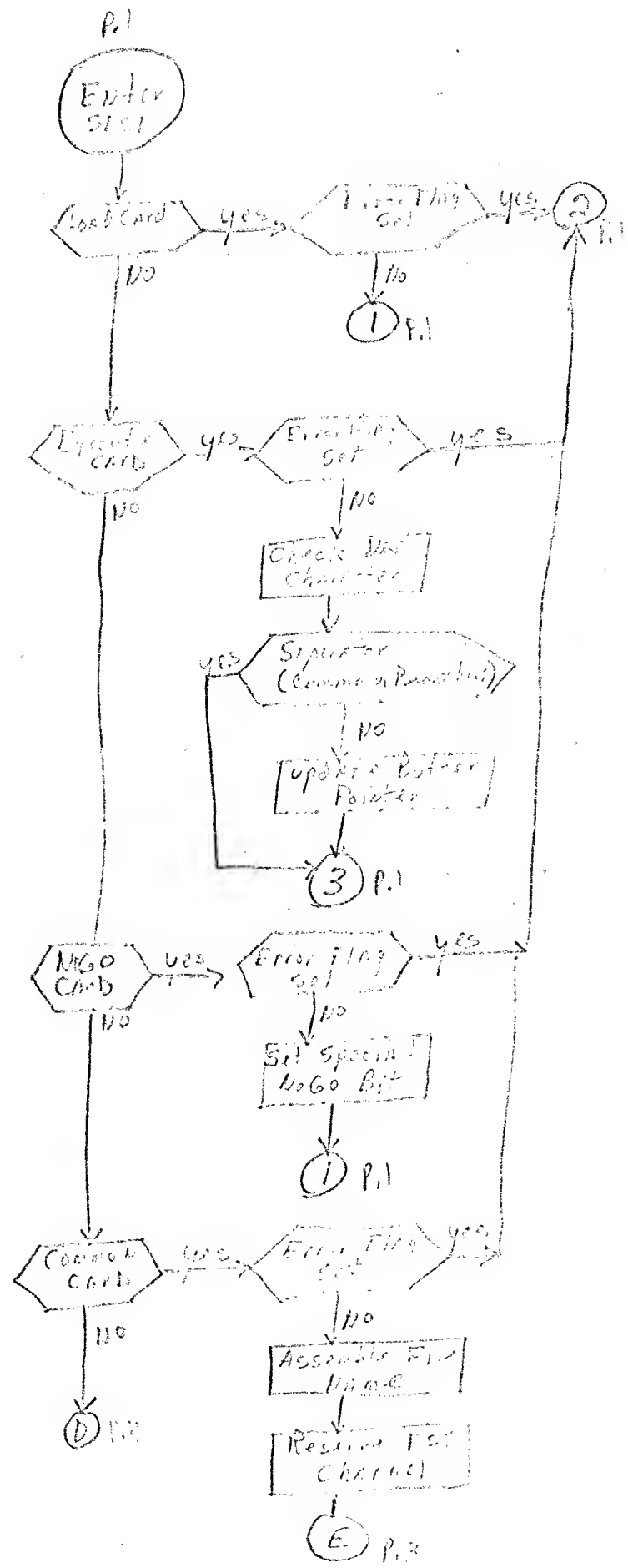
RJM
RLS

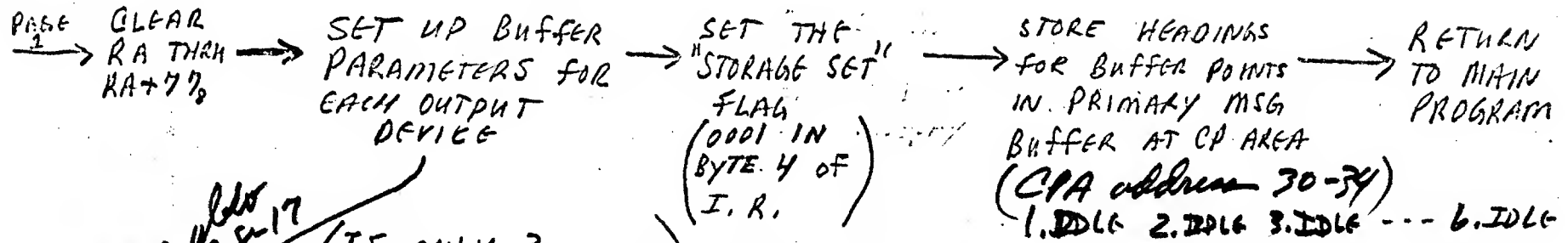












see also P. 8-17
(IF ONLY 3 OUTPUT DEVICES, THEN SET UP ONLY 3, ETC.)

RA+11 = 0 — 0100 ⇒ FIRST
RA+12 = 0 — 0100 ⇒ IN
RA+13 = 0 — 0100 ⇒ OUT
RA+14 = 0 — 01100 ⇒ LIMIT } OUTPUT DEVICE AT BUFFER POINT #1

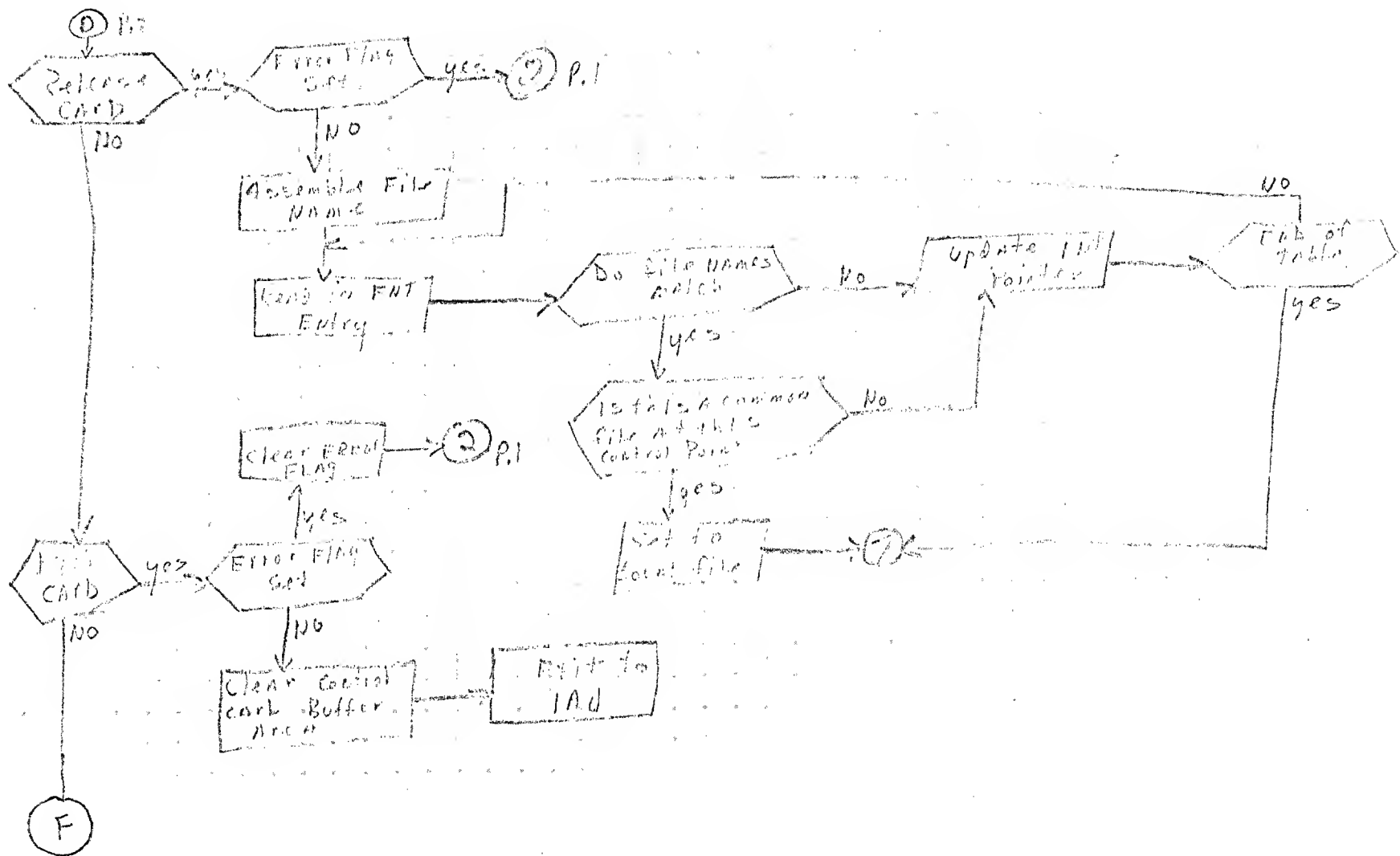
RA+21 = 0 — 01100 ⇒ FIRST
RA+22 = 0 — 01100 ⇒ IN
RA+23 = 0 — 01100 ⇒ OUT
RA+24 = 0 — 02100 ⇒ LIMIT } OUTPUT DEVICE AT BUFFER POINT #2

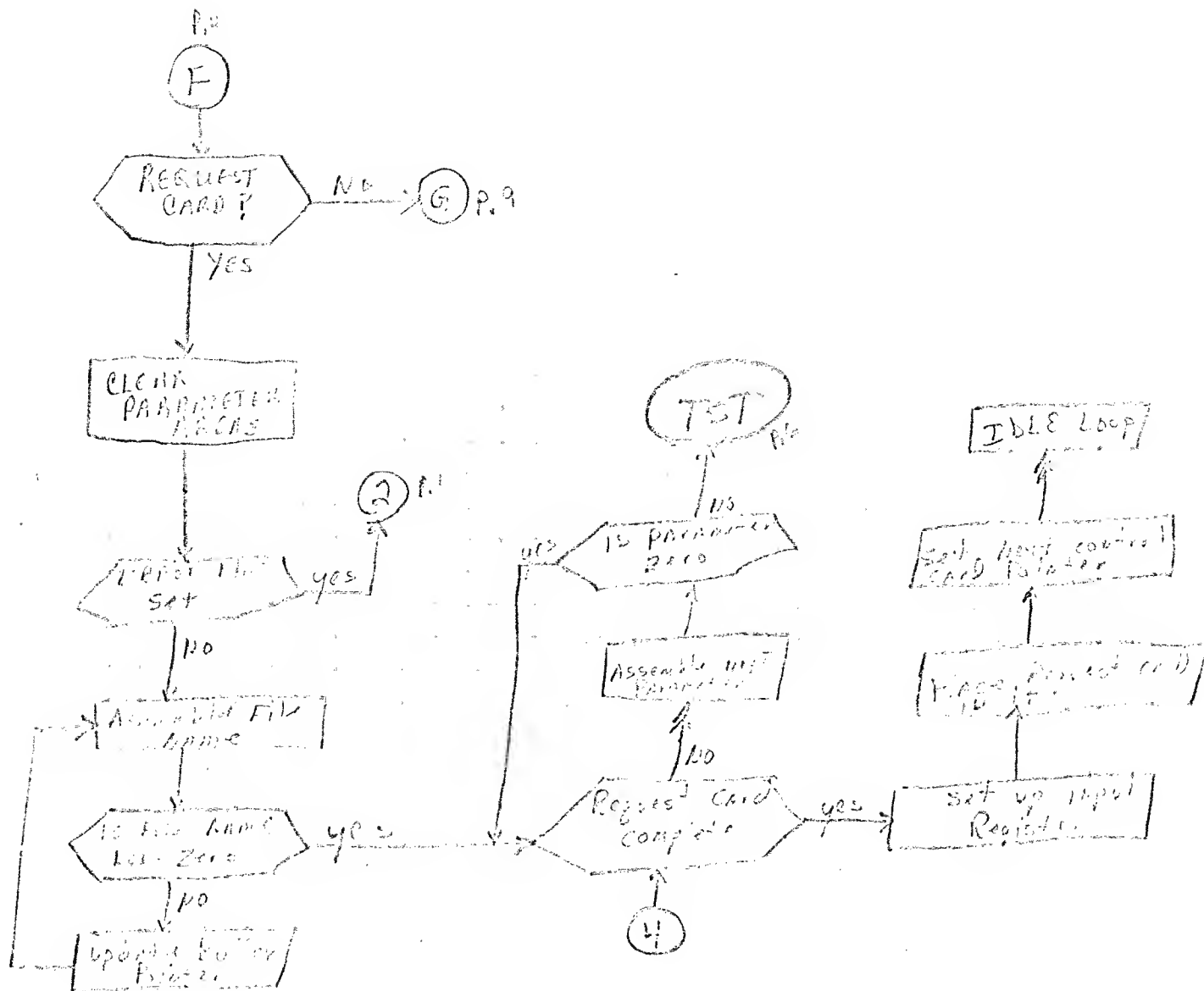
RA+31 = 0 — 01100 ⇒ FIRST
RA+32 = 0 — 01100 ⇒ IN
RA+33 = 0 — 01100 ⇒ OUT
RA+34 = 0 — 02100 ⇒ LIMIT } OUTPUT DEVICE AT BUFFER POINT #3

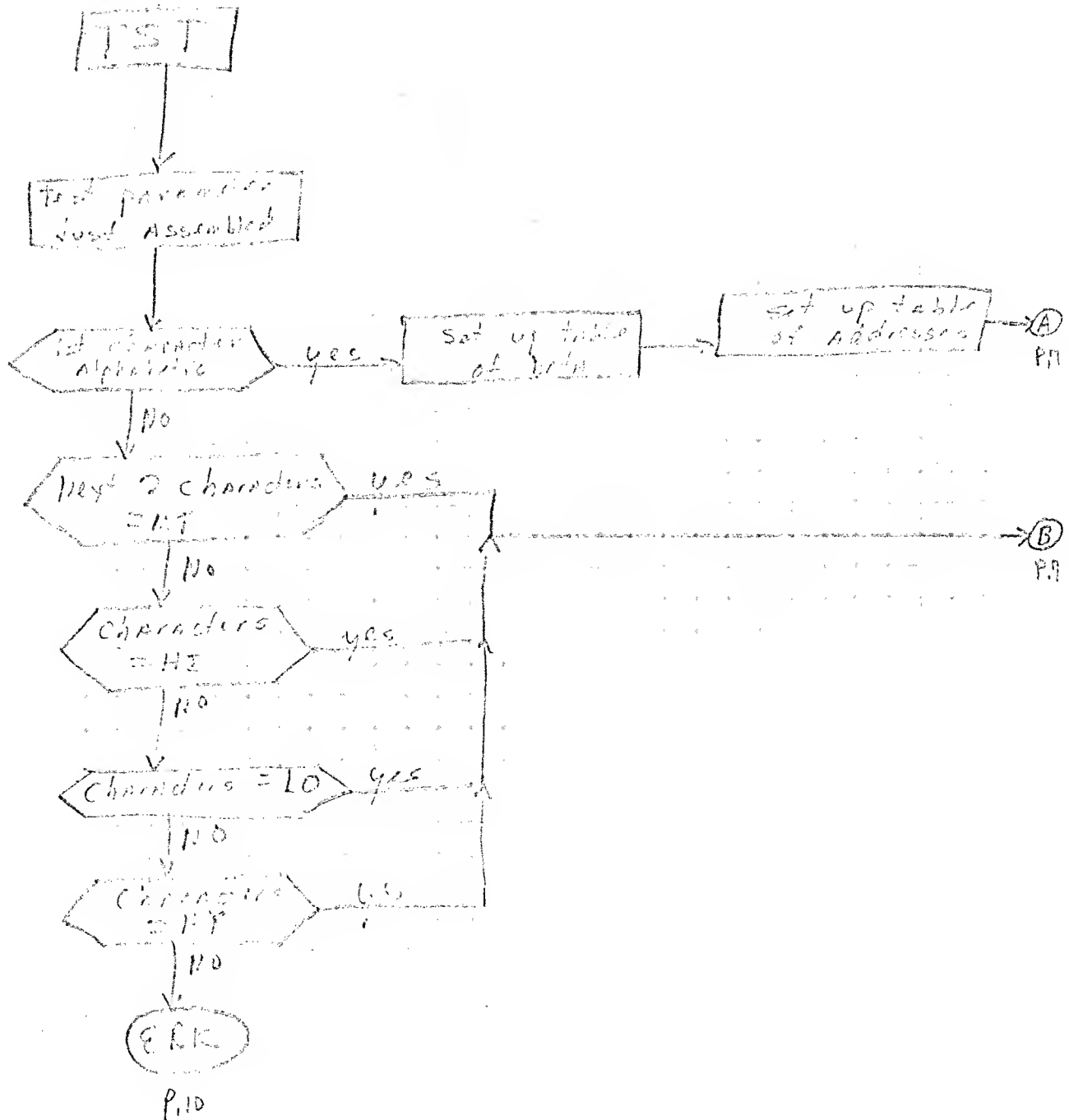
RA+41 = 0 — 01100 ⇒ FIRST
RA+42 = 0 — 01100 ⇒ IN
RA+43 = 0 — 01100 ⇒ OUT
RA+44 = 0 — 02100 ⇒ LIMIT } OUTPUT DEVICE AT BUFFER POINT #4

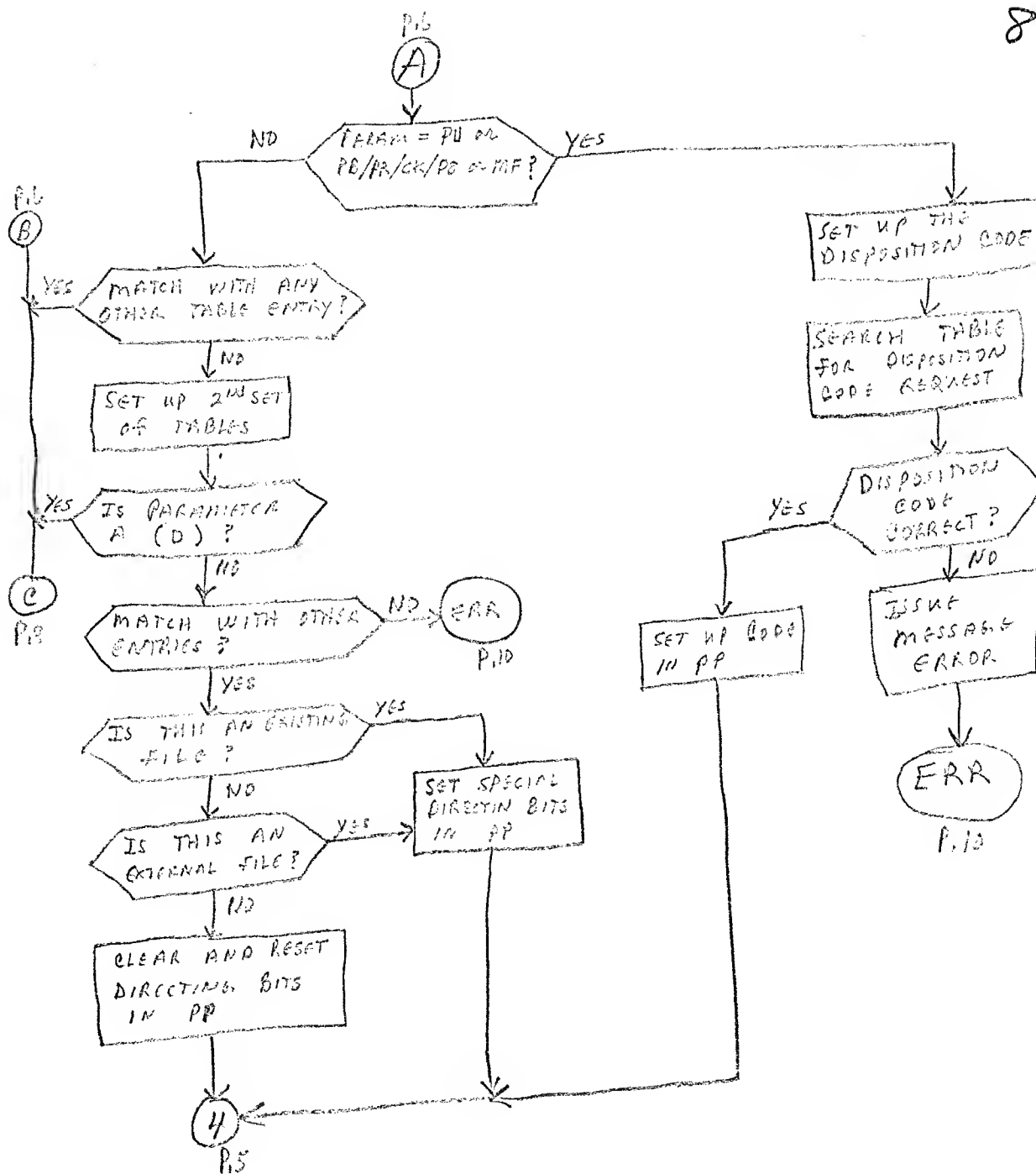
RA+51 = 0 — 01100 ⇒ FIRST
RA+52 = 0 — 01100 ⇒ IN
RA+53 = 0 — 01100 ⇒ OUT
RA+54 = 0 — 02100 ⇒ LIMIT } OUTPUT DEVICE AT BUFFER POINT #5

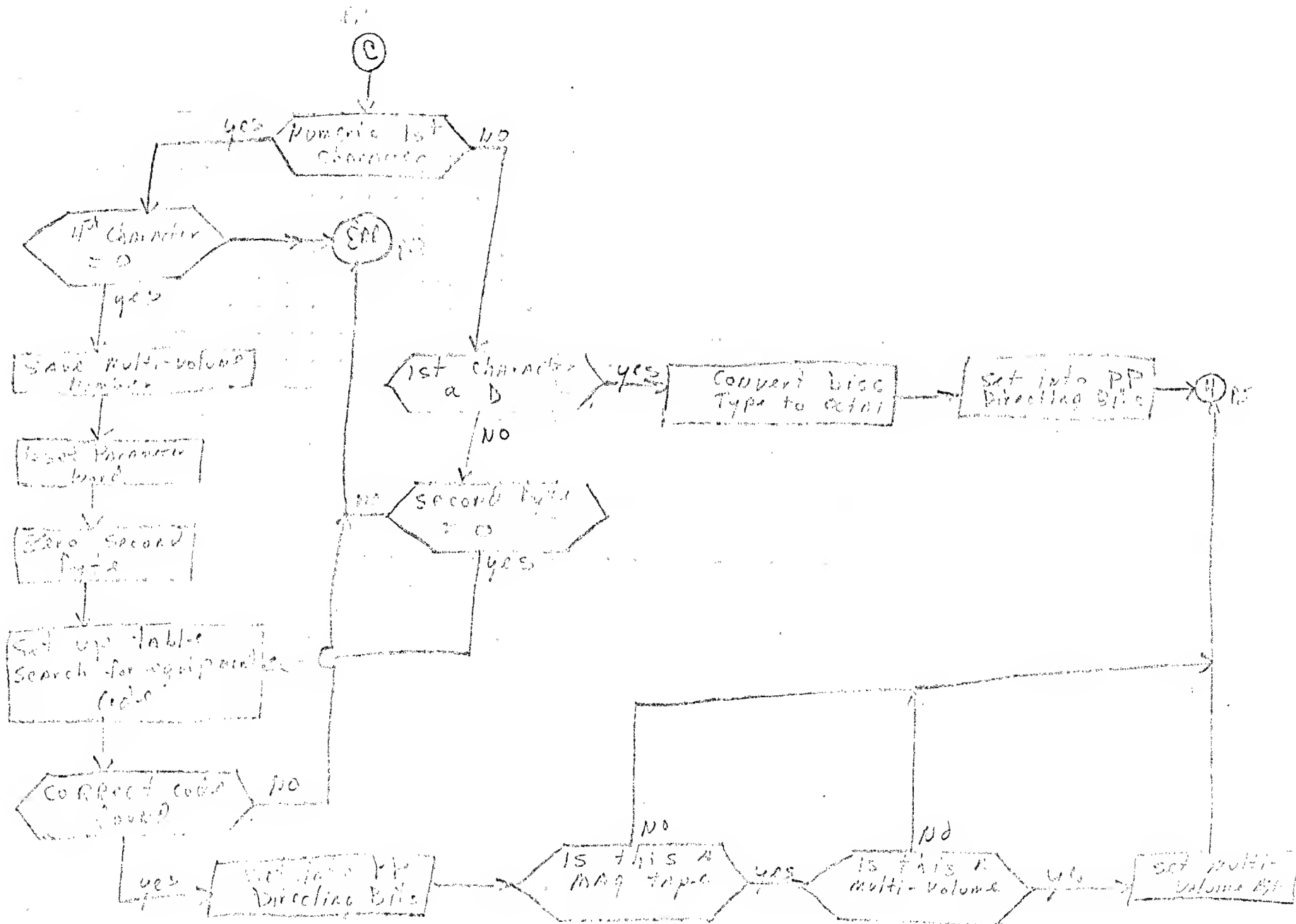
RA+61 = 0 — 05100 ⇒ FIRST
RA+62 = 0 — 05100 ⇒ IN
RA+63 = 0 — 05100 ⇒ OUT
RA+64 = 0 — 06100 ⇒ LIMIT } OUTPUT DEVICE AT BUFFER POINT #6

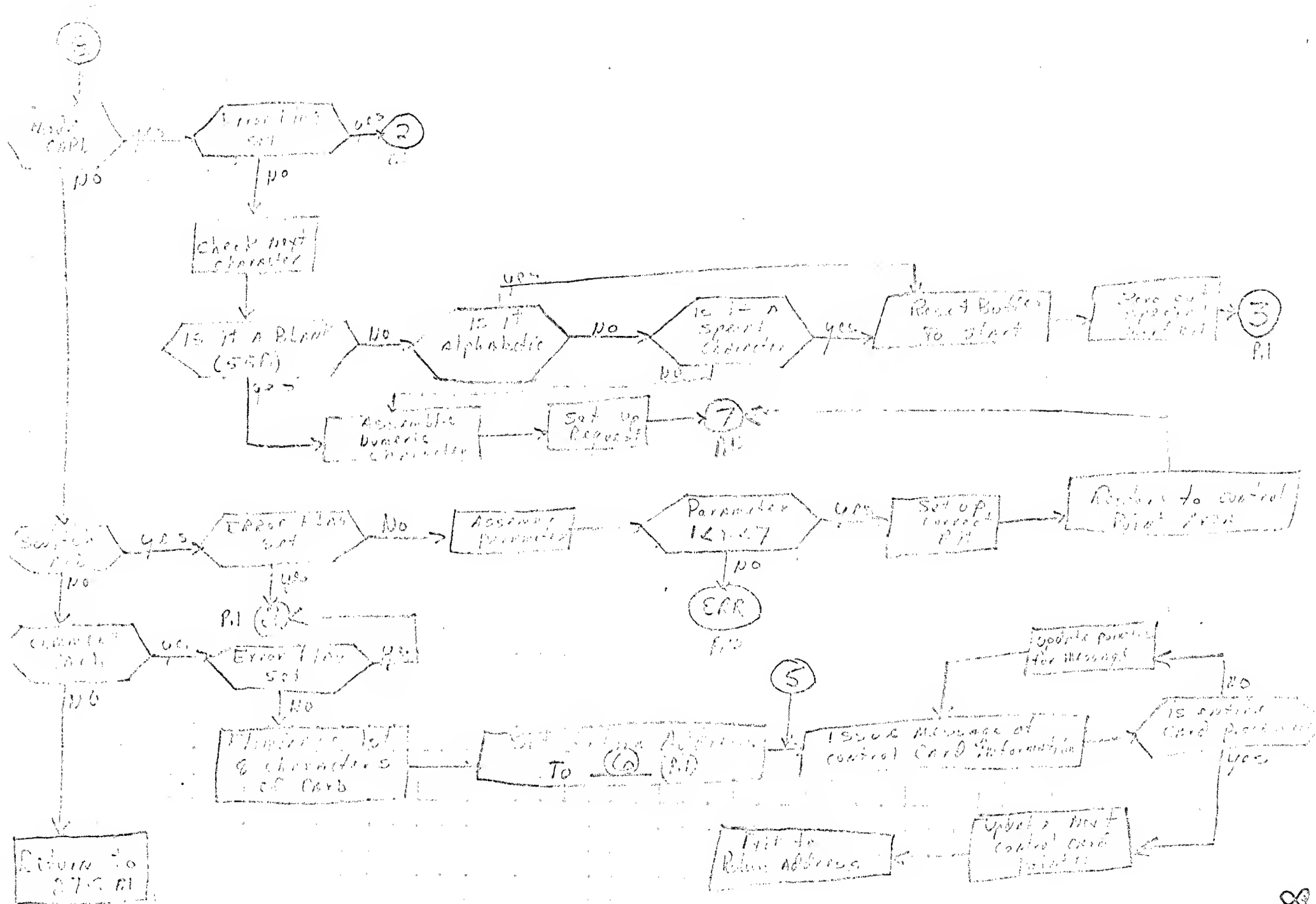


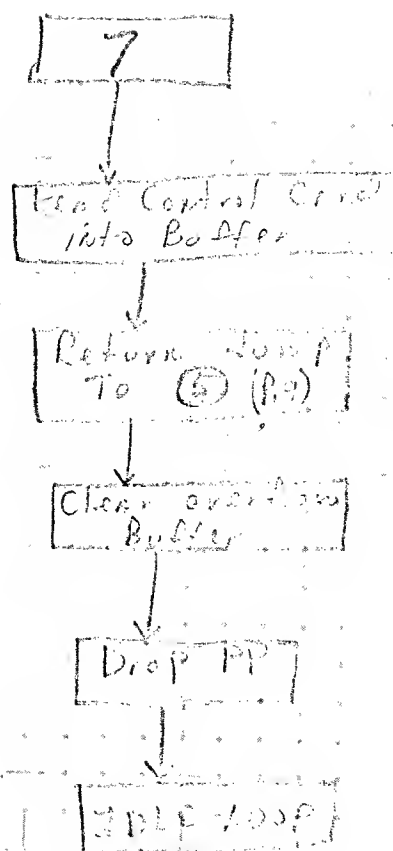
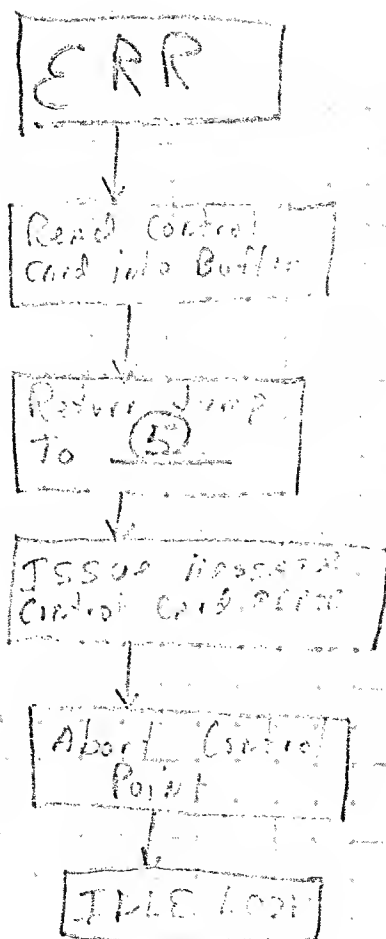












LOT MAIN PROGRAM

LOT TRANSIENT
LOADED INTO
PPR

READ
IR

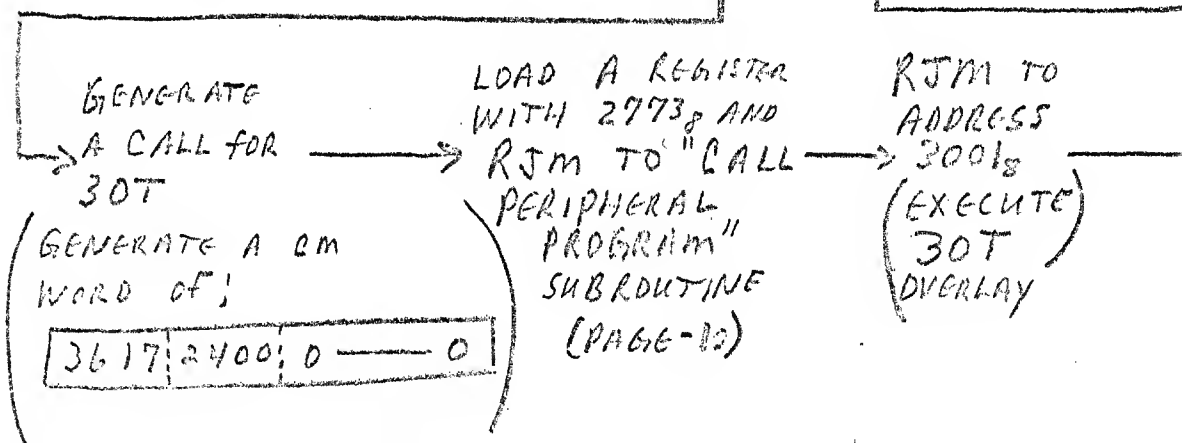
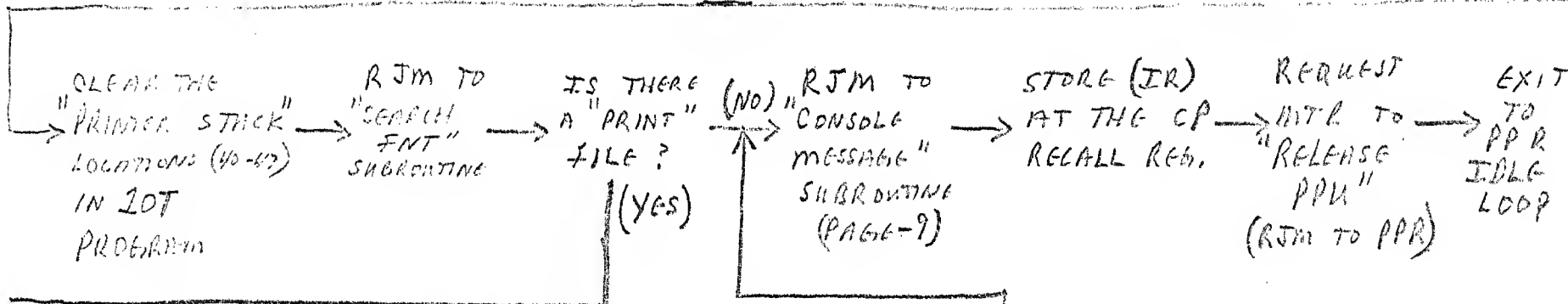
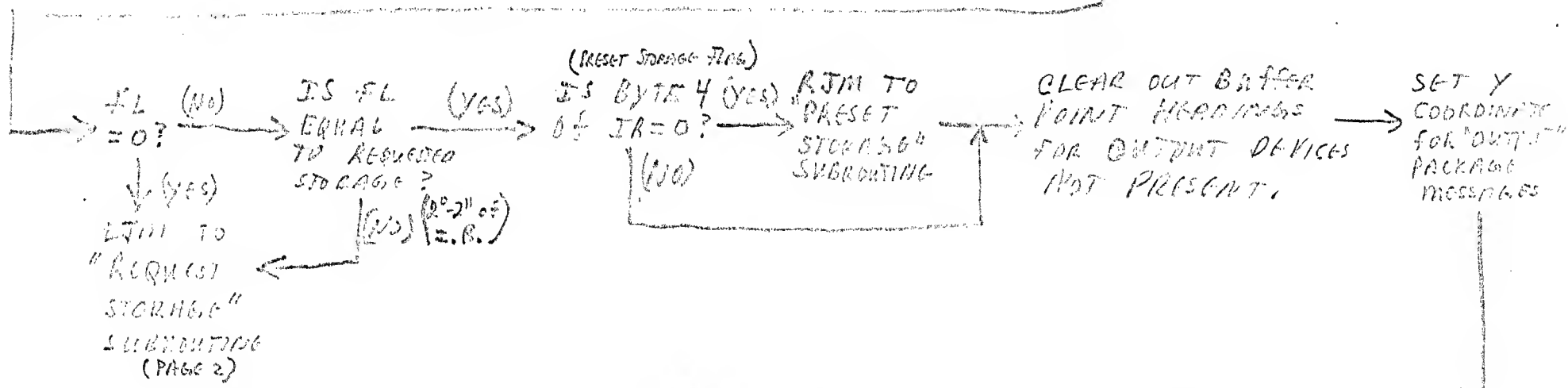
OBTAIN
CP ADDRESS
FROM CP

READ R1,
FL, + LEAD
+ FLAG BYTES

ERROR (NO)
FLAG
+ 1?

REQUEST
MTR TO
RELEASE
PPR

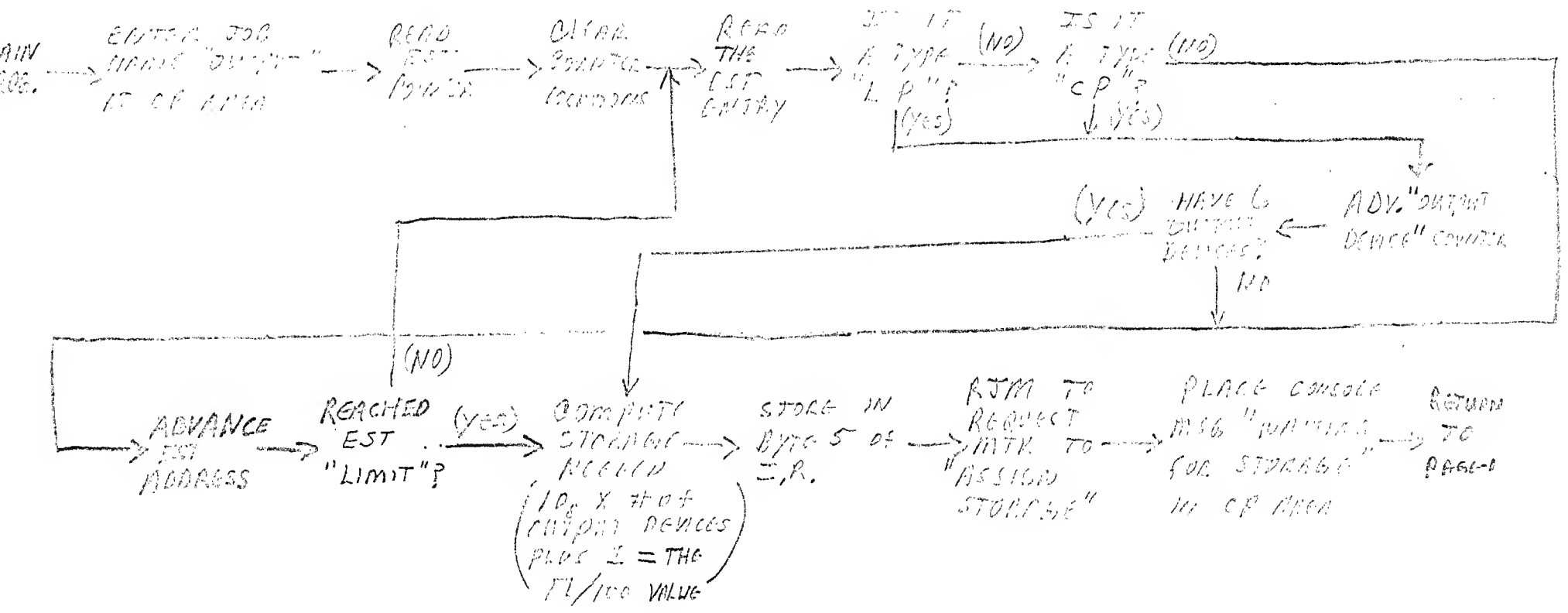
TO
IDLE
LOOP IN
PPR



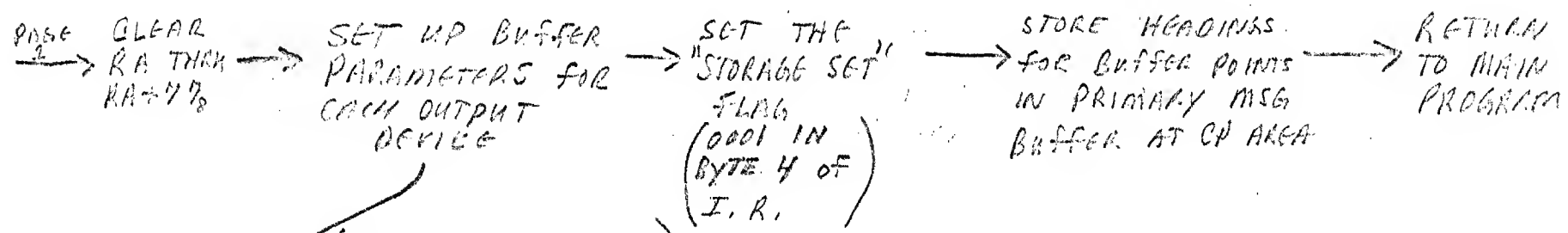
LOT TRANSIENT
MAIN PROGRAM
SCOPE 3.0
3/5/68

REQUEST STORAGE SUBROUTINE

FOR LOT TRANSFER PROGRAM SCOP 3



87-8



(IF ONLY 3 OUTPUT DEVICES, THEN SET UP ONLY 3, ETC.)

$RA+11 = 0 \text{ --- } 0100 \Rightarrow \text{FIRST}$
 $RA+12 = 0 \text{ --- } 0100 \Rightarrow \text{IN}$
 $RA+13 = 0 \text{ --- } 0100 \Rightarrow \text{OUT}$
 $RA+14 = 0 \text{ --- } 01100 \Rightarrow \text{LIMIT}$

} OUTPUT DEVICE AT BUFFER POINT #1

$RA+21 = 0 \text{ --- } 01100 \Rightarrow \text{FIRST}$
 $RA+22 = 0 \text{ --- } 01100 \Rightarrow \text{IN}$
 $RA+23 = 0 \text{ --- } 01100 \Rightarrow \text{OUT}$
 $RA+24 = 0 \text{ --- } 02100 \Rightarrow \text{LIMIT}$

} OUTPUT DEVICE AT BUFFER POINT #2

$RA+31 = 0 \text{ --- } 01100 \Rightarrow \text{FIRST}$
 $RA+32 = 0 \text{ --- } 01100 \Rightarrow \text{IN}$
 $RA+33 = 0 \text{ --- } 01100 \Rightarrow \text{OUT}$
 $RA+34 = 0 \text{ --- } 02100 \Rightarrow \text{LIMIT}$

} OUTPUT DEVICE AT BUFFER POINT #3

$RA+41 = 0 \text{ --- } 01100 \Rightarrow \text{FIRST}$
 $RA+42 = 0 \text{ --- } 01100 \Rightarrow \text{IN}$
 $RA+43 = 0 \text{ --- } 01100 \Rightarrow \text{OUT}$
 $RA+44 = 0 \text{ --- } 02100 \Rightarrow \text{LIMIT}$

} OUTPUT DEVICE AT BUFFER POINT #4

$RA+51 = 0 \text{ --- } 01100 \Rightarrow \text{FIRST}$
 $RA+52 = 0 \text{ --- } 01100 \Rightarrow \text{IN}$
 $RA+53 = 0 \text{ --- } 01100 \Rightarrow \text{OUT}$
 $RA+54 = 0 \text{ --- } 02100 \Rightarrow \text{LIMIT}$

} OUTPUT DEVICE AT BUFFER POINT #5

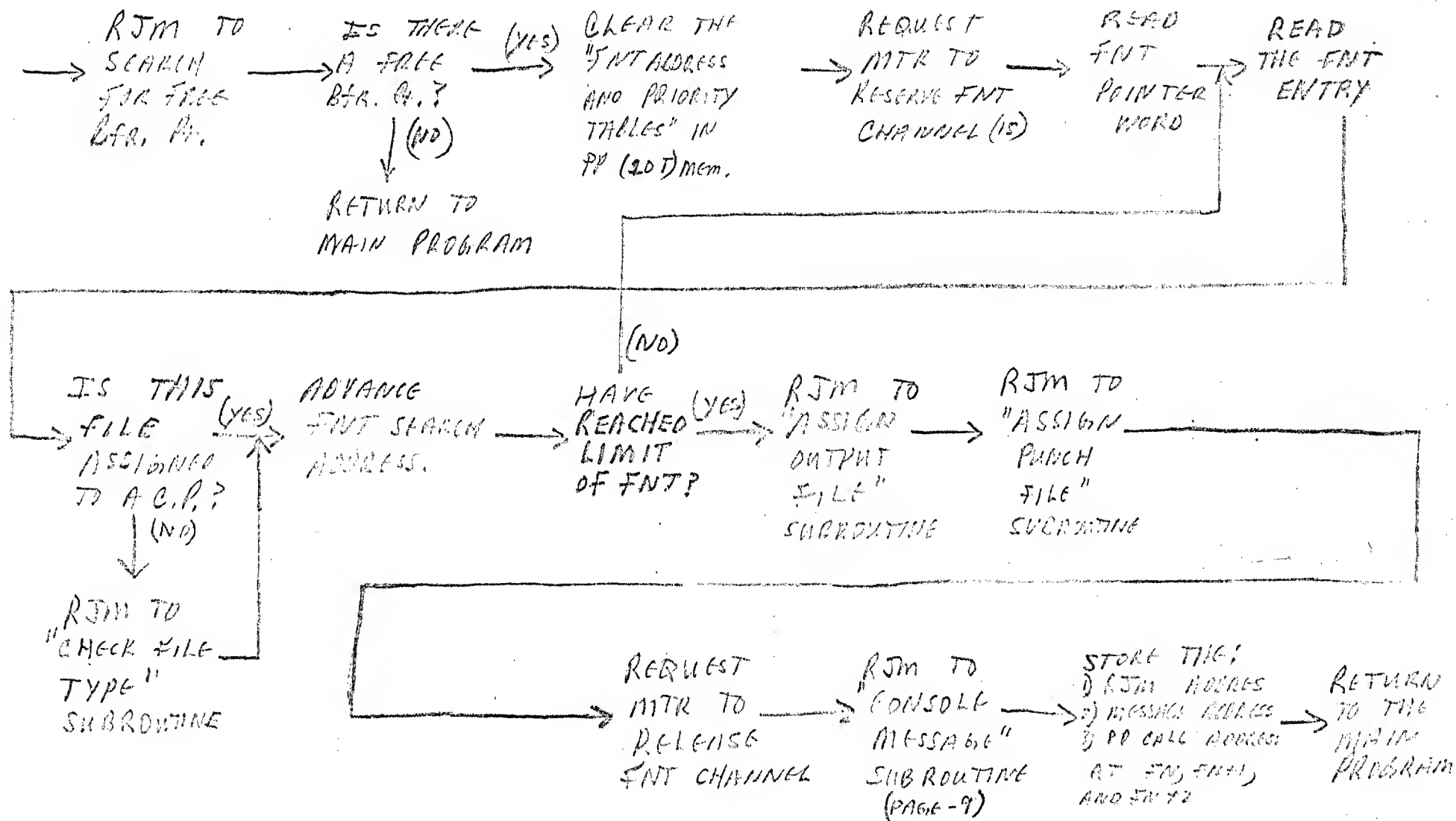
$RA+61 = 0 \text{ --- } 05100 \Rightarrow \text{FIRST}$
 $RA+62 = 0 \text{ --- } 05100 \Rightarrow \text{IN}$
 $RA+63 = 0 \text{ --- } 05100 \Rightarrow \text{OUT}$
 $RA+64 = 0 \text{ --- } 06100 \Rightarrow \text{LIMIT}$

} OUTPUT DEVICE AT BUFFER POINT #6

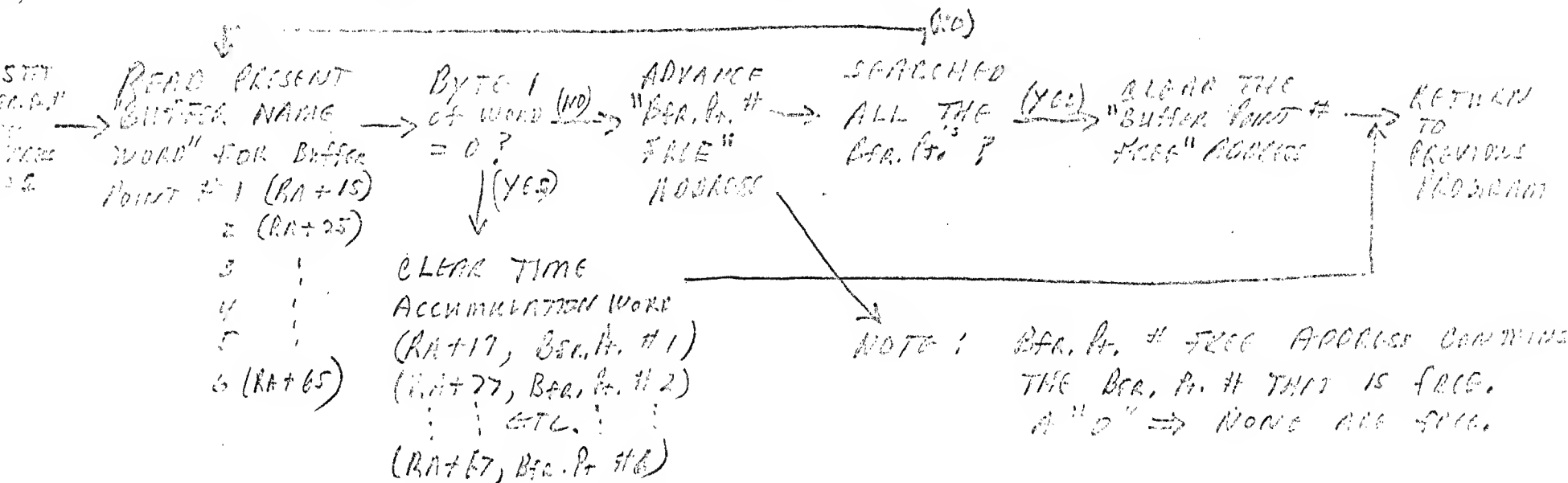
677-8

SEARCH FNT SUBROUTINE

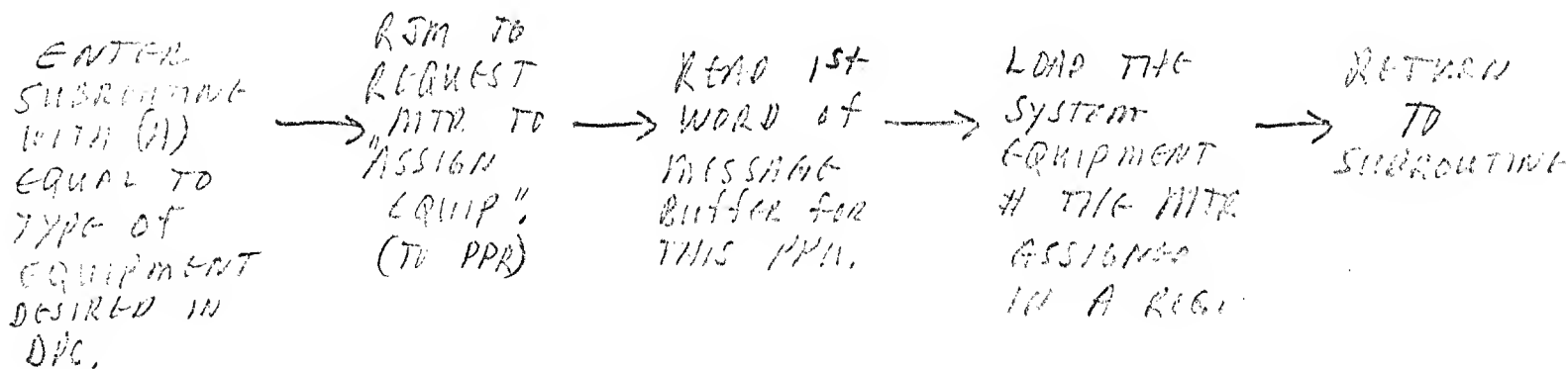
LOT TRANSIENT PROGRAM, SCOPE 3



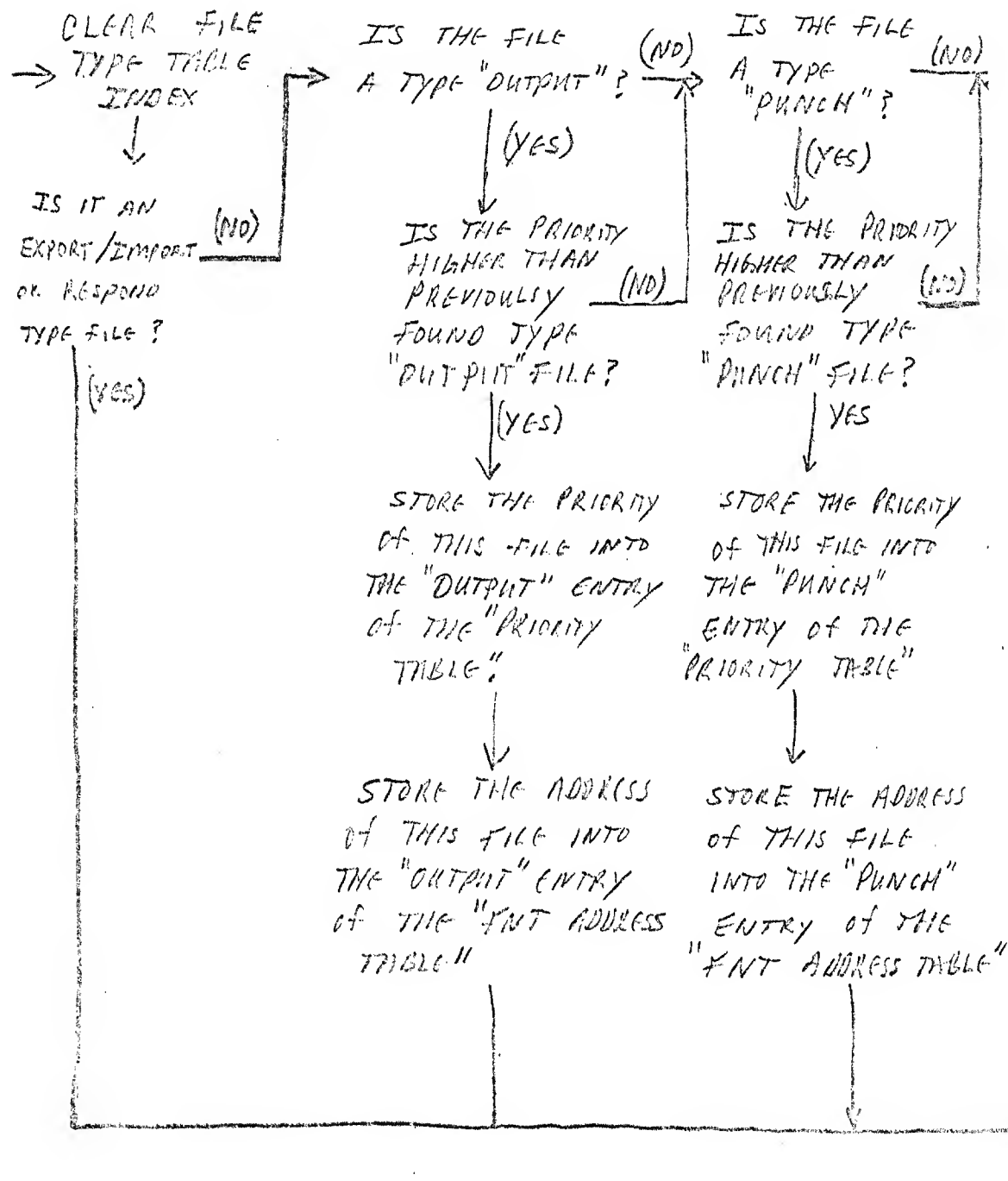
SEARCH FOR FREE BUFFER SUBROUTINE 10T TRANSIENT PROGRAM, SCOPE 3



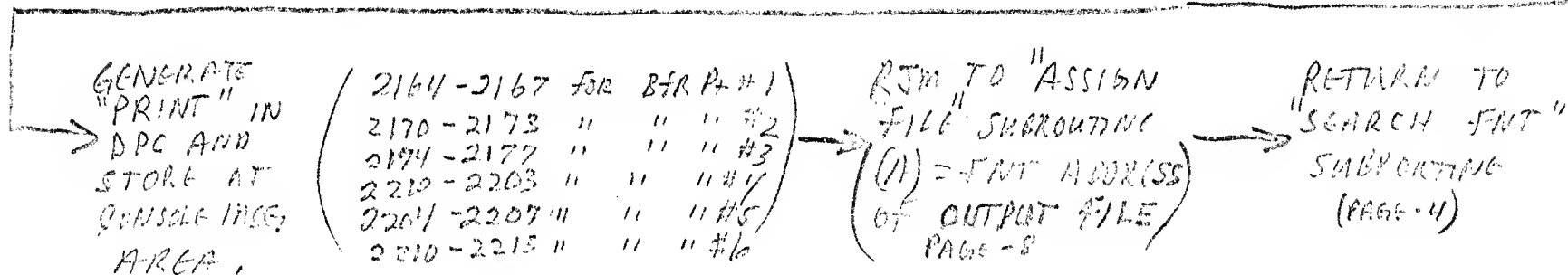
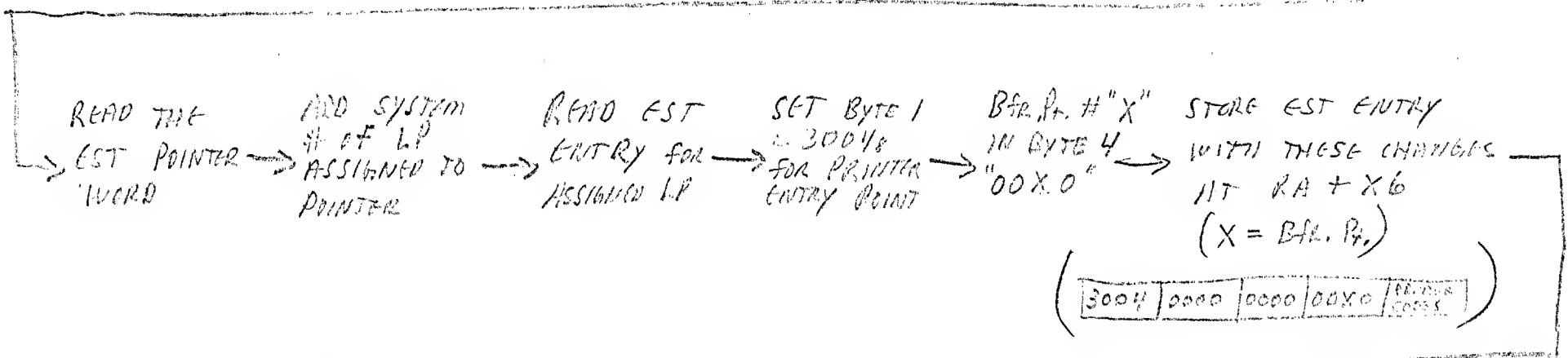
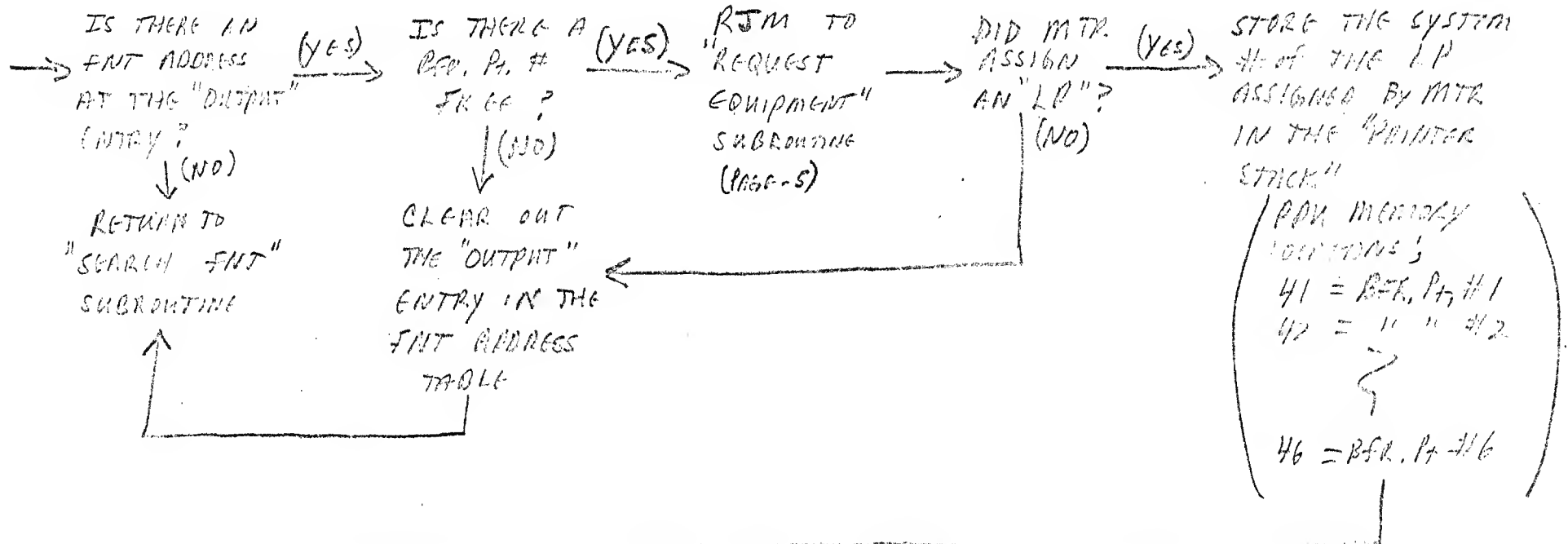
REQUEST EQUIPMENT SUBROUTINE 10T TRANSIENT PROGRAM, SCOPE 2.0, PSR # 47 4/20/67



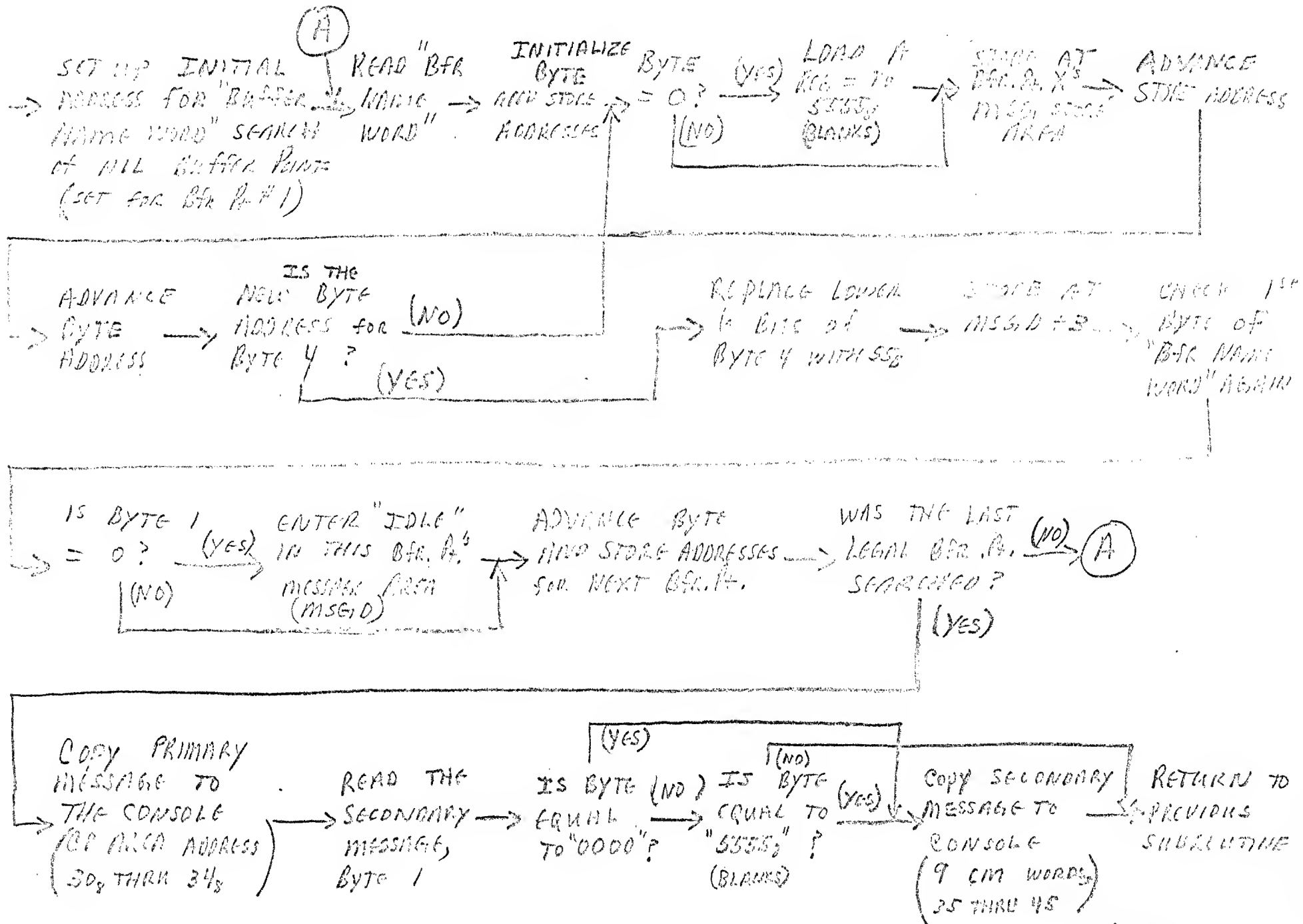
CHECK FILE TYPE SUBROUTINE 10T TRANSIENT PROGRAM



ASSIGN OUTPUT FILE SUBROUTINE LDT TRANSIENT PROGRAM, SCOPE 3 (ASSIGN "PRINT" FILE)

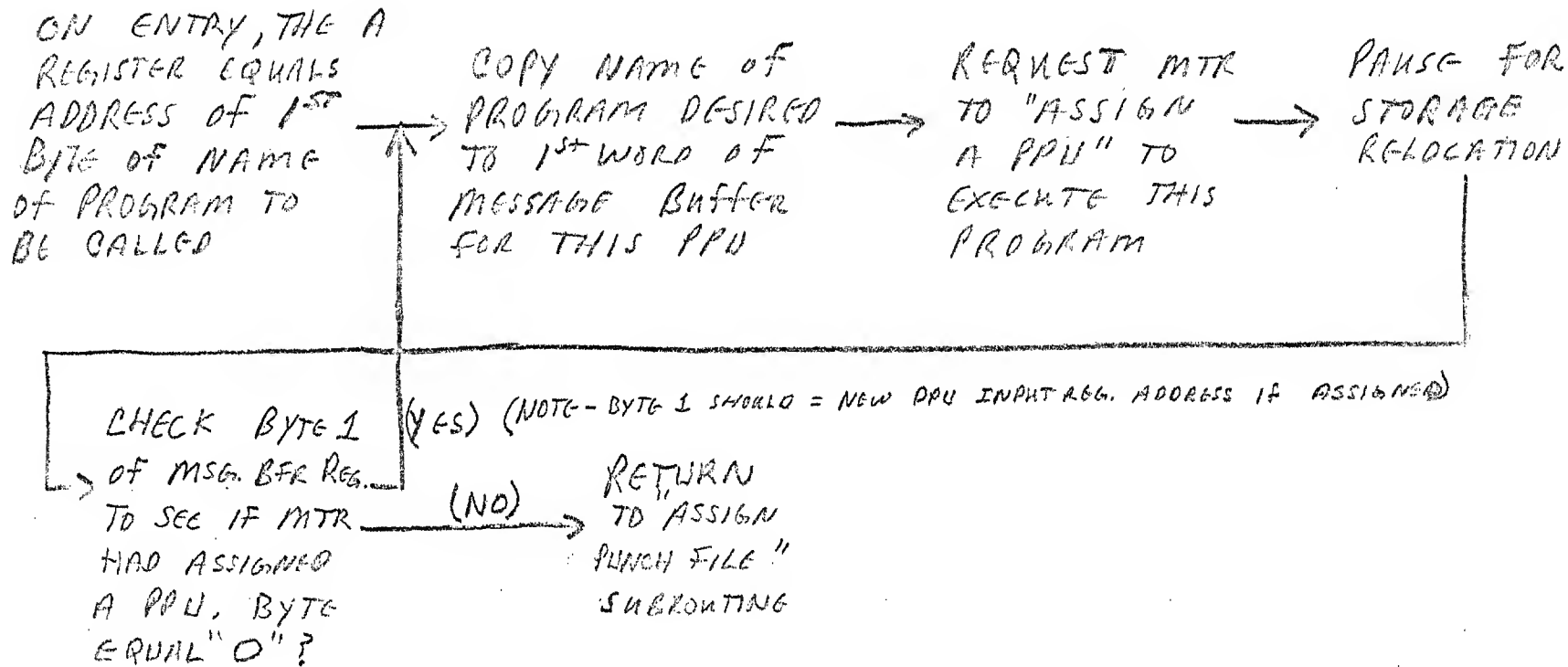


CONSOLE MESSAGE SUBROUTINE | LOT TRANSFER PROGRAM, SCOPE 3



9-9

CALL PERIPHERAL PROGRAM 10T TRANSIENT PROGRAM, SCOPE 3



30T

30T

•PRS•

insert CP # in CIO idls

3200 cycle count

Buffer empty? \xrightarrow{NO} LPD4

\xrightarrow{YES}

FET status contain

End-of-Information?

\xrightarrow{YES}

•PTIM•
TIME ACCOUNTING

LIM
TER

1st time for this file? \xrightarrow{NO} set OUT = FIRST + IN = FIRST + 5

\xrightarrow{YES}

call CIO

To READ DISK FILE

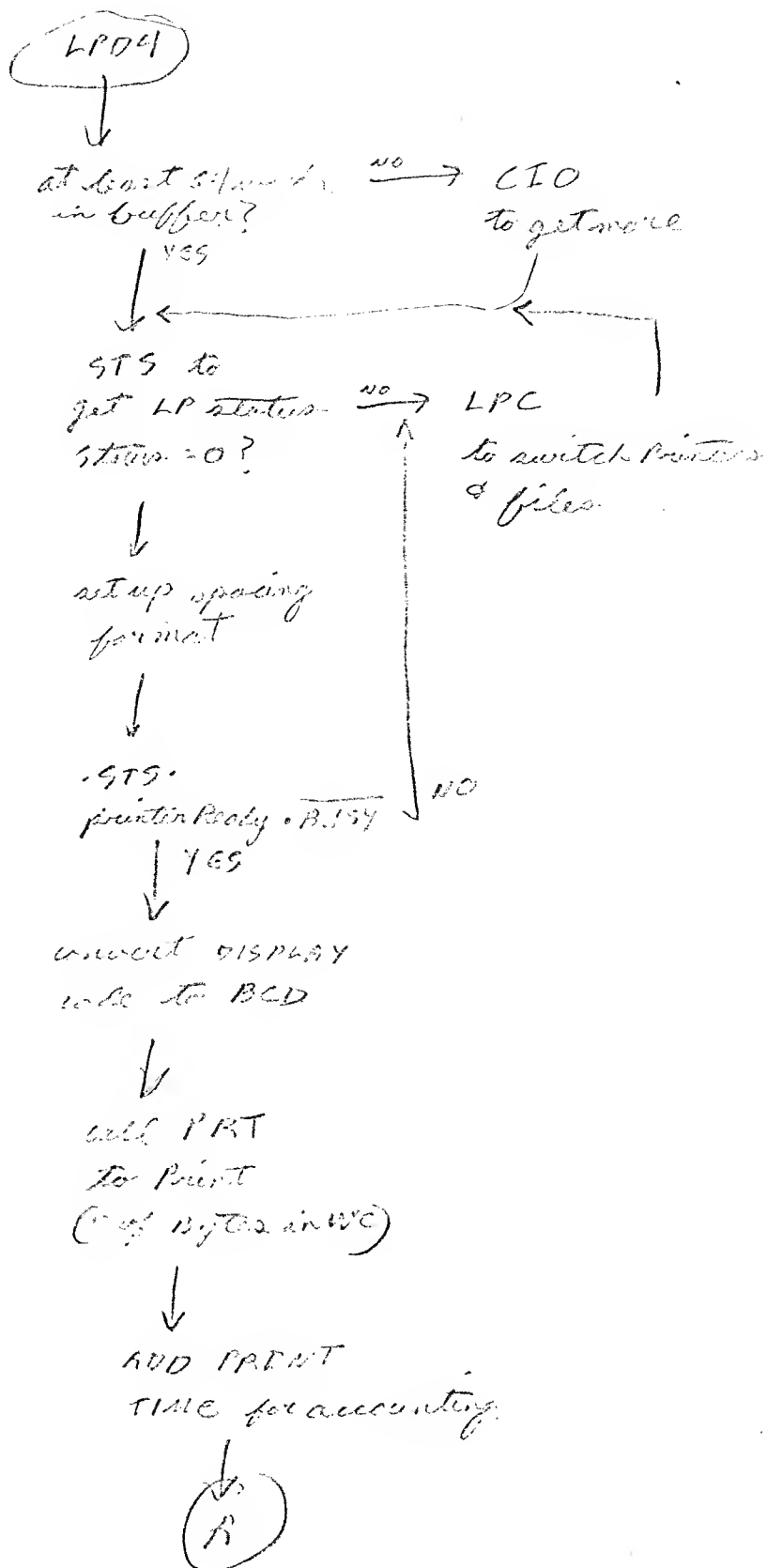
LPC

Try to change to another

printer + file

Another LP & FILE READY • BUSY

1



TER

in REPEAT flag set at LP+0?
 NO → get EST ordinal for LP → call CLP in LOT to DAOP PRINTER

YES
 clr flags
 wait for File Busy
 clr EOI flag in FET

↓
 LPC
 check for active LP

↓
 CALL CIO To Rewind file

A

LPC

store return addr. in FET+6

↓
 clr PR
 increase CC by 1

↓
 CC=100
 YES → call FNT in LOT to find new print file
 NO → ADD 1 to PS (which FET we are concerned with) if its too high, set PS=1.

set return addr. from FET+6

are bits 0-11 of RA+0=0?
 YES → EXIT to LPD cycle for new FET
 NO → PROCESS FLAG

Read new FET → PP

NO
 is (PS+FET#) = zero?
 → increase PR# by 1 to find a new printer
 → is PR# total number of FET'S?
 YES → all FET'S have been checked for an active LP. EXIT 30T BACK TO 10T
 NO → ADD 1 to PS (which FET we are concerned with) if its too high, set PS=1.

ADD 1 to PS (which FET we are concerned with) if its too high, set PS=1.

all FET'S have been checked for an active LP. EXIT 30T BACK TO 10T

6000 SCOPE LOADER

(SEE ALSO APPENDIX D OF SCOPE 3.1 REFERENCE MANUAL)

9-1

RA

ERROR EXIT .

RA+1

COMMUNICATIONS

RA+2

PARAMETERS FROM THE "PROGRAM CALL" CARD UPON
AVAILABLE TO USER DURING EXECUTION FOR
FILE NAMES AND OTHER USAGE

RA+63

RA+64

PROGRAM CALL OR FILE NAME

NUMBER OF
PARAMETERS

RA+65

FWA SEGMENT
TABLE

CORE NEXT

RA+66

FWA LOADER
TABLES

DEBUG
CONTROL
BITS(31-28)

U

V

FWA PROGRAM

RA+67

W

X

Y

FWA LOADER

RA+70

CARD IMAGE OF CARD WHICH CALLED
FOR EXECUTION

RA+77

U (4 BITS) LDR DIRECTIVES

V (6 BITS) LEVEL OF INCOMING OVERLAY

W SEGMENT LOAD - ADDRESS OF LOWEST LINKAGE TABLE

OVERLAY LOAD - ADDRESS OF BLANK COMMON

X (5 BITS)

E	D	C	B	A
---	---	---	---	---

Y (3 BITS) LAST CONTROL CARD

A RSS BIT

B NO MAP BIT

C REQUEST EXIT FLAG

D END OF LOAD

E PARTIAL MAP BIT

0 PROGRAM CALL

1 LOAD

2 EXECUTE

4 NOGO

RA+100

* Section Name	0 0 0 0 0 0
Program Name	0 0 0 0 0 0
Program Name	0 0 0 0 0 0
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0 0 0 0 0 0
* Segzero Name	0 0 0 0 0 0
Program/Section Name	0 0 0 0 0 0
Program/Section Name	0 0 0 0 0 0
* Segment Name	0 0 0 0 0 0
Program/Section Name	0 0 0 0 0 0
Program/Section Name	0 0 0 0 0 0
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	
* $2^{59} = 1$	

SECTION table

SEGMENT table

FWA of
Segzero

RA+100

Labeled Common declared
in Program A

Program A

Labeled Common if Program B
declares Labeled Common not
in Program A

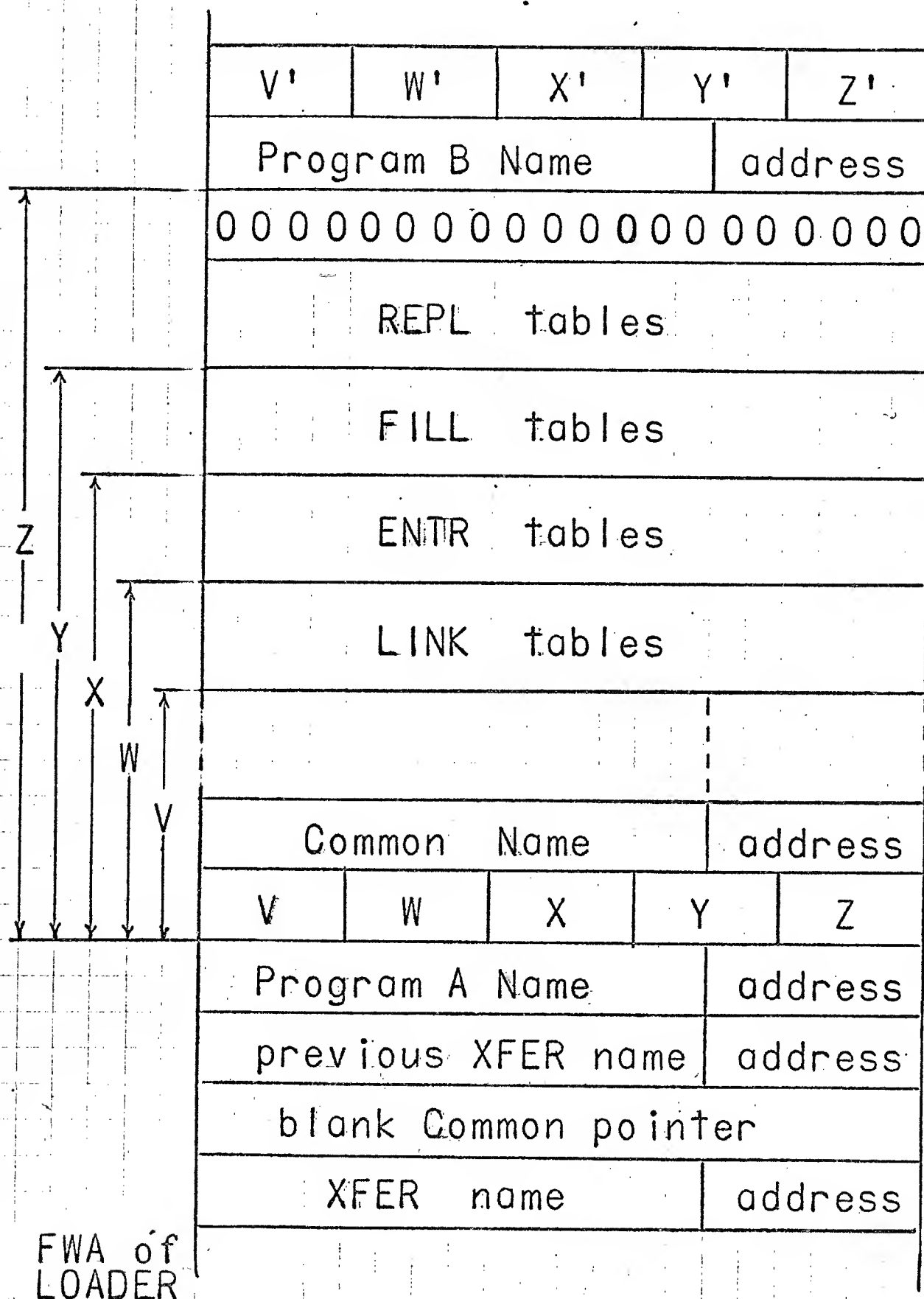
Program B

BEGINNING OF
BLANK COMMON

Loader Tables

LOADER


RA+FL-1



next table															
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0															
nal		abs. loc.			P	R		rel. loc.							
P	R		rel. loc.			2nd		exter-							
P	R		rel. loc.			P	R		rel. loc.						
1st external									abs. loc.						
pointer to next LINK table															

next table															
0000000000000000															
												reloc	rel. loc.		
3rd entry												abs. loc.			
												reloc	rel. loc.		
2nd entry												abs. loc.			
												reloc	rel. loc.		
1st entry												abs. loc.			
pointer to next) ↑															

pointer to next
ENTR table



next table											
00000000000000000000											
P	R	rel. loc			P	R	rel. loc				
P	R	rel. loc			control byte						
control byte					P	R	rel. loc				
pointer to next FILL table											

next table															
0000000000000000															
C				B				D R				D			
				I				S R				S			
C				B				D R				D			
				I				S R				S			
pointer to next) ↑															

pointer to next
REPL table

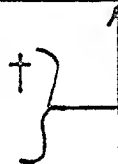


TABLE FORMAT

CN	00	WC	000	LR	L

CN Code Number - identifies type of table

WC Word Count - number of words in table
(total table length = WC + 11)

LR and L used only with TEXT table

PIDL TABLE

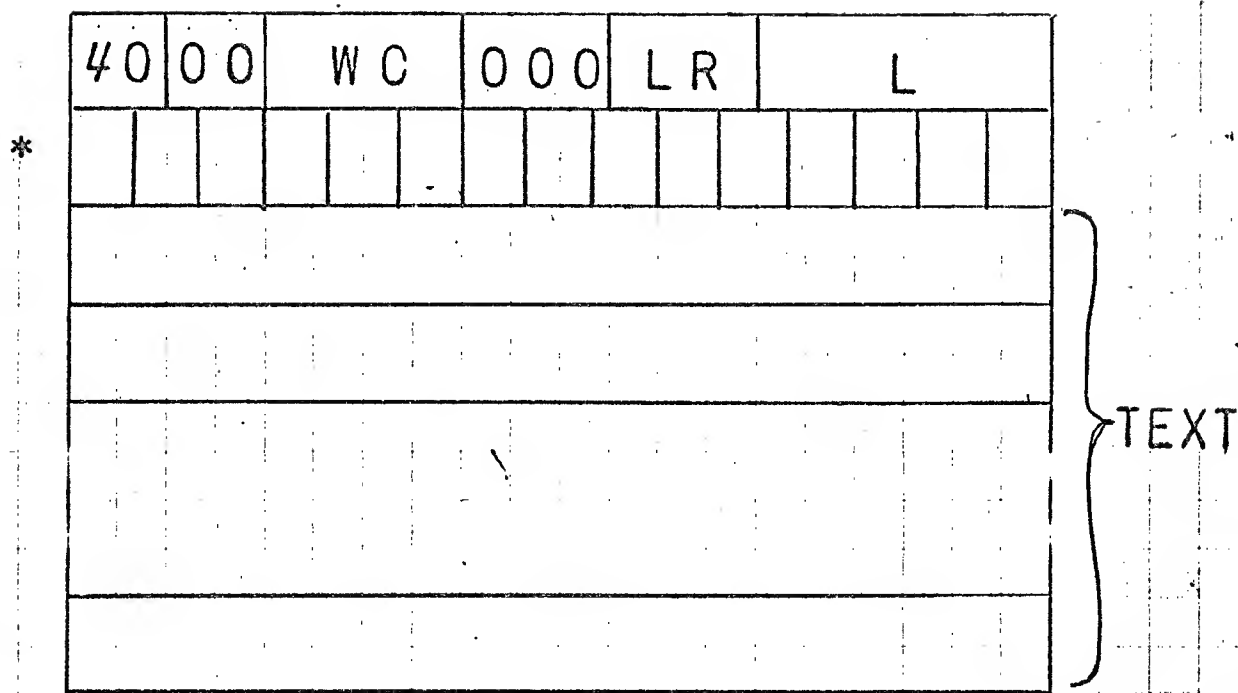
34	00	WC	000	0000	000000
PROGRAM NAME					PROGRAM LENGTH
COMMON NAME					BLOCK LENGTH
COMMON NAME					BLOCK LENGTH
COMMON NAME					BLOCK LENGTH

First entry must be the program name

Blank Common same as any common (name is 55555555555555)

All names left justified with zero fill

TEXT TABLE



* Relocation bytes (4 bits)

WC limited to maximum of 16
more text requires more
TEXT tables

LR specifies area into which text is
to be loaded
0 = absolute (relative to RA)
1 = program
3 - n = common

L first word address for text
within LR area

RELOCATION BYTES

1000	upper+
1100	upper-
1010	upper+ lower+
1011	upper+ lower-
1110	upper- lower+
1111	upper- lower-
0010	lower+
0011	lower-
0100	middle+
0110	middle-
0000	no relocation

FILL TABLE

4200	WC	000	0000	00000000
00000000	reloc	P	R	L
P	R	L	P	R
P	R	L	00000000	reloc
		P	R	L

reloc - relocation

R - area

0 absolute

0 absolute

1 program

1 program

2 negative prog.

3-n common

3-n common

L - first word address in area

P - position in address

100 lower

101 middle

110 upper

	IDENT	PROG
	ORG	2
P1	VFD	30/5LINPUT,12/0,18/BE1
P2	VFD	36/6LOUTPUT,6/0,18/BE2
P3	VFD	42/7LSCRATCH,18/BE3
	USE	/ALPHA/
AL1	BSS	2000B
AL2	BSS	2000B
	USE	/BETA/
BE1	BSS	5
BE2	BSS	5
BE3	BSS	5
	USE	//
BUFFER	BSS	2000B
	USE	
	SA1	P1-2
	SB1	A1
	MX3	42
LOOP	SA1	A1+B1
	ZR	X1,OUT
	BX6	X1*X3
	SA6	X1

ZR LOOP
USE /BETA/
ORG 0
VFD 30/5LINPUT,30/0
VFD 60/ALI
VFD 60/AL2
VFD 36/6LOUTPUT, 24/0
VFD 60/AL2
VFD 60/AL2+2000B
VFD 42/7LSCRATCH,18/0
VFD 60/BUFFER
VFD 60/BUFFER+2000B
USE *

OUT

3	4	0	0	0	0	0	4	0	0	0	0	0	0	0	0	0	0	0	0
2	0	2	2	1	7	0	7	0	0	0	0	0	0	length					
0	1	1	4	2	0	1	0	0	1	0	0	0	0	0	0	4	0	0	0
0	2	0	5	2	4	0	1	0	0	0	0	0	0	0	0	0	0	1	7
5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	0	0	2	0	0
<hr/>																			
4	0	0	0	0	0	0	7	0	0	0	0	0	0	1	0	0	0	0	2
*	0	0	0	0	4	1	0	4	0	0	0	0	0	0	0	0	0	0	0
1	1	1	6	2	0	2	5	2	4	0	0	0	0	0	0	0	0	0	0
1	7	2	5	2	4	2	0	2	5	2	4	0	0	0	0	0	0	0	5
2	3	0	3	2	2	0	1	2	3	0	3	1	0	0	0	0	0	1	2
5	1	1	0	0	0	0	0	0	0	6	4	1	1	0	4	3	3	5	2
5	4	1	1	1	0	3	0	1	0	0	0	0	1	0	1	1	6	1	3
5	3	6	1	0	0	4	0	0	0	0	0	0	0	5	4	6	0	0	0
<hr/>																			
4	0	0	0	0	0	1	2	0	0	0	0	0	4	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	1	1	6	2	0	2	5	2	4	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	2	0	0
1	7	2	5	2	4	3	0	2	5	2	4	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	2	0	0

* 0 | 0 | 0 | 0 | 4 | 1 | 0 | 4 | 0 | 0
 00000000000000001000001000010000000000
 no | no | no | up | mid | mid | no

0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	4	0	0	0
2	3	0	3	2	2	0	1	2	4	0	3	1	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	2	0	0
<hr/>																			
4	2	0	0	0	0	0	6	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	3	4	0	0	4	0	0	0	0	0	1
4	0	0	4	0	0	0	0	0	2	4	0	0	4	0	0	0	0	0	4
4	0	0	4	0	0	0	0	0	5	0	0	0	0	0	0	0	0	0	4
4	0	0	1	0	0	0	0	0	2	4	0	0	1	0	0	0	0	0	3
4	0	0	1	0	0	0	0	0	4	0	0	0	0	0	0	0	0	0	5
4	0	0	4	0	0	0	0	0	7	4	0	0	4	0	0	0	0	1	0
<hr/>																			
4	0	0	0	x	x	x	x	0	0	0	0	0	1	0	0	0	0	1	0

LINK TABLE

44	00	WC	000	000	000000
name of external					000000
P	R	L	P	R	L
P	R	L	name of ext-		
ernal		000000	P	R	L
P	R	L	P	R	L

P position in word
 4 lower
 5 middle
 6 upper

R area
 0 absolute
 1 program
 3-n common

L location reference made from

ENTR TABLE

9-20

36	00	WC	00	0000	0000000
name of entry point					0000000
000000000000				R	L
name of entry point					0000000
000000000000				R	L
name of entry point					0000000
000000000000				R	L

R area
 0 absolute
 1 program
 3-n common

L location in area

	IDENT	GREEK
	ENTRY	ALPHA,BETA,GAMMA
	EXT	ALEPH,BETH,GIMMEL
ALPHA	DATA	0
	RJ	ALEPH
DELTA	SX6	B7
	SA6	A0
	RJ	GIMMEL
	SA1	A0
	JP	ALPHA
BETA	DATA	0
SA1	BETA	
	BX6	X1
	SA6	ALPHA
	RJ	BETH
	JP	DELTA
GAMMA	DATA	0
	SA1	GAMMA
	BX6	X1
	SA6	GIMMEL
	JP	GIMMEL+1

[illegible]

0	0	0	0	2	0	1	0	4	2	0	0	4	0
0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	0000	00
no	no	no	mid	no	up	up	up	no	up	no			

REPL TABLE

4300	WC	000	0000	0000000
	I		SR	S
C	B		DR	D
	I		SR	S
C	B		DR	D

SR source area

S address in area

B block size

DR destination area

D address in area

C number of replications

I replication interval

SR and DR

0 absolute

1 program

3-n common

4-25

4	6	0	0	0	0	1	M	M	D	D	Y	Y
name										000000		

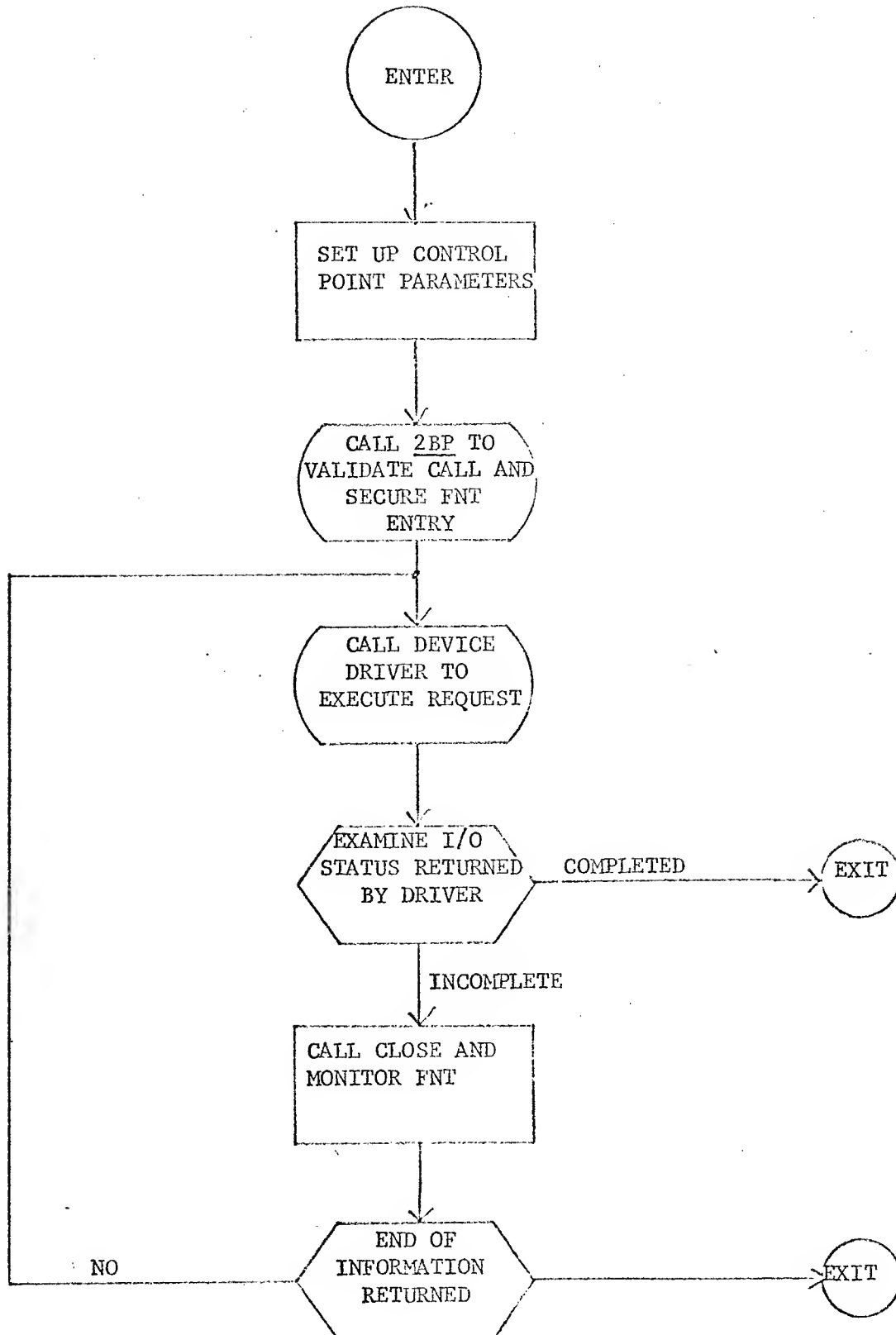
9-26

words 2 through
end of logical
record are absolute
text

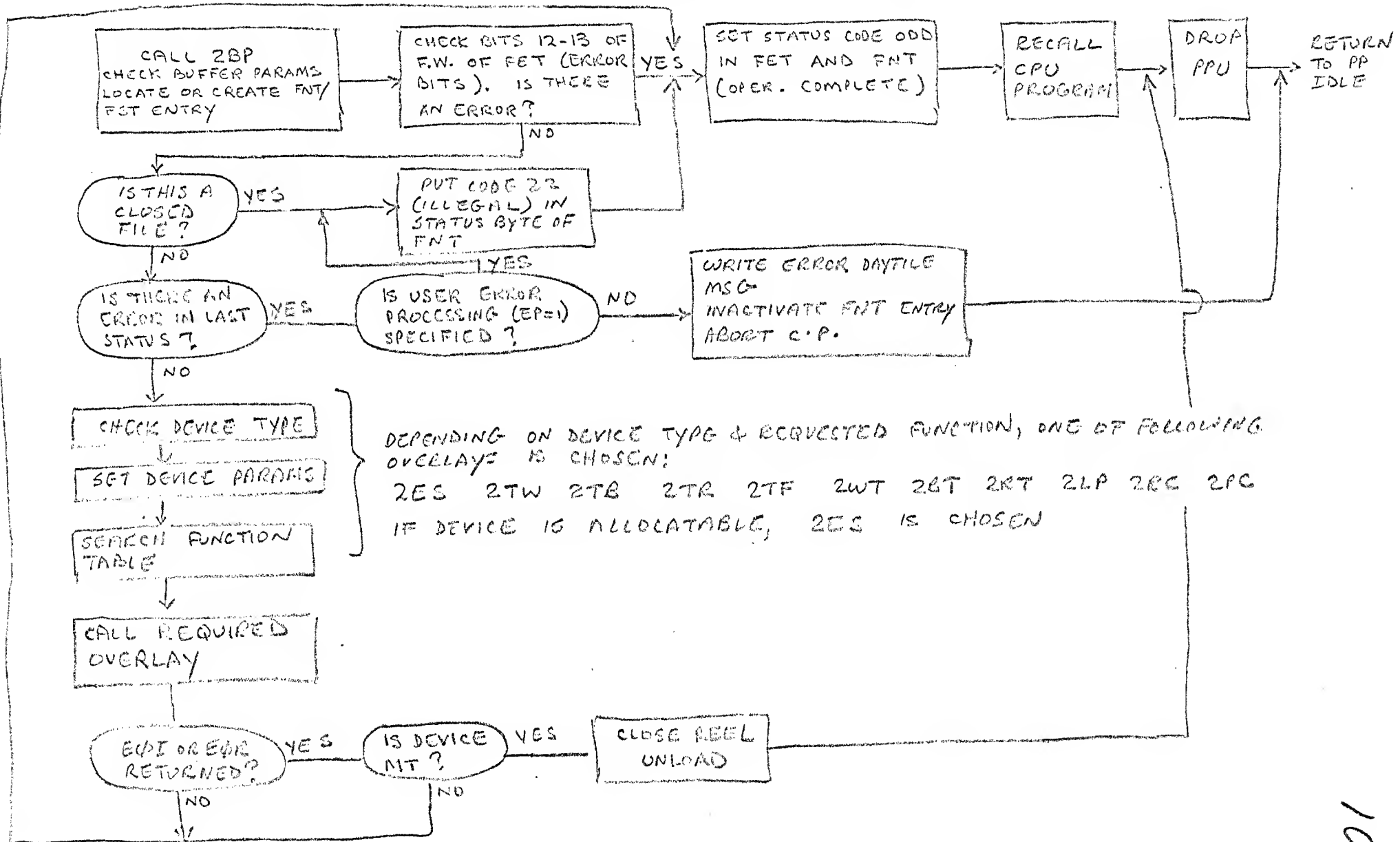
CIO

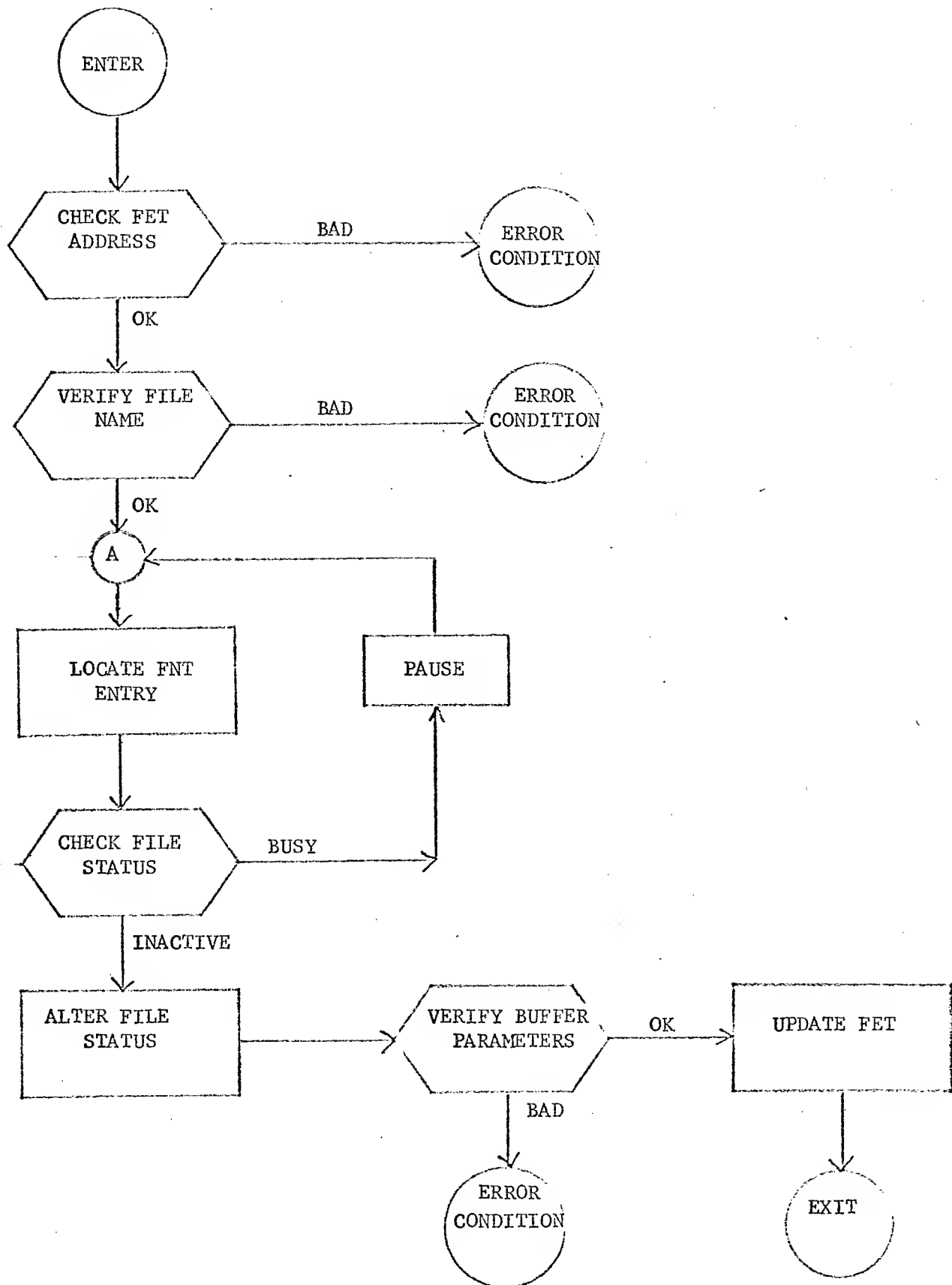
CIO

10-1
□



CICP (CIRCULAR I/O)

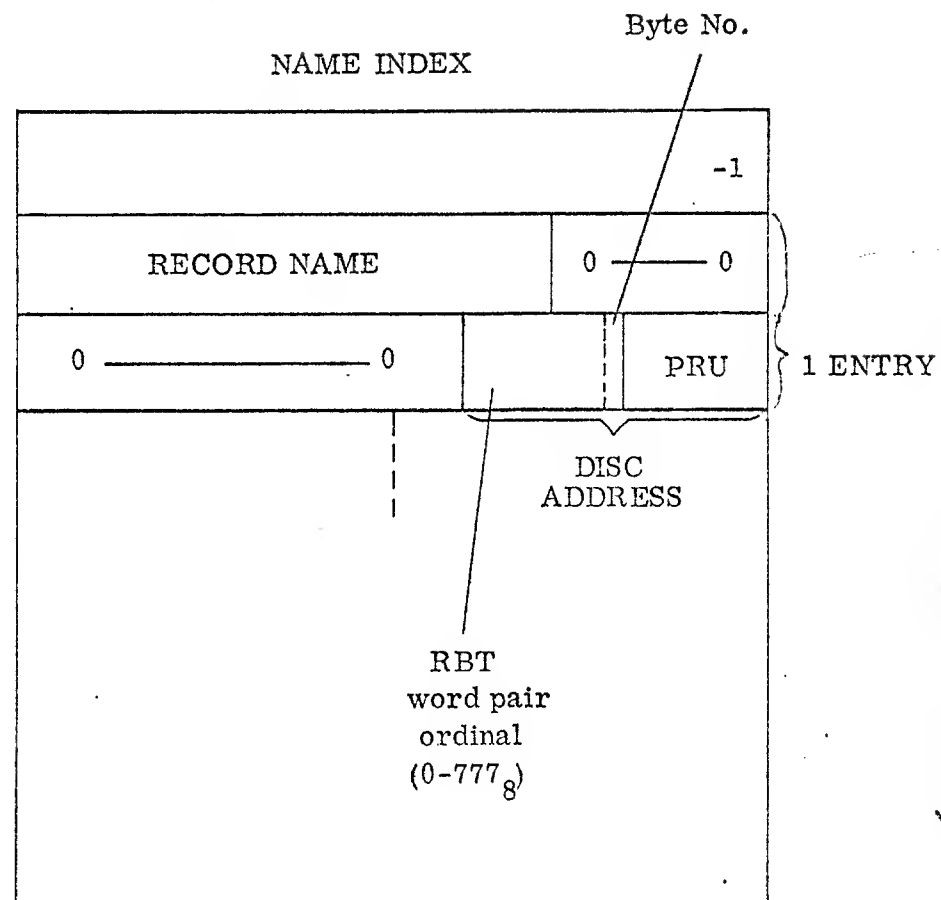
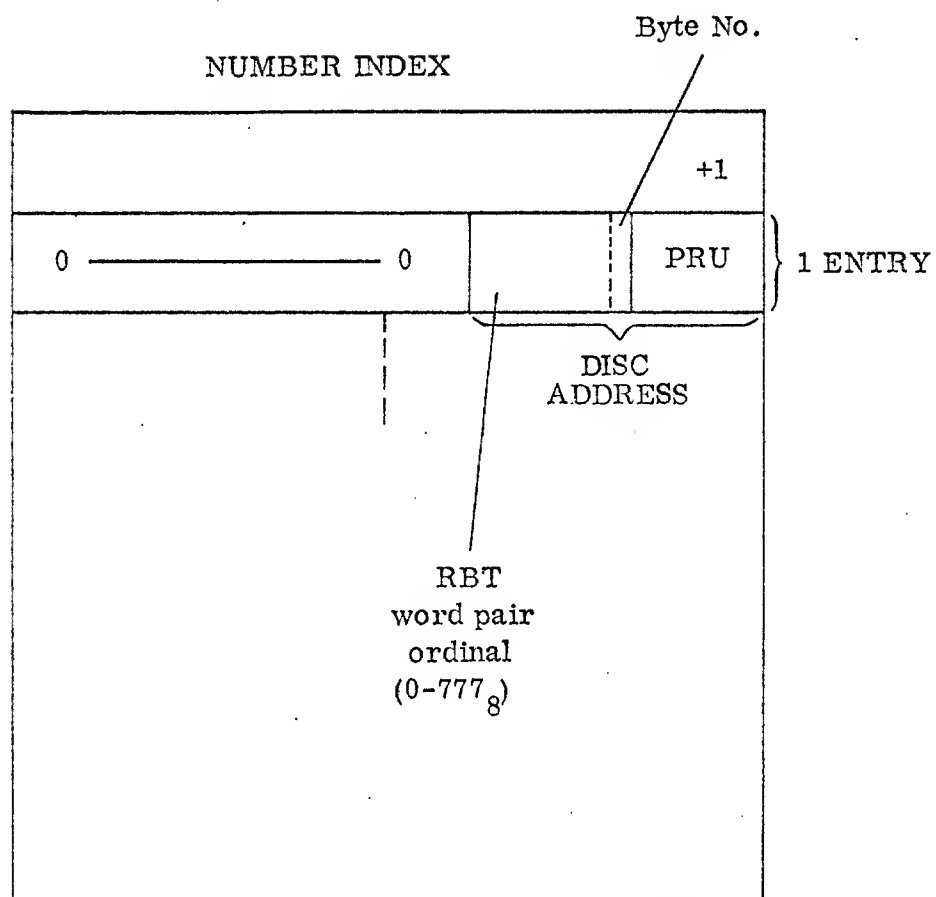




FILE ENVIRONMENT TABLE

Bits	59	47	44	35	32	29	23	17	0	Words
logical file name (lfn)									code and status	1
device type	r	n	u	e	disposition code		L	FIRST	2	
0									IN	3
0									OUT	4
FNT pointer	record block size			physical record unit size			LIMIT		5	
				working storage fwa				working storage lwa+1	6	
						record request/return information			7	
record number					index length		index address		8	
				EOI address				error address	9	
Label file name (first 10 chars)										10
Label file name (last 10 chars)										11
edition number	retention cycle					creation date				12
position number				multi-file name			reel number			13

RANDOM ACCESS FILE INDEXES



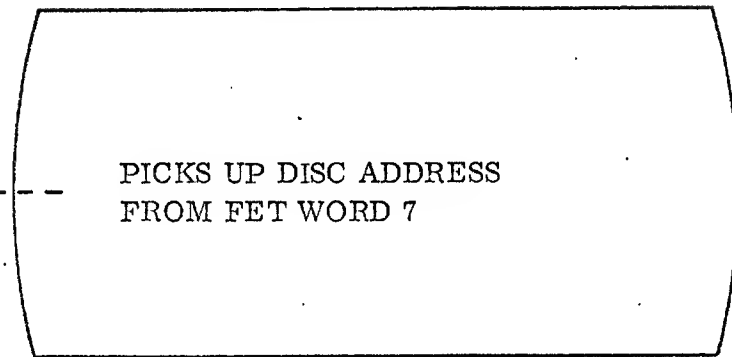
SCOPE I/O (TAPE + DISK)

10-4

READ FILE

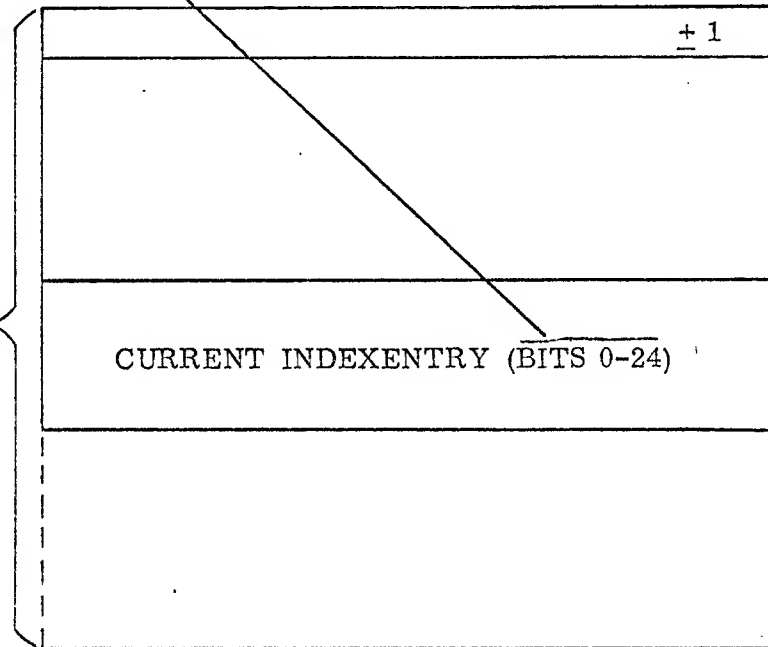
1	LFN		
2			FIRST
3			IN
4			OUT
5			LIMIT
6			
7			DISC ADDRESS
8	REC #		
		INDEX LENGTH	INDEX ADDRESS

PP ROUTINE



INDEX

1st WORD



FET

WRITE FILE

1

LFN

2

FIRST

3

IN

4

OUT

5

LIMIT

6

7

Index Entry Address

8

REC #

INDEX
LENGTH

INDEX
ADDRESS

PP ROUTINE

PLACES DISC ADDRESS INTO
BITS 0-24 OF INDEX WORD
SPECIFIED IN FET WORD 7

INDEX

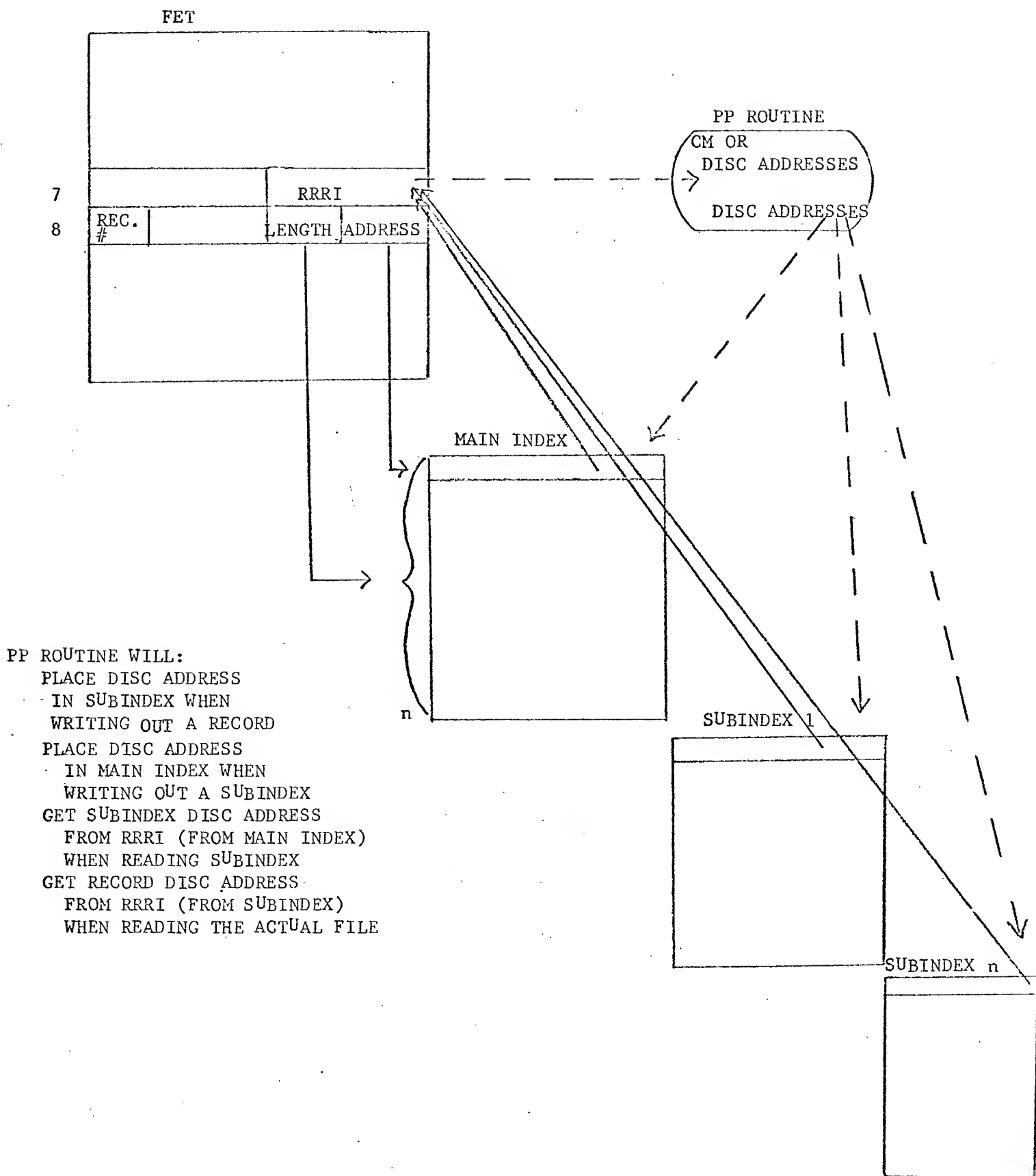
1st WORD

+ 1

CURRENT INDEX ENTRY

9-01/10-6

MULTI INDEX



MAJOR TAPE ROUTINES

10-8

CIO: GENERAL CHECKING AND ROUTING ROUTINE FOR ANY INPUT/
OUTPUT REQUEST.

2TR: READ DRIVER FOR $\frac{1}{2}$ " MAGNETIC TAPE TRANSPORTS.

2TW: WRITE DRIVER FOR $\frac{1}{2}$ " MAGNETIC TAPE TRANSPORTS.

2TB: POSITION BACKWARD ROUTINE FOR $\frac{1}{2}$ " MAGNETIC TAPE.

2TF: POSITION FORWARD ROUTINE FOR $\frac{1}{2}$ " MAGNETIC TAPE.

4LB: READ AND WRITE ROUTINE FOR $\frac{1}{2}$ " MAGNETIC TAPE WHICH PROCESSES
LABEL INFORMATION.

OPEN: GENERAL FILE INITIALIZING ROUTINES. FOR LABELED TAPES, OPEN
CALLS THE APPROPRIATE ROUTINES TO POSITION TAPE AND READ OR
WRITE LABELS.

CLOSE: GENERAL FILE TERMINATING ROUTINE. FOR LABELED TAPES,
CLOSE CALLS THE APPROPRIATE ROUTINES TO POSITION TAPE AND
READ OR WRITE LABELS.

NOTES:

DIFFERENCE IN PHYSICAL AND LOGICAL TAPE MARKS

10-9

	HOW WRITTEN	HOW RECORDED	WHEN WRITTEN	ACTION ON READING
PHYSICAL EOF (TAPE MARK)	DEVICE FUNCTION	1 CHAR. 17 ₈ EVEN PARITY <i>and 6" gap</i>	IN CONJUNCTION WITH TAPE LABEL <i>[difference in label]</i>	EOI & EOF RETURNED TO CIO
LOGICAL EOF	DATA WRITE	8 CHAR. BIN BCD 0000 2020 0000 2020 0017 2017	WRITEF	SETS EOF BITS IN CODE & STATUS

SINGLE REEL FILE

VOL1
HDR1
TAPE MARK
DATA
TAPE MARK
EOF1
TAPE MARK
TAPE MARK

MULTI-REEL FILE

VOL1
HDR1
TAPE MARK
FIRST VOLUME DATA
TAPE MARK
EOV1
TAPE MARK
TAPE MARK

⋮

VOL1
HDR1
TAPE MARK
LAST VOLUME DATA
TAPE MARK
EOF1
TAPE MARK
TAPE MARK

MULTI-FILE REEL

7-8
10-11

VOL1
HDR1
TAPE MARK
FILE A DATA
TAPE MARK
EOF1
TAPE MARK
HDR1
TAPE MARK
FILE B DATA
TAPE MARK
:
:
:
TAPE MARK
EOF1
TAPE MARK
TAPE MARK

NOTES:

10-12

MULTI-REEL MULTI-FILE

REEL 1

VOL 1
HDR 1
TAPE MARK
FILE A
TAPE MARK
EOF1
TAPE MARK
HDR1
TAPE MARK
FILE B
TAPE MARK
EOV1
TAPE MARK
TAPE MARK

REEL 2

VOL 1
HDR 1
TAPE MARK
FILE B
CONTINUED
TAPE MARK
EOV1
TAPE MARK
TAPE MARK

REEL 3

VOL 1
HDR 1
TAPE MARK
LAST OF FILE B
TAPE MARK
EOF1
TAPE MARK
HDR1
TAPE MARK
FILE C
TAPE MARK
EOV1
TAPE MARK
TAPE MARK

LABELS - 80 CHARACTERS

10-13

VOLUME HEADER LABEL

3 CHAR.	1 CHAR.	6 CHAR.	1 CHAR.	1 CHAR.	60 CHAR.
V O L	1	VISUAL REEL NUMBER (VRN)	BLANK /	BLANK 0 1 2	19 SPACES FOR OPERATING SYSTEM 49 SPACES FOR FUTURE STAN.
LABEL IDENT.		LABEL NUMBER		SECURITY	DENSITY

FILE HEADER LABEL

1 3 CHAR.	1 CHAR.	20 CHAR.
H D R	1	FILE LABEL NAME
LABEL IDENTIFIER LABEL		

21	25 3 CHAR.	28 4 CHAR.	32 4 CHAR.	36 3 CHAR.	
	MULTI FILE IDENTIFICATION (MFN)	REEL NUMBER	MULTI-FILE POSITION NUMBER	RESERVED "SPACES"	ED. NO.

41	42	43	5 CHAR.	48	49	5 CHAR.	54	55	6 CHARS.
			CREATION DATE YYDDD			EXPIRATION DATE YYDDD			BLOCK COUNT "ZEROS"
"SPACE"			"SPACE"			SECURITY			

61	20 CHARACTERS
RESERVED FOR FUTURE STANDARDIZATION "SPACES"	

10-14

1 3 CHAR	1 CHAR	20 CHAR	20
EOF or EOV	1	FILE LABEL NAME	

LABEL IDENTIFIER LABEL

#

21	25 3 CHAR.	28 4CHAR.	36 3CHAR.	40
	MULTI FILE IDENTIFICA- TION (MFN)	REEL NUMBER	MULTI FILE POSITION NUMBER	RESERVED "SPACES"
				ED. NO.

41	42	43	5 CHAR.	48	49	5 CHAR.	54	55	6 CHAR.	60
	SPACE		CREATION DATE YYDDD	SPACE		EXPIRATION DATE YYDDD	SECURITY		BLOCK COUNT NON ZERO = PRU COUNT	

61	20 CHARACTERS	80
RESERVED FOR FUTURE STANDARDIZATION		
"SPACES"		

10-15

FET TAPE LABEL INFORMATION

59				17	0
1	LOGICAL FILE NAME				
2	<div> <div>FILE</div> <div>PP</div> </div>			FIRST	
3				IN	
4				OUT	
5				LIMIT	
6					
7					
8					
9					
10	FILE LABEL NAME (FIRST 10 CHARACTERS)				
11	FILE LABEL NAME (LAST 10 CHARACTERS)				
12	EDITION NUMBER	RETENTION CYCLE	CREATION DATE		
13	POSITION NUMBER	MULTI-FILE NAME	REEL NUMBER		
59	47	41	29	23	0

7-23 → 7-25

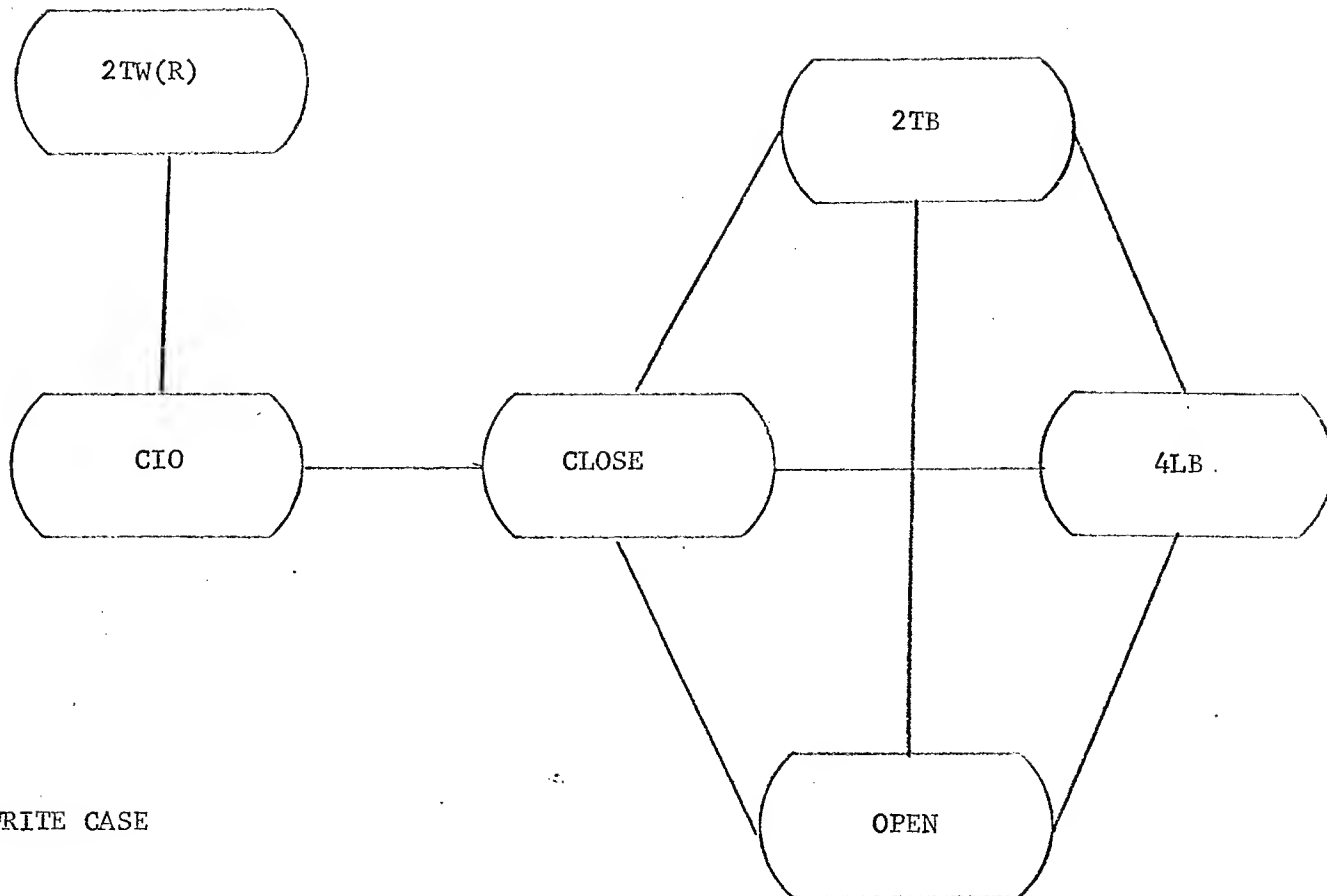
LABEL REQUESTS

REQUEST, LFILE, N, 2LD

REQUEST, MFI, MF, N.

END OF TAPE PROCESSING

10-16



A. WRITE CASE

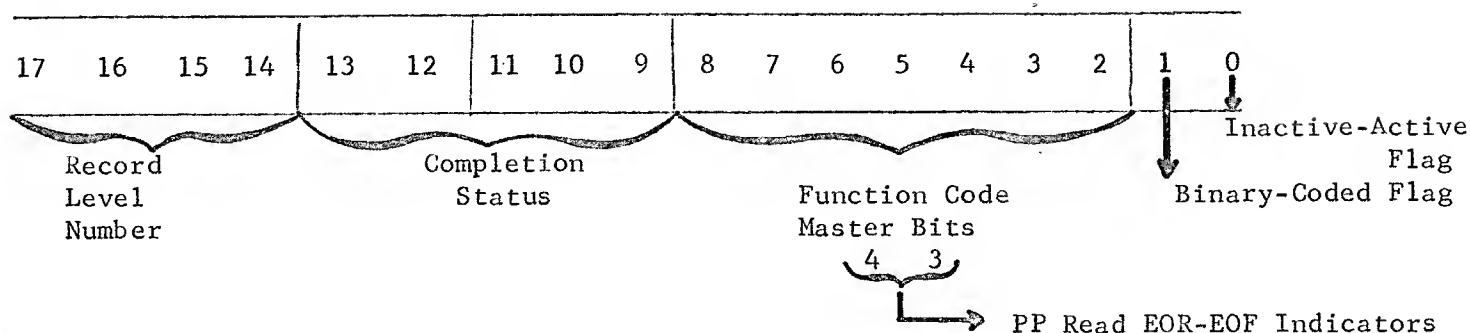
LABELLED TAPE	Y	N	Y	N
UP BIT SET	N	N	Y	Y
	1	2	3	4

B. READ CASE

LABELLED TAPE	Y	N	Y	N
UP BIT SET	N	N	Y	Y
	5	6	7	8

FET-FNT

Code & Status

COMPLETION STATUS (BITS 9-13)

01 = EOI on For. or Back. Read
 02 = End-of-Reel on For. or Back. Operation when up bit is set.
 04 = Unrec. parity error
 10 = Device of buffer capacity exceeded.
 20 = Open function redundant
 21 = Close function redundant
 22 = Illegal function
 23 = Index full
 24 = FNT full

ADDITIONAL STACKORDER CODES

02 = Read-CM-Program (3 wrds lost)
 10 = PP Read
 11 = PP Overlay Read
 14 = PP Write No EOR
 15 = PP Write EOR
 17 = Release Chain

FUNCTION CODE

(Bits 0-8)

Bits 2-8

Stack
Order
Code

*00X = Illegal	00-01	--
*010 = Read	02	00
*014 = Write	03	04
020 = Read-Skip	04	01
*024 = Write-EOR	05	05
*034 = Write EOF	07	05
*040 = Backspace	10	13
044 = Backspace PRU	11	16
*050 = Rewind	12	
*054 = Rewind	13	
*060 = Unload	14	
*064 = Unload	15	
100 = Open	20	
114 = Evict	114	
240 = Skip-For.	50	12
640 = Skip-Back.	150	13

INACTIVE-ACTIVE (BIT 0)

*0 = Active
 *1 = Inactive

BINARY-CODED (BIT 1)

*0 = Coded
 *1 = Binary

PP READ (BITS 3-4)

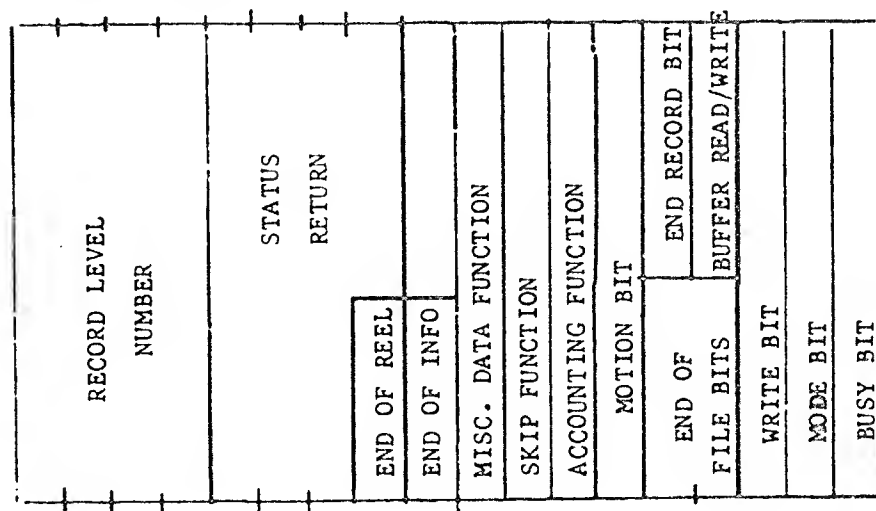
*10 = EOR Read
 *11 = EOF Read

* Identical Meaning in Scope 2.0

2.0/3.0 CIO operation code comparison

CODE	2.0 USAGE	3.0 USAGE	3.0 SYSTEM MACRO FORM	
00	NOT USED	READ PHYSICAL REC.	RPHR	LFN, RECALL
04	NOT USED	WRITE PHYSICAL REC.	WPHR	LFN, RECALL
10	READ	READ	READ	LFN, RECALL
14	WRITE	WRITE	WRITE	LFN, RECALL
20	NOT USED	READ SKIP	READSKP	LFN, 1, RECALL
24	WRITE RECORD	WRITE RECORD	WRITER	LFN, 1, RECALL
30	SEARCH FILE FORW.	ILLEGAL		
34	WRITE END FILE	WRITE END FILE	WRITEF	LFN, RECALL
40	BACKSPACE	BACKSPACE	BKSP	LFN, RECALL
44	BACKSPACE	BACKSPACE PHYSICAL	BKSPRU	LFN, N, RECALL
50	REWIND	REWIND	REWIND	LFN, RECALL
60	UNLOAD	UNLOAD	UNLOAD	LFN, RECALL
114	NOT USED		EVICT	LTN, RECALL
130	SEARCH FILE FORWARD	ILLEGAL		
240	NOT USED	SKIP FORWARD	SKIPF	LFN, N, 1, RECALL
640	NOT USED	SKIP BACKWARDS	SKIPB	LFN, N, 1, RECALL

BIT 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0



LAST CODE AND STATUS
Found in FET(1) and
FNT/FST(3)

BIT	USE	SYMPTOM	
0	BUSY BIT	0	BUSY
		1	FREE
1	MODE BIT	0	CODED
		1	BINARY
2	WRITE BIT	0	READ
		1	WRITE
3,4	END OF FILE BITS	0	NO EOF
		3	EOF
4	END OF RECORD BIT	0	NO EOR
		1	EOR
5	MOTION BIT	0	FORWARD
		1	REVERSE
6	ACCOUNTING BIT	0	NO OPEN/CLOSE
		1	OPEN/CLOSE ACTION
7	SKIP FUNCTION BIT	0	NO SKIP
		1	SKIP
9	END OF INFO BIT	0	NO EOI
		1	EOI

STATUS-OPCODE BREAKDOWN-3.0 CIO CALL

10	END OF TAPE BIT	0	NO END OF TAPE
		1	END OF TAPE
9-13	STATUS RETURN	04	UNRECOVERABLE PART
		10	CAPACITY ERROR
		20	OPEN REDUNDANT
		21	CLOSE REDUNDANT
		22	ILLEGAL FUNCTION
		23	INDEX FULL
		24	FNT FULL

- EXAMPLES -

1. BASIC DECK STRUCTURE
2. Simple FORTRAN JOB.

JOB CARD

REQUEST CARDS

COMPILE (OR. ASSEMBLE) CARD(S)

EXECUTION CARDS

ERROR CONTROL CARDS

7-8-9

PROGRAM EXAMPLE (INPUT, OUTPUT)

END

7-8-9

DATA

7-8-9

6-7-8-9

SEE FOLLOWING EXAMPLES FOR FORTRAN
AND ASSEMBLER RUNS.

Dec 16, 1967
14-De

To run a Compass program on 6000 Scope 3 without using the Fortran compiler -

Col 11

(1st card) → IDENT NAME
ENTRY MOVE
MOVE (1st instruction)

"PERIPH" CARD IF
PROGRAM IS PP TYPE

(Last instruction) (NOT A "PS")
 END RUN (STORES "END" IN RA+1)
 AREA B DATA 47 -- etc

DATA, Common area, etc.

(Last cmd) → . END MOVE

3
9

TERMINATE ASSEMBLER
AND LOOKS FOR TRANSFER
ADDRESS. (MUST HAVE)

CORE DUMPS AND RELATED DEBUGGING AIDS

Robert W. Bartlett
Control Data Corporation

When a program terminates abnormally, the SCOPE monitoring system provides a one page dump of information consisting of:

1. The contents of all 24 operating registers
2. The contents of 7 core memory locations whose addresses are in registers A1 thru A7
3. A 200B word dump of the 100B words before and after the core address where the error was detected
4. A variable number of words beginning at location 000000 which contain system information, including the names and buffer parameter addresses of all files referenced by the program.

Only in rare instances will the average applications programmer be able to make use of items (3) and (4) in the above list. Rather, he should typically make use of the following tools:

- A. Items (1) and (2) above
- B. A core map of routines, entry points, and references
- C. A compilation listing of applicable routines

The concerted use of these tools should allow the programmer to pinpoint the FORTRAN statement that was executing when the error occurred. In addition, it is often possible to determine the current value of several variables used in the statement. From then on, the programmer's intimate knowledge of his own program should allow him to deduce what combination of faulty input, or coding, or logic caused this particular error to occur at this particular place. Let us now discuss in greater detail what can be found in these three debug tools.

A. One page dump (Figure 4)

The 24 operating registers consist of:

- 8 address registers (A0 thru A7) - 18 bits each
- 8 increment registers (B0 thru B7) - 18 bits each
- 8 operand registers (X0 thru X7) - 60 bits each

The contents of the A registers should correspond to addresses of memory locations within the field length of the job. An address in an A register equal to or greater than the field length specified on the job card would normally be associated with a mode 1 arith error.

To the right of the A and B register contents on the dump is a list of contents of memory locations specified by the associated A registers. For example, the contents of memory location 123B referenced in address register A3 are listed as C(A3)=
1721 6000 0000 0000 0000.

The contents of the X registers correspond to the values of operands involved in arithmetic or alpha-numeric operations. The programmer's eye should be trained to recognize the following error indications in the C and X lists.

+ Indefinite	1777	----	----	----	----
+ Infinity	3777	----	----	----	----
- Infinity	4000	----	----	----	----
- Indefinite	6000	----	----	----	----

B. Core map (Figure 3)

From the core map of the attached sample program (figure 1), we can learn the following things which are true of core maps in general.

The first two lines tell us that the job loaded between addresses 100B and 2060B and that storage for BLANK COMMON begins at 2056B and runs for a length of 2 words.

The table headed - PROGRAM----ADDRESS - tells us that program DMPREAD began loading at address 103B and subroutine SYSTEM began at address 1143B. To the right of this table are the names and beginning addresses of the labeled common blocks associated with each routine. For example, the labeled common block named BLOK loaded at address 100B preceeding the program (DMPREAD) that first declared it.

The table headed --ENTRY-----ADDRESS- tells us that program DMPREAD has an entry point named DMPREAD at address 104B. The remaining eight entry points all belong to subroutine SYSTEM

since they all occur at locations greater than the loading address of SYSTEM at 1143B. They are:

Q8NTRY	at location	1144B
SYSTEM	at location	1307B
SYSTEMC	at location	1254B
SYSTEMP	at location	1302B
END	at location	1200B
STOP	at location	1227B
EXIT	at location	1221B
ABNORMAL	at location	1237B

In addition we notice that entry point Q8NTRY is referenced by location 105B found in DMPREAD and the entry point END is referenced by location 116B also found in DMPREAD. Since DMPREAD began loading location 103B, these calls are actually from relative address 2 and 13B in DMPREAD.

The table headed ---- UNSATISFIED EXTERNALS----- tells us that we forgot to load the subroutine SUB with this run. This subroutine is called from relative location 11B in DMPREAD, and the program would have quit on a mode 1 arith error at that point if other errors had not caused us to exit sooner.

C. Compilation listing (Figure 1)

The numbers along the left hand side of the compilation listing are the relative locations within the program at which the corresponding FORTRAN statements begin execution. For example, the set of machine language instructions necessary to perform the calculation $AO = BO + CO * 2$ begins execution at relative location 6 in DMPREAD or at absolute location 111B in memory since DMPREAD begins at location 103B according to the core map.

On the page of information below the compilation listing is the following useful information. (Figure 2)

The list headed BLOCK NAMES AND LENGTHS tells us there are two common blocks associated with program DMPREAD. Block 01 is blank COMMON (the block name has been left blank) which is 2 words long, and block 02 is labeled COMMON which is 3 words long and is named BLOK.

The names and relative storage locations of variables used in program DMPREAD are:

AO at relative address 26B and absolute address 131B
 CO at relative address 25B and absolute address 130B
 CO at relative address 01 in common block 02 and absolute address 101B
 BO at relative address 01 in common block 01 and absolute address 2057B
 BN, CN, and CP do not appear among the list of variables because they do not appear in any executable statement within the program.

Now that we know what can be found in core dumps, core maps, and compilation listings, let us see how they can be used together to find the source of error in the attached example. (Figure 1)
 But before we start let us list the three distinct arith errors that occur.

MODE1 Trying to reference an address that is equal to or greater than the field length.

MODE2 Trying to operate with an infinite operand (3777---
 -----or 4000-----)

MODE4 Trying to operate with an indefinite operand (1777---
 -----or 4000-----)

MODE3, MODE5, MODE6, MODE7 are algebraic combinations of these three i.e. MODE5 means MODE1 and MODE4.

Thus we will begin debugging by reading the DAYFILE message (figure 5) telling us that we have a MODE4 arith error at address 115B. Looking at the core map we see that program DMPREAD begins at location 103B and subroutine SYSTEM begins at address 1143B. Therefore, the error occurred in program DMPREAD and is in fact at relative address twelve in DMPREAD.

000115B
 -000103B
 =12B

If we now look at the listing for program DMPREAD, we see that location 12 corresponds on the processing of the END card. To escape from this dilemma, we must apply the following rule of thumb corrections to 6400 and 6600 addresses.

6600 Subtract 3 from error address
6400 Subtract 1 from error address

Then we find that the error occurred at relative address 7 while processing the FORTRAN statement $AO=BO+CO**2$.
 000012B
 -000003B
 = 7B

This is because the machine is looking ahead toward execution of future instructions before it discovers the error condition from the previous operation. Because of parallel operation the 6600 is looking further ahead than the serial operating 6400.

Now, knowing the logic of his program and the statement where the error showed up, the programmer might be in a position to say to himself as follows:

"The value of BO or CO must be indefinite since arith errors MODE1 thru 7 only occur during arithmetic operations on specious operands. But CO is equal to 3.0 and BO is equal to CO and both are good values. Oops - wait a minute there is a mispunch, BO = C0 and C0 is not previously defined on the left hand side of an arithmetic replacement statement. It is good that SETINDF (LRC programing manual, Q1.001) helped me to catch this mistake or there is no telling what answers I would have gotten on this or other machines."

In a longer program where the logic was not as easy to trace back, the debugging programmer might go on to look at the register dumps for clues. Indeed, in our example we see that the contents of the memory locations referenced in address registers A1, A4, and A6 have the distinctive 1777 byte characteristic of indefinite operands.

C (A1) = Contents of 002057 = 1777 0000 0000 0000 0130B
 C(A4) = Contents of 000130 = 1777 0000 0000 0000 0130B
 C(A6) = Contents of 000131 = 1777 0000 0000 0000 0000B

Making use of the core map and doing the familiar octal subtractions, we find that A4 and A6 are referencing the 25th and 26th locations in DMPREAD, and A1 is referencing the 01 element in BLANK COMMON.

A4	A6	A1
000130B	000131B	002057B
-000103B	-000 103B	002056B
= 25B	=26B	=1B

Then looking at the list of variable assignments on the page below the compilation listing for DMPREAD, we see that

A1 is referencing BO - 000001C01
 A4 is referencing C0 - 000025
 A6 is referencing AO - 000026

By some thunderbolt technique we should soon discover that C0 is unwanted and undefined.

Final mention should be made of the use of a library routine called SETINDF as an aid in debugging this example. The control card SETINDF. placed after the RUN(S) and before LGO. Writes the indefinite indicator 1777 into the upper byte of every memory location from address 100B up to your field length. In addition, it writes each memory location's own address into the bottom half of each word. For example, core address 200B contains 1777 0000 0000 0000 0200B. Thus, when your program references a variable that has not been defined, the preset indefinite value will be used and should soon cause a MODE4 arith error to occur. An indefinite value caused by arithmetic operations of the machine will look like 1777 0000 0000 0000 0000B, i.e. the 1777 indefinite flag will be followed by all zeroes to the right. Thus, the telltale 130 in C(A4) and C(A1) should have sent us immediately looking for the undefined variable C0.

AA.005
CORE DUMP

```
PROGRAM DMPREAD(OUTPUT=1001)
000003 COMMON BN , RO
000003 COMMON /BLOK/ CN , CO , CP
000003 CO = 3.0
000005 RO = CN
000006 AU = BN+CN**2
000011 CALL SUB
000012 END
```

8-14-67
RWB
Page 7 of 11

14-1

Figure 1 Compilation Listing of Sample Program

PROGRAM LENGTH INCLUDING I/O BUFFERS

001040

FUNCTION ASSIGNMENTS

STATEMENT ASSIGNMENTS

BLOCK NAMES AND LENGTHS

- 000002 BLUK - 000003

VARIABLE ASSIGNMENTS

A0 - 000026 H0 - 000001C01 C0 - 000001C02 C0 - 000025

START OF CONSTANTS

000014

START OF TEMPORARIES

000022

START OF INDIRECTS

000025

UNUSED COMPILER SPACE

006600

Figure 2 Program Information

2-11

AA.005
CORE DUMP

CORE MAP	11.42.26.	NORMAL	CONTROL	000100	002060	002056	000002
---	TIME---	LOAD MODE	--L1--L2---	TYPE---	USER---	CALL---	----
FWA LOADER	035045	FWA TABLES	034771	----	FWA LOAD---	LWA LOAD---	BLNK COMN--LENGTH--
PROGRAM---	ADDRESS---	----	LABELED---	COMMON---	----	----	----
DMPREAD	000103	----	HLOK	000100	----	----	----
SYSTEM	001143	----	----	----	----	----	----
ENTRY---	ADDRESS---	----	----	REFERENCES	----	----	----
DMPREAD	000104	----	----	----	----	----	----
QENTRY	001144	DMPREAD	000105	----	----	----	----
SYSTEM	001307	----	----	----	----	----	----
SYSTEMC	001254	----	----	----	----	----	----
SYSTEMP	001302	----	----	----	----	----	----
END	001200	DMPREAD	000116	----	----	----	----
STOP	001227	----	----	----	----	----	----
EXIT	001221	----	----	REFERENCES	----	----	----
ABNORMI	001237	----	----	----	----	----	----
----	UNSATISFIED EXTERNALS----	----	----	----	----	----	----
SUB	----	----	000114	----	----	----	----

Figure 3 Core Map

000021	14071	70000	00000	00047	00000	00000	00000	22676	000025	00000	00000	00000	32357
000031	11162	62520	00000	00021	00000	00000	00000	22676	00000	00000	00000	00000	37015
00000	00000	00000	00000	00000	000055	00000	00000	22676	00000	00000	00000	00000	00002
000062	00000	00000	00000	00000	00000	00000	00000	40000	00000	00000	00000	00000	35045
00000	00000	00000	00000	02060	00000	00000	00000	00103	00000	00000	00000	00000	55522
05301	12457	55000	00000	00000	22052	12505	23245	50001	22054	15610	11575	00000	00000
03040	35522	05051	45510	00000	00000	00000	00000	40000	00000	00000	00000	00000	00101
17214	00000	00000	00000	00000	17770	00000	00000	00102	04152	02205	01040	00101	00130
01000	01140	00000	00000	00000	01100	00121	10630	40000	01600	00101	51600	00130	40000
51500	00101	00255	40000	00000	00100	00057	30712	20007	01600	00131	40000	40000	00027
51300	00124	10730	40000	00000	01600	01700	07000	00103	00000	00000	00000	00000	00000
00000	00000	00000	10000	00000	17216	00000	00000	00107	17216	00000	00000	00000	00142
000126	17770	00000	00000	00126	17770	00000	00000	00000	00000	00000	00000	00000	00142
00000	00000	00000	00000	00000	00000	00000	00000	10000	17770	00000	00000	00000	00147
17770	00000	00000	00145	17770	00000	00000	00000	00146	17770	00000	00000	00147	00153
17770	00000	00000	00151	17770	00000	00000	00000	00152	17770	00000	00000	00157	00163
17770	00000	00000	00155	17770	00000	00000	00000	00156	17770	00000	00000	00167	00173
17770	00000	00000	00161	17770	00000	00000	00000	00162	17770	00000	00000	00177	00203
17770	00000	00000	00165	17770	00000	00000	00000	00166	17770	00000	00000	00207	00213
17770	00000	00000	00171	17770	00000	00000	00000	00172	17770	00000	00000	00000	00000
17770	00000	00000	00175	17770	00000	00000	00000	00176	17770	00000	00000	00000	00000
17770	00000	00000	00201	17770	00000	00000	00000	00202	17770	00000	00000	00000	00000
17770	00000	00000	00205	17770	00000	00000	00000	00206	17770	00000	00000	00000	00000
17770	00000	00000	00211	17770	00000	00000	00000	00212	17770	00000	00000	00000	00000
00000	00000	00000	00000	17252	42025	24000	00132	00000	00000	00000	00000	00000	00000

14.35.41. PTB1967. READ.
 14.35.42. PTR1967. PP 000 SEC.
 15.00.06. PTB1967. PTR1967.
 15.00.06. PTR1967. LRC COMPUTER COMPLEX
 15.00.06. PTR1967.
 15.00.06. PTR1967. JOB.01.100.40000. P9970.12234.2.
 15.00.06. PTR1967. R W HARTLETT , CDC ,
 15.00.07. PTR1967. RUN(S)
 15.00.11. PTR1967. SETINDF.
 15.00.11. PTR1967. LGO.
 15.00.13. PTR1967. ARITH ERROR.
 15.00.13. PTR1967. MODE = 4. ADDRESS = 000115
 15.00.13. PTR1967. CP 000.106 SEC.
 15.00.13. PTR1967. PP 002.453 SEC.
 15.00.22. PTR1967. PRINT-PP 000 SEC.
 06/01/67. - SCOPE 2.0 - LRC 6600 - B -05/17/67.

AA.005
 CORE DUMP

VER 1

SETINDF

061067L006

		000013	IDENT	SETINDF
			PROGRAM LENGTH	
			BLOCKS	
	CCCC00	000013	PROGRAM* LOCAL	
			ENTRY POINTS	
			000006 SETINDF	
			ENTRY	SETINDF
000000	56620		SET SA6	B2
	36661		IX6	X6+X1
	66221		SB2	B2+B1
000001	C7230CCCC00 +		LT	B2,B3,SET
	7170C51604		SX7	516048
000002	20752		LX7	42
	56710		SA7	B1
	10644		BX6	X4
	54640		SA6	A4
000003	36661		IX6	X6+X1
	54661		SA6	A6+B1
	36661		IX6	X6+X1
	54661		SA6	A6+B1
000004	36661		IX6	X6+X1
	54661		SA6	A6+B1
	36661		IX6	X6+X1
	54661		SA6	A6+B1
000005	23052411160406CCCC00		VFD	42/0HSETINDF,18/0
000006	46000		SETINDF NO	
	6110C00001		SB1	1
	76110		SX1	B1
000007	6120000005 +		SB2	SETINDF-1
	76220		SX2	B2
	64300		SB3	A0
000010	5140000000 +		SA4	SET
	74040		SX0	A4
000011	7140C01777		SX4	17773
	20460		LX4	48
	36642		IX6	X4+X2
000012	36440		IX4	X4+X0
	C2C00000000 +		JP	SET
000013			END	SETINDF

6000 SLOPE 3

14-13

CREATE EDITSYM LISTING OF A SIMPLE PROGRAM

PAGE - 1 →

TEST006

PAGE - 2 →

TEST006

PAGE - 3 IS BLANK

PAGE - 4 →

EDITSYM CONTROL CARDS

*DECK, TRYIT

PAGE - 5 THRU 8 ARE THE FOLLOWING PAGES.

PAGE - 9 IS BLANK

PAGE - 10 IS THE "IDENT TRY COMP" PAGE

PAGE - 12 IS THE "CORE MAP"

PAGE - 13 IS THE "DMPX"

PAGE - 14 IS THE "DAYFILE"

EDITSYM LIST

	PROGRAM TRYIT(INPUT,OUTPUT)	TRYIT00001
	COMMON I,J,K,L,M	TRYIT00002
	READ 10,I,J,K,L	TRYIT00003
10	FORMAT(4(I5,X1))	TRYIT00004
	CALL TEST(I,J,K,L,M)	TRYIT00005
	PRINT 11,M	TRYIT00006
11	FORMAT(1H1,I7)	TRYIT00007
	STOP	TRYIT00008
	END	TRYIT00009
	IDENT TRYCOMP	TRYIT00010
	ENTRY TEST	TRYIT00011
	LIST L,R	TRYIT00012
	VFD 42/7LTRYCOMP,18/5 .GENERATE THE HES	TRYIT00013
TEST	BSSZ 1 .RESERVE A WORD FOR THR RETURN JUMP INSTRUCTION	TRYIT00014
	SA1 B1 .LOAD THE FIRST CONSTANT FROM I (B1)	TRYIT00015
	SA2 B2	TRYIT00016
	SA3 B3	TRYIT00017
	SA4 B4	TRYIT00018
	IX5 X1-X2 .DIFF OF I AND J TO X5	TRYIT00019
	ZR X5,ADD .JUMP TO K+L IF X5 IS ZERO	TRYIT00020
	PL X5,SUB .JUMP TO K-L IF X5 IS POSITIVE	TRYIT00021
	IX6 X4-X3 .SUB L-M	TRYIT00022
	EQ STORE	TRYIT00023
ADD	IX6 X3+X4 .ADD K TO L	TRYIT00024
	EQ STORE	TRYIT00025
SUB	IX6 X3-X4 .SUBTRACT K FROM L	TRYIT00026
STORE	SA6 B5	TRYIT00027
	EQ TEST .RETURN TO THE RETURN JUMP INSTRUCTION	TRYIT00028
	END	TRYIT00029

1/21-194

PROGRAM TRYIT(INPUT,OUTPUT)

TRYIT00001

000000 0 L00002
000001 6110 L00002 L00001
6120 C00001
000002 0100 S00100

COMMON I,J,K,L,M
PEAK 10,I,J,K,L

TRYIT00002
TRYIT00003

000003 6130 N00001
6120 C00005
000004 67202
66100
000005 0100 S00200 L00005
0700 L00002
000006 66200
6110 V00001
000007 0100 S00200
000010 66200
6110 V00002
000011 0100 S00200
000012 66200
6110 V00003
000013 0100 S00200
000014 66200
6110 V00004
000015 0100 S00200
000016 6110 777776
0100 S00200

10 FORMAT(4(I5,/))
CALL TEST(I,J,K,L,M)

TRYIT00004
TRYIT00005

000017 6110 V00001
6120 V00002
000020 6130 V00003
6140 V00004
000021 6150 V00005
000022 0100 S00300 L00007
0705 L00002

PRINT 11,M

TRYIT00006

000023 6130 N00002
6120 C00012
000024 67202
66100
000025 0100 S00400 L00012
0700 L00002
000026 66200
6110 V00005
000027 0100 S00400
000030 6110 777776
0100 S00400

11 FORMAT(11H1,I7)

TRYIT00007

STOP

TRYIT00008

000031 5130 C00015

10733

000032 0100 S00500 L00014

0700 L00002

END

TRYIT00009

000033 5140 C00016

10744

000034 0100 S00600 L00015

0700 L00002

000034

PROGRAM LENGTH INCLUDING I/O BUFFERS

004121

FUNCTION ASSIGNMENTS

STATEMENT ASSIGNMENTS

10 000044 11 000050

BLOCK NAMES AND LENGTHS

000005

VARIABLE ASSIGNMENTS

I 000000C01 J 000001C01 K 000002C01 L 000003C01
M 000004C01

START OF CONSTANTS

000035

START OF TEMPORARIES

000053

START OF INDIRECTS

000055

UNUSED COMPILER SPACE

043600

14-77

IDENT-TRYCOMP

TRYITO

000011 PROGRAM LENGTH

BLOCKS

000000 000011 PROGRAM* LOCAL

ENTRY POINTS

000001 TEST

ENTRY TEST

000000 24223103171520000005

000001

000002 56110

TEST

56220

56330

56440

000003 37512

0305000006 +

000004 0325000007 +

37643

000005 0400000010 +

000006 36634

ADD

0400000010 +

000007 37634

SUB

000010 56650

STORE

0400000001 +

000011

LIST L,R

VFD 42/7LTRYCOMP,18/5

BSSZ 1

SA1 B1

SA2 B2

SA3 B3

SA4 B4

IX5 X1-X2

ZR X5,ADD

PL X5,SUB

IX6 X4-X3

EQ STORE

IX6 X3-X4

EQ STORE

IX6 X3-X4

SA6 B5

EQ TEST

END

.GENERATE THE MES

.RESERVE A WORD FOR THR RETURN JUMP INSTRUCTION

.LOAD THE FIRST CONSTANT FROM I (B1)

.DIFF OF I AND J TO X5

.JUMP TO K+L IF X5 IS ZERO

.JUMP TO K-L IF X5 IS POSITIVE

.SUB L-K

.ADD K TO L

.SUBTRACT K FROM L

.RETURN TO THE RETURN JUMP INSTRUCTION

TRYIT00011

TRYIT00012

TRYIT00013

TRYIT00014

TRYIT00015

TRYIT00016

TRYIT00017

TRYIT00018

TRYIT00019

TRYIT00020

TRYIT00021

TRYIT00022

TRYIT00023

TRYIT00024

TRYIT00025

TRYIT00026

TRYIT00027

TRYIT00028

TRYIT00029

053330

UNUSED STORAGE

20 STATEMENTS

4 SYMBOLS

SYMBOLIC REFERENCE TABLE

ADD	0000006	PROGRAM*	000003	
STORE	0000010	PROGRAM*	000005,	000006
SUB	0000007	PROGRAM*	000004	
TEST	0000001	PROGRAM*	000010	

CORE MAP 0003738 NORMAL CONTROL 000100 010515 010506 000005
 ---TIME---LOAD MODE --L1--L2---TYPE-----USER---+---CALL-----FWA LOAD--LWA LOAD--BLNK COMM--LENGTH--

FWA LOADER 074614 FWA TABLES 074403

--PROGRAM---ADDRESS-- --Labeled--COMMON--

TRYIT 000100
 TRYCOMP 004221
 SYSTEM 004232 SCOPE2 004232
 INPUTC 005207
 OUTPTC 006306
 SIO. 007640
 GETBA 010467

--ENTRY---ADDRESS-- REFERENCES

TRYIT 000101
 TEST 004222 TRYIT 000122
 QBNTY 004233 TRYIT 000102
 SYSTEM 004433 INPUTC 005246 006117
 OUTPTC 006322 007451

SYSTEMC 004400
 SYSTEMP 004426
 END 004324 TRYIT 000134

STOP 004353 TRYIT 000132

EXIT 004345
 ABNORML 004363 INPUTC 005247 006120
 OUTPTC 006323 007452

INPUTC 005211 TRYIT 000105 000107 000111 000113 000115 000116

KRAKER 005313
 OUTPTC 006310 TRYIT 000125 000127 000130

KODER 006444
 BKSPRU 010136
 FI74AK 010147
 POSFIL 010205 OUTPTC 006340

RQPRU 010215
 DAT 010237 INPUTC 005265 005233 005256
 OUTPTC 006362 006376 006343

CIOI 010103
 OPEN 007642 SYSTEM 004704
 INPUTC 005231
 OUTPTC 006330

SIO 007755 INPUTC 005261
 OUTPTC 006375

GETBA 010467 INPUTC 005217
 OUTPTC 006316

---UNSATISFIED EXTERNALS---

REFERENCES

64146

000
20
000
000
+10
433
027
000
502
000

02/05/68 SCOPE V3.1 6500-00 131K 11/21/67

00.50.49.TEST006

00.50.49.TEET,P17,T100,CM100000.

```
00.50.49.EDITSYM(I,C,L,,NPL)
```

•GENERATE AN EDITSYM LISTING OF THE

00.50.49.PROGRAM

```
00.50.52.RUN(M,,,COMPILE)
```

00.51.03.LGO.

00.51.06. ARITH ERROR.

00.51.06.MODE = 1. ADDRESS =100001

00.51.06.CP 005.159 SEC.

00.51.06.PP 014.537 SEC.

the following listing illustrates the results of using UPDATE to obtain an assembly of a system program for later modifications.

the 1st Page is -

1-----
2--- SCPD01F-----
3-----
4-----
5-----
6-----

the 2nd Page is -

1-----
2--- SCPD01F-----
3-----
4-----
5-----
6-----

the 3rd page is blank

the following pages are the 4th and on.

*IDENT+RAK
COMPILE.MSG

COMP-ILE	ISPISQ	FMTSKCH	READPP	BUFINIT	RELOC
PPIME	CUMCOMP	TRNCOM	TAPEDS	DISCUS	DUMPOS
RECOV	PREPP	PRECP	STLPP	STLCP	CMR
MR	DSU	CIU	LDR	MSG	OPE
1AJ	1BJ	1LJ	1OT	1SP	1SW
2HP	2DF	2FS	2TR	2TS	2TW
7SO	ATS	CHK	CKP	CLL	CLO
CLS	CPL	CTS	DIS	DMP	EXU
LHC	LOC	LOD	MDI	MEM	PBC
KHR	KEW	RFL	RST	SRB	TIM
WBR	1BT	1CO	1CY	1OF	1EE
1LT	1MF	1MR	1MW	1PL	1PO
1KA	1KU	1SA	1TU	2BT	2CA
2CF	2CJ	2EF	2LA	2LB	2LE
2LP	2PC	2PC	2RT	2TB	2TF
2TU	2WT	3OT	4LB	5DT	7TP
LOAVER	OVERLOD	QXX	EDITLIB	EOITSYM	COPYBCD
CUTL	CPC	IORANDM	IO	COMPARE	RESTART
BKSP	COPY	COPYBF	COPYN	COPYSBF	REWIND
UPDATE	CVRT	ACGOER	DBLE	EXP	GETHA
IOALEX	LEGVAR	LOCF	SINCOS	SNGL	SQRT
SYSTEM	TAN	XRCL	ALNLOG	ASINCOS	ATAN
ATAN2	BACKSP	BUFFE1	BUFFE0	CAHS	CBALEX
CCUS	CEXP	CLOG	CSIN	CSQRT	DABS
DAFAN	DBADEX	DBALEX	DEXP	DISPLA	DLNLOG
DMOD	DSIGN	DSINCOS	DSQRT	DUMP	DVCHK
ENUTIL	IDINT	IFENDF	INPUTB	INPUTC	INPUTN
INPUTS	IOCHEC	IOCHEK	LENGTH	OUTPTB	OUTPTC
OUTPTN	OUTPTS	OVERFL	OVERLAY	PAUSE	RANF
RBALEX	RBAHEX	REMAKK	REWIND	SCOPE2B	SECOND
SEGMENT	SIOB	SLITE	SLITET	SSWTCB	START
TANN	TIME	COMPASS	CVDEL	RUN	RUN1
UBUAGP	TRANS	SYSTEXT	SCOP31A	SCOP31B	IPARAMS
SCPTXT	SCUP31C	SCUP31E	SCUP31D	SERLO	87S0019
86C0186	86S0351	86S0204	86S0248	86S0257	86S0259
86S0275	86S0297	86S0393	86S0412	87S0445	86S0449
87S0480	87S0481	87S0482	87S0483	87S0524	87S0525
87S0527	86S0528	86S0529	86S0530	87S0531	86S0532
86S0533	86C0534	86F0535	91S0025	91S0041	89S0123
89S0253	89S0273	89S0365	91S0366	91S0402	89S0432
89S0446	87S0456	89S0458	89S0485	86S0495	89S0541
89S0550	87S0552	90S0582	89S0583	87S0588	87S0589
87S0590	87S0592	89S0595	89S0597	89S0598	89S0605
89S0607	89S0614	91S0664	91S0666	91S0667	91S0668
91S0669	91S0671	91S0672	91S0674	91S0675	91S0676
87F0333	87F0403	90F0406	86C0214	86C0217	86C0237
86C0244	86C0263	86C0281	86C0314	87C0490	87C0567
87C0568	87C0569	87C0580	91C0652	91C0663	91C0665
94S0747	SYSALT	91S0670	RAN		

COMPILE	ISPLSQ	ENTSRCH	READPP	BUFINIT	RELOC
PPIME	COMCOMP	TINCOM	TAPEDS	DISCOS	DUMPPDS
RECUV	PREPP	PRECP	STLPP	STLCP	CMR
MTR	DSU	CIO	LDR	MSG	OPE
IAU	IBJ	ILJ	IOI	ISP	ISW
ZBP	ZDF	ZES	ZTH	ZTS	ZTW
7SD	ATS	CHK	CKP	CLL	CLO
CLS	CPL	CTS	DIS	DMP	EXU
LNC	LUC	LNU	MDI	MEM	PBC
MBR	MEG	HFL	HST	SRB	TIM
WBR	IBT	ICO	ICY	IDF	IEE
ILT	IMF	IMR	IMW	IPL	IPU
IMA	IRU	ISX	ITU	2BT	2CA
ZCF	ZCU	ZEF	ZLA	2LB	2LE
ZLP	ZPC	ZHC	ZHT	2TB	2TF
ZTJ	ZWT	3OT	4LB	5DT	7TH
LOADEM	OVERLOD	QXX	EDITLIB	EDITSYM	COPYRCD
COPYL	CPC	IOHANDM	IO	COMPARE	RESTART
BRSP	COPY	COPY9F	COPYN	COPYSBF	REWINO
UPDATE	CVRT	ACGNR	UHL	EXP	GETBA
IBATEX	LEGVAR	LOC	SINCS	SNGL	SQRT
SYSTEM	TAN	XRCI	ALNLOG	ASINCS	ATAN
AIAN2	BACKSP	BUFEI	BUFEQ	CABS	CBAEX
CCUS	CEAP	CLUG	CSIN	CSQRT	UABS
UAIAN	UHADEX	UHAEX	UEAP	DISPLA	UNLOG
UMOU	USIGN	USINCS	USQRT	DUMP	DVCHK
ENDFIL	IDINT	IFENDF	INPUTB	INPUTC	INPUTN
INPUTS	LOCHEC	INCHEK	LENGTH	OUTPTB	OUTPTC
OUTPTN	OUTPTS	OVERFL	OVERLAY	PAUSE	RANF
KHAEX	KBAEX	REMARK	REWINM	SCOPE2B	SECOND
SEGMENT	SIU\$	SLITE	SLITET	SSWCH	START
TANT	TIME	COMPASS	CVUEL	RUN	RUN1
QADLAGP	TRANS	SYSTEM	IPARAMS	SCPTXT	USLCOM
CPSYS					

*WEOR-S, *CALLS-S AND MODIFICATIONS 02/08/68 PAGE NO. 00004

*YANK+SCOP31E

SERLO 00001

14-216

COMMON DECKS ENCOUNTERED

02/08/68

PAGE

NO.

00005

DSL COM

IPARAMS

ISP ISO

FNTSRCH

READPP

BUF INIT

RELUC

PPTIME

COMCOMP

TRNCOM

14-27

LIST OF DECKS WRITTEN ONTO COMPILE FILE 02/08/68 PAGE NO. 00006

MSU

30-71

IDENT	MSG.C.PPFWA	MSG	00002
001157	PROGRAM LENGTH		
	BLOCKS		
000000	001157	PROGRAM*	ABSOLUTE
	PERIPH	MSG	00003
	FUNCTION	MSG	00005
	MSG IS USED TO ENTER MESSAGES FROM A CENTRAL MEMORY PROGRAM	MSG	00006
	INTO THE DAYFILE	MSG	00007
		MSG	00008
		MSG	00009
		MSG	00010
		MSG	00011
		MSG	00012
	REVISIONS	MSG	00013
		MSG	00014
	REVISION 1.	MSG	00015
	8/2/66. TO PUT THE DAYFILE MESSAGE COUNTER IN	MSG	00016
	MSG INSTEAD OF MTR.	MSG	00017
	REVISION 2. (M. C. STEELE)	MSG	00018
	CONVERSION TO SCOPE 3.0	MSG	00019
		MSG	00020
	ENTRY-AND-EXIT INFORMATION	MSG	00021
	ENTRY	MSG	00022
	BITS 0-17 OF RA+1 SPECIFY THE ADDRESS OF THE MESSAGE	MSG	00023
	TO BE ENTERED. BITS 24-35 SPECIFY WHETHER THE	MSG	00024
	MESSAGE WILL BE ENTERED IN THE DAYFILE OR MERELY	MSG	00025
	DISPLAYED ON THE SCOPE. IF THESE BITS ARE ZERO, MESS.	MSG	00026
	WILL BE ENTERED IN DAYFILE NONZERO-JUST DISPLAYED.	MSG	00027
		MSG	00028
	EXIT	MSG	00029
	NONE	MSG	00030
		MSG	00031
		MSG	00032
	REGISTER AND BUFFER FORMATTING	MSG	00033
		MSG	00034
	BITS 42-59 OF IR SPECIFY THE PROGRAM NAME, BITS 36-38 SPECIFY	MSG	00035
	CONTROL POINT, BITS 24-35 SPECIFY WHETHER MESSAGE WILL BE	MSG	00036
	ENTERED INTO DAYFILE OR JUST DISPLAYED, BITS 0-17 SPECIFY	MSG	00037
	ADDRESS OF MESSAGE.	MSG	00038
		MSG	00039
	MSG PLACES THE RECEIVED MESSAGE INTO THE MESSAGE BUFFER	MSG	00040
		MSG	00041
		MSG	00042
		MSG	00043
	SST	SCOP31C	00260

19-7-68

1000		ORG	C.PPFWA	MSG	00046
1000	1401	LDN	1	SCOP31A	00419
1001	3470	STU	D.PPONE	SCOP31A	00420
1002	3075	LDU	D.PPIR	SCOP31A	00421
1003	6050	CRU	D.PPIRB	SCOP31A	00422
1004	3053	LDU	D.PPIRB+3	MSG	00048
1005	1237	LPN	378	SCOP31A	00423
1006	3464	STU	O.OUT	MSG	00049
1007	1014	SHN	12	MSG	00050
1010	3154	ADD	D.PPIRB+4	MSG	00051
1011	3405	STU	D.OUT+1	MSG	00052
1012	0200 0634	RJM	R.TFL	MSG	00053
1014	0732	MJN	MS3	SCOP31A	00424
1015	6040	CRU	D.BA	SCOP31A	00425
1016	3051	LDU	D.PPIRB+1	MSG	00055
1017	1220	LPN	208	MSG	00056
1020	3462	STU	D.IN	MSG	00057
1021	0412	ZJN	MSG4	MSG	00058
1022	3042	LDU	D.BA+2	MSG	00065
1023	1377	SCN	778	MSG	00066
1024	1006	SHN	6	MSG	00067
1025	3341	LMU	D.BA+1	MSG	00068
1026	1006	SHN	6	MSG	00069
1027	3405	STU	D.OUT+1	MSG	00070
1030	1063	SHN	-12	MSG	00071
1031	1237	LPN	378	SCOP31A	00426
1032	3404	STU	D.OUT	MSG	00072
1033	1410	LDN	8	SCOP31A	00427
1034	3401	STU	D.Z1	SCOP31A	00428
1035	3404	STU	D.Z4	SCOP31A	00429
1036	2000 1155	LDC	HUF-1	SCOP31A	00430
1040	3402	STU	D.Z2	SCOP31A	00431
1041	3004	LDU	D.OUT	MSG	00079
1042	1014	SHN	12	SCOP31A	00432
1043	3165	AOU	D.OUT+1	MSG	00083
1044	0200 0634	RJM	R.TFL	SCOP31A	00433
1046	0732	MJN	MS4	SCOP31A	00434
1047	6170 1156	CRU	HUF.D.PPONE	SCOP31A	00435
1051	1405	LDN	5	SCOP31A	00436
1052	3403	STU	D.Z3	SCOP31A	00437
1053	5500 1050	KAM	MS3A+1	SCOP31A	00438
1055	3002	AOU	D.Z2	SCOP31A	00439
1056	4002	LDI	D.Z2	SCOP31A	00440
1057	1071	SHN	-6	MSG	00088
1060	1700	SHN	608	MSG	00089
1061	0617	PJN	MS4	SCOP31A	00441
1062	4002	LDI	D.Z2	MSG	00091
1063	1277	LPN	778	MSG	00092
1064	1700	SHN	608	MSG	00093
1065	0613	PJN	MS4	MSG	00094
1066	3703	SOU	D.Z3	SCOP31A	00442
1067	0505	NJN	MS2	SCOP31A	00443
1070	3701	SOU	D.Z1	SCOP31A	00444
1071	0415	ZJN	MS5	SCOP31A	00445
1072	4002	LDI	D.Z2	SCOP31A	00446
1073	0453	ZJN	MS3A	SCOP31A	00447
1074	3005	AOU	O.OUT+1	MSG	00102

14-30

1075	1003	SHN	-12	MSG	00103
1076	3004	KAD	D. OUT	IF NOT, GET ADDRESS OF NEXT WORD	MSG 00104
1077		UJK	MS1		SCOP31A 00448
1100	2000 1142	LDC	MSG1	MESSAGE FORMAT ERROR	MSG 00113
1102	0200 0650	RJM	R.DFM	PROCESS DAYFILE MESSAGE	MSG 00114
1104	1413	MS12	LDN	M.ABORT	MSG 00115
1105	0322	UJN	MS7	ABORT CP	MSG 00116
1106	3052	MS5	LDD	D.PPIRB+2 IS THIS FOR DISPLAY ON SCOPE ONLY	MSG 00117
1107	0524	NJN	MS9	YES, JUMP	MSG 00118
1110	2000 1156	LDC	BUF		SCOP31A 00449
1112	0200 0650	RJM	R.DFM	OUTPUT DAYFILE MESSAGE	SCOP31A 00450
1114	3002	MS10	LDD	D.IN	MSG 00123
1115	0411	ZJN	MS10A	NO-RECALL	MSG 00124
1116	3044	AOD	D.BA+4	SET COMPLETEION BIT FOR CPC	MSG 00125
1117	3053	LDD	D.PPIRB+3		MSG 00126
1120	1237	LPN	378		SCOP31A 00451
1121	1014	SHN	12		SCOP31A 00452
1122	3154	ADD	D.PPIRB+4		MSG 00130
1123	0200 0634	RJM	R.TFL	RELOCATE ADDRESS	SCOP31A 00453
1125	6240	CWD	D.BA	TELL CPC	MSG 00131
1126	1412	MS10A	LDN	M.DPP	MSG 00132
1127	0200 0450	MS7	RJM	R.PROCES	MSG 00133
1131	0100 0100	LJM	R.IDLE	SEND REQUEST TO OUTPUT REGISTER	MSG 00134
1133	3051	MS9	LDD	D.PPIRB+1	SCOP31A 00454
1134	1207	LPN	L.CPNUM	EXIT TO IDLE LOOP	SCOP31A 00455
1135	1007	SHN	7	EXTRACT CP NUMBER	SCOP31A 00456
1136	1630	ADN	W.CPDM		MSG 00140
1137	6304 1156	CWM	BUF, D.Z4	WRITE MESSAGE FOR DISPLAY ONLY	SCOP31A 00457
1141	0352	UJN	MS10	EXIT	MSG 00142
1142	1505	MS01	DIS	MESSAGE-FORMAT-ERROR.	MSG 00143
1156	0000	RUF	DATA	0	MSG 00144
1157			END		MSG 00146

030274 UNUS-ED STORAGE 139 STATEMENTS 475 SYMBOLS

14-31

SYMBOLIC REFERENCE TABLE

BUF	0001156	001036,	001047,	001110,	001137.
CH.UMP	0000023				
CH.FNT	0000015				
CH.FST	0000014				
CH.LIB	0000016				
CH.RBT	0000017				
CP.ECSM	0002034				
CP.IDLE	0002000				
CP.JOBN	0002020				
CP.MOVE	0002000				
CP.SM	0002023				
C.CPAR	0000001				
C.CPECFL	0000004				
C.CPEF	0000001				
C.CPERT	0000004				
C.CPFL	0000004				
C.CPFP	0000003				
C.CPFST	0000002				
C.CPNCS	0000004				
C.CPNUM	0000001				
C.CPOUT	0000003				
C.CPPRI	0000000				
C.CPRA	0000003				
C.CPSM	0000002				
C.CPSTAI	0000000				
C.CPTIML	0000002				
C.DFBMS	0000001				
C.DIRCMA	0000001				
C.DIRFWA	0000000				
C.DIRPRU	0000004				
C.DIRPTR	0000000				
C.DIRKBA	0000002				
C.DIRKBN	0000003				
C.DIRUNT	0000001				
C.FCIB	0000000				
C.FCPNUM	0000003				
C.FCS	0000003				
C.FDC	0000002				
C.FEQP	0000000				
C.FFRBA	0000001				
C.FLBL	0000003				
C.FLOCK	0000003				
C.FLPRU	0000004				
C.FLRBEB	0000003				
C.FLRBWP	0000002				
C.FNAME	0000000				
C.FPDEV	0000002				
C.FPRI	0000004				
C.FSC	0000003				
C.FSDEV	0000001				
C.FTYPE	0000003				
C.PPFWA	0001000	000000			
C.PPSWA	0002000				
C.PPTWA	0003000				
C.PP4WA	0004000				
C.PP5WA	0005000				
C.PP7WA	0007000				

14-34

C.RBRA 0000003

C.RBRAU 0000000

C.RBR LAV 0000003

C.RHRTPA 0000000

C.RBRUNF 0000001

C.RBTAL 0000002

C.RBTFB 0000001

C.RBTLPK 0000000

C.RHTLRK 0000004

C.RHTPRU 0000003

C.RBTBKH 0000001

C.RBTWPL 0000000

C.RQSCS 0000000

C.RQSFS 0000002

C.RSTA 0000002

C.RSTRA 0000001

C.RSTU 0000004

C.RSTWP 0000003

C.RWPPCC 0000003

C.RWPPCF 0000000

C.RWPPPL 0000002

C.RWPPST 0000003

C.RWPPWC 0000004

C.RWPPWI 0000001

C.STAICP 0000002

C.STCPU 0000004

C.STEI 0000004

C.STFB 0000003

C.STU 0000003

C.SIPFW 0000002

C.SIPLW 0000004

C.SIPMS 0000001

C.SIPPRU 0000002

C.SIPRBA 0000000

C.SIPRBN 0000001

C.SIPWC 0000000

D.BA 0000040

D.CPAU 0000074

D.DTS 0000037

D.EST 0000032

D.FA 0000057

D.FF0 0000050

D.FF1 0000051

D.FF2 0000052

D.FF3 0000053

D.FF4 0000054

D.FF5 0000055

D.FF6 0000056

D.FF7 0000057

D.FIKST 0000060

D.FL 0000056

D.FNT 0000020

D.FR0 0000040

D.FR1 0000041

D.FR2 0000042

D.FR3 0000043

D.FR4 0000044

001015, 001022, 001025, 001116, 001125

14-37

SYMBOLIC REFERENCE TABLE

14-34

D.T6	0000016	
D.T7	0000017	
D.Z0	0000000	
D.Z1	0000001	001034, 001070
D.Z2	0000002	001040, 001055, 001056, 001062, 001072
D.Z3	0000003	001052, 001066
D.Z4	0000004	001035, 001137
D.Z5	0000005	
D.Z6	0000006	
D.Z7	0000007	
F.EKAK	0000002	
F.EKCP	0000004	
F.EKOD	0000006	
F.ERPCE	0000005	
F.EKPP	0000003	
F.EHTL	0000001	
LE.FNT	0000003	
L.CPNUM	0000007	001134
L.PPHDR	0000005	
MSG1	0001053	001021
MSG1	0001142	001100
MS1	0001041	001077
MS10	0001114	001141
MS10A	0001126	001115
MS12	0001104	
MS2	0001055	001067
MS3	0001046	001014
MS3A	0001047	001053, 001073
MS4	0001100	001046, 001061, 001065
MS5	0001106	001071
MS7	0001127	001105
MS9	0001133	001107
M.ABORT	0000013	001104
M.AEWP	0000033	
M.CCPA	0000035	
M.CDF	0000011	
M.CPUST	0000036	
M.DCH	0000003	
M.DCP	0000016	
M.DEQP	0000023	
M.DFM	0000001	
M.DPP	0000012	001126
M.DTAPE	0000032	
M.EREQS	0000034	
M.NTIME	0000014	
M.OPDROP	0000030	
M.PAUSE	0000017	
M.PPTIME	0000004	
M.RCH	0000002	
M.RCLCP	0000021	
M.RCP	0000015	
M.KEM	0000025	
M.REQP	0000022	
M.RPJ	0000037	
M.RPP	0000020	
M.RPHI	0000024	
M.RSTOR	0000010	

14-75

SYMBOLIC REFERENCE TABLE

M.HTAPE	0000031
M.SEF	0000030
M.STEP	0000005
N.CP	0000007
OV.ATS	0012423
OV.CHK	0031013
OV.CIO	0031117
OV.CKP	0031320
OV.CLL	0031414
OV.CLO	0031417
OV.CLS	0031423
OV.CPL	0032014
OV.CIS	0032423
OV.DIS	0041123
OV.DMP	0041520
OV.EAU	0053025
OV.LBC	0140203
OV.LDK	0140422
OV.LUC	0141703
OV.LUD	0141704
OV.MOI	0150411
OV.MEM	0150515
OV.MSG	0152307
OV.OPE	0172005
OV.PBC	0200203
OV.RBR	0220222
OV.REQ	0220521
OV.HFL	0220614
OV.RST	0222324
OV.SRB	0232202
OV.WBR	0270222
OV.IAJ	0340112
OV.IBJ	0340212
OV.IBT	0340224
OV.ICU	0340317
OV.ICY	0340331
OV.IDF	0340406
OV.IEE	0340505
OV.IEI	0340511
OV.IIX	0341130
OV.IIJ	0341412
OV.III	0341424
OV.IIM	0341522
OV.IIM	0341527
OV.IIU	0341704
OV.IUT	0341724
OV.IPO	0342017
OV.IRI	0342211
OV.IRO	0342217
OV.IRU	0342221
OV.ISP	0342320
OV.ITU	0342404
OV.IXP	0343020
OV.2BP	0350220
OV.2BT	0350224
OV.2CA	0350301
OV.2CF	0350306

OV.2CJ 0350312
OV.2DF 0350406
OV.2EF 0350506
OV.2ES 0350523
OV.2LA 0351401
OV.2LB 0351402
OV.2LE 0351405
OV.2LP 0351420
OV.2PC 0352003
OV.2RC 0352203
OV.2RQ 0352221
OV.2RT 0352224
OV.2TB 0352402
OV.2TF 0352406
OV.2TJ 0352412
OV.2TK 0352422
OV.2TS 0352423
OV.2TW 0352427
OV.2WT 0352724
OV.30T 0361724
OV.4LB 0371402
OV.5DT 0400424
OV.7TP 0422420
O.BPRU 0000016
O.RCHN 0000017
O.RCMPR 0000002
O.RDP 0000010
O.RDPNP 0000011
O.RDSK 0000001
O.HEAD 0000000
O.SKB 0000013
O.SKF 0000012
O.WRP 0000014
O.WRPR 0000015
O.WRT 0000004
O.WRTR 0000005
P.CAT 0000011
P.CST2 0000015
P.CST3 0000016
P.CST4 0000017
P.DFB 0000003
P.EST 0000005
P.FNT 0000004
P.INS 0000007
P.LIB 0000001
P.RBR 0000002
P.RBT 0000002
P.RQS 0000013
P.SPI 0000012
P.ZERO 0000000
R.CPFL 0000627
R.CPRA 0000631
R.DCH 0000714
R.DFM 0000650
R.EREQS 0000300
R.ERQ 0000334
R.IDLE 0000100

001102, 001112

001131

14-37

SYMBOLIC REFERENCE TABLE

R.MTR	0000450
R.OVL	0000124
R.OVLJ	0000111
R.PAUSE	0000430
R.PROCES	0000450
R.RCH	0000704
R.READP	0000460
R.RWP	0000505
R.RWPP	0000530
R.STB	0000620
R.STBMSK	0000611
R.STEP	0000734
R.TFL	0000634
R.WAIT	0000410
R.WHITEP	0000470
S.DIRPR	0000010
S.DIRPT	0000004
S.FNTEQP	0000006
S.FNTLK	0000005
S.FNTTYP	0000003
S.RBKUNT	0000006
S.RBTRBR	0000003
S.RBTRKL	0000007
S.RBTRNU	0000006
S.STF	0000006
S.STFA	0000000
S.STFEOP	0000004
S.STFETP	0000002
S.STFNTP	0000003
S.STFPM1	0000005
S.STFKEL	0000004
T.C10LE	0000023
T.CLK	0000030
T.CPA1	0000200
T.CPA2	0000400
T.CPA3	0000600
T.CPA4	0001000
T.CPA5	0001200
T.CPA6	0001400
T.CPA7	0001600
T.CPS	0000040
T.CPT1	0000056
T.CPZ	0000020
T.CP8	0002000
T.CST	0000014
T.DATE	0000031
T.JDATE	0000027
T.MON	0000021
T.MSC	0000040
T.MSP	0000037
T.P10LE	0000024
T.PPC1	0000060
T.PPC2	0000070
T.PPC3	0000100
T.PPC4	0000110
T.PPC5	0000120
T.PPC6	0000130

001127

001012, 001044, 001123

14-38

T.PPC7 0000140
T.PPC8 0000150
T.PPC9 0000160
T.PPR 0000025
T.PPS0 0000041
T.PPS1 0000042
T.PPS2 0000043
T.PPS3 0000044
T.PPS4 0000045
T.PPS5 0000046
T.PPS6 0000047
T.PPS7 0000050
T.PPS8 0000051
T.PPS9 0000052
T.SLAB1 0000031
T.SLAB2 0000032
T.SLAB3 0000033
T.SLAB4 0000034
T.SLAB5 0000035
T.SLAB6 0000036
T.STATCP 0000056
T.STU 0000022
T.TMP 0000055
T.UAS 0000056
W.CPAR 0000157
W.CPCAF 0000051
W.CPCAL 0000150
W.CPDFM 0000030
W.CPECS 0000022
W.CPEF 0000020
W.CPENC 0000155
W.CPERT 0000040
W.CPFL 0000020
W.CPJNAM 0000021
W.CPOAE 0000153
W.CPOUT 0000153
W.CPPRI 0000022
W.CPRCL 0000025
W.CPRES1 0000156
W.CPRES2 0000160
W.CPSM 0000020
W.CPSTAT 0000020
W.CPTBUF 0000041
W.CPTBUL 0000050
W.CPTIME 0000023
W.CPTIML 0000022
W.CPVRNO 0000154
W.EWP 0000027
W.FSTCC 0000151
W.FSTNR 0000152
W.FTYPE 0000000
W.PPIR 0000000
W.PPMES1 0000002
W.PPMES2 0000003
W.PPMES3 0000004
W.PPMES4 0000005
W.PPMES5 0000006

001136

14-39

SYMBOLIC REFERENCE TABLE

W.PPMES6	0000007
W.PP0K	0000001
W.PPTIME	0000024
W.RBKLA	0000001
W.RBKTPA	0000000
W.RHRUNT	0000000
W.RWPPCW	0000002
W.SSW	0000026
W.STCPU	0000000
W.STEI	0000000
W.STFB	0000001
W.STO	0000000
W.STPFW	0000001
W.STPLW	0000001
W.STPMS	0000001
W.STPPRU	0000000
W.STPRBA	0000000
W.STPRBN	0000000
W.STPWC	0000001

02/08/68 SCOPE V3.1 6500-00 131K 11/21/67

01.25.36.SCPD01F

01.25.36.SCPDOC,P17,1500.CM60000.

01.25.52.REQUEST(OLUPL)

01.25.52. (S3 ASSIGNED)

01.25.52.REWIND(OLUPL)

01.25.53.UPDATE(Q)

01.26.13. UPDATING FINISHED

01.26.13.COMPASS(I=COMPILE,S=SCPTXT)

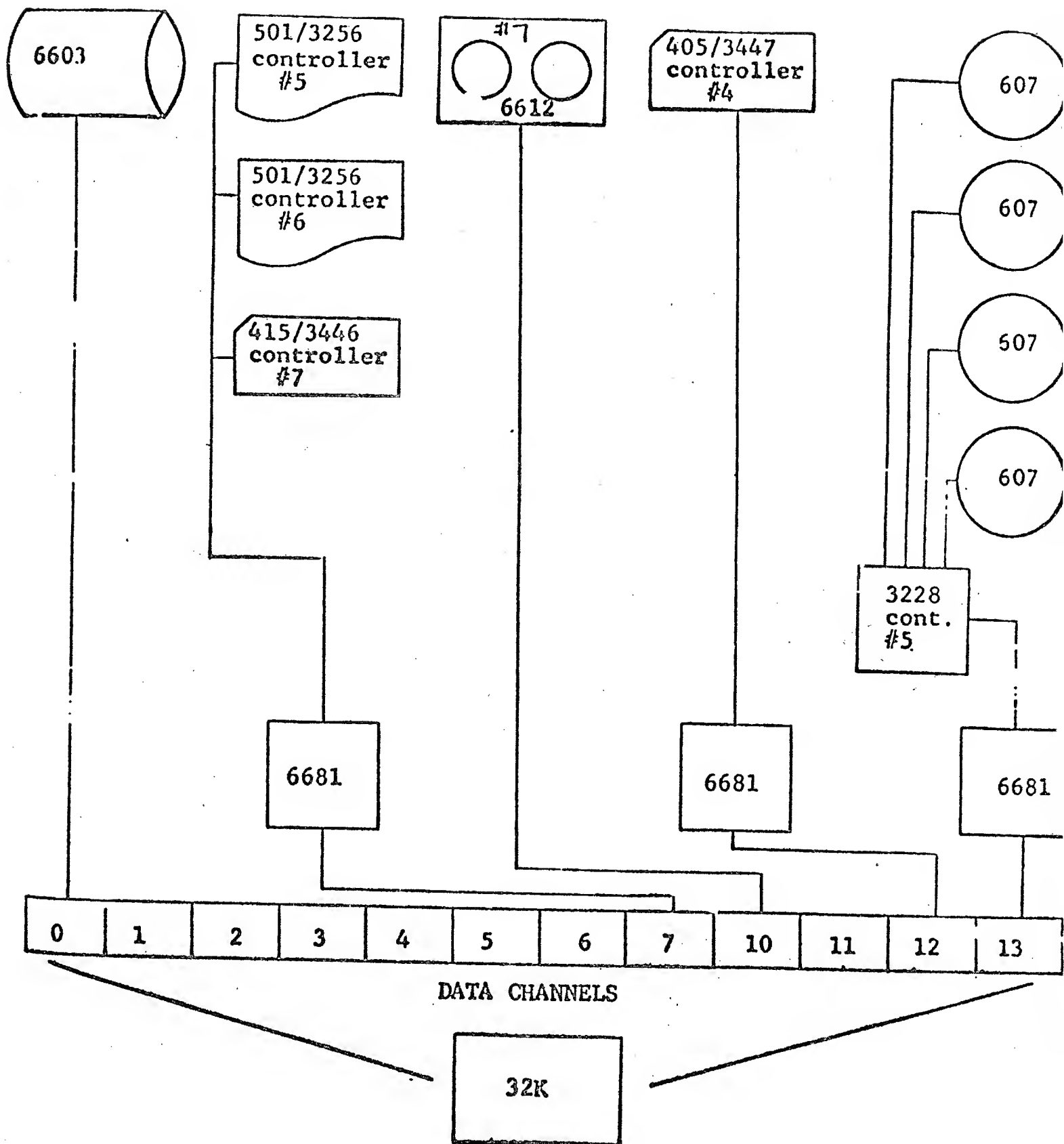
01.26.23.CP 013.068 SEC.

01.26.23.PP 041.638 SEC.

11-11

UNIVERSITY OF ARIZONA

6400



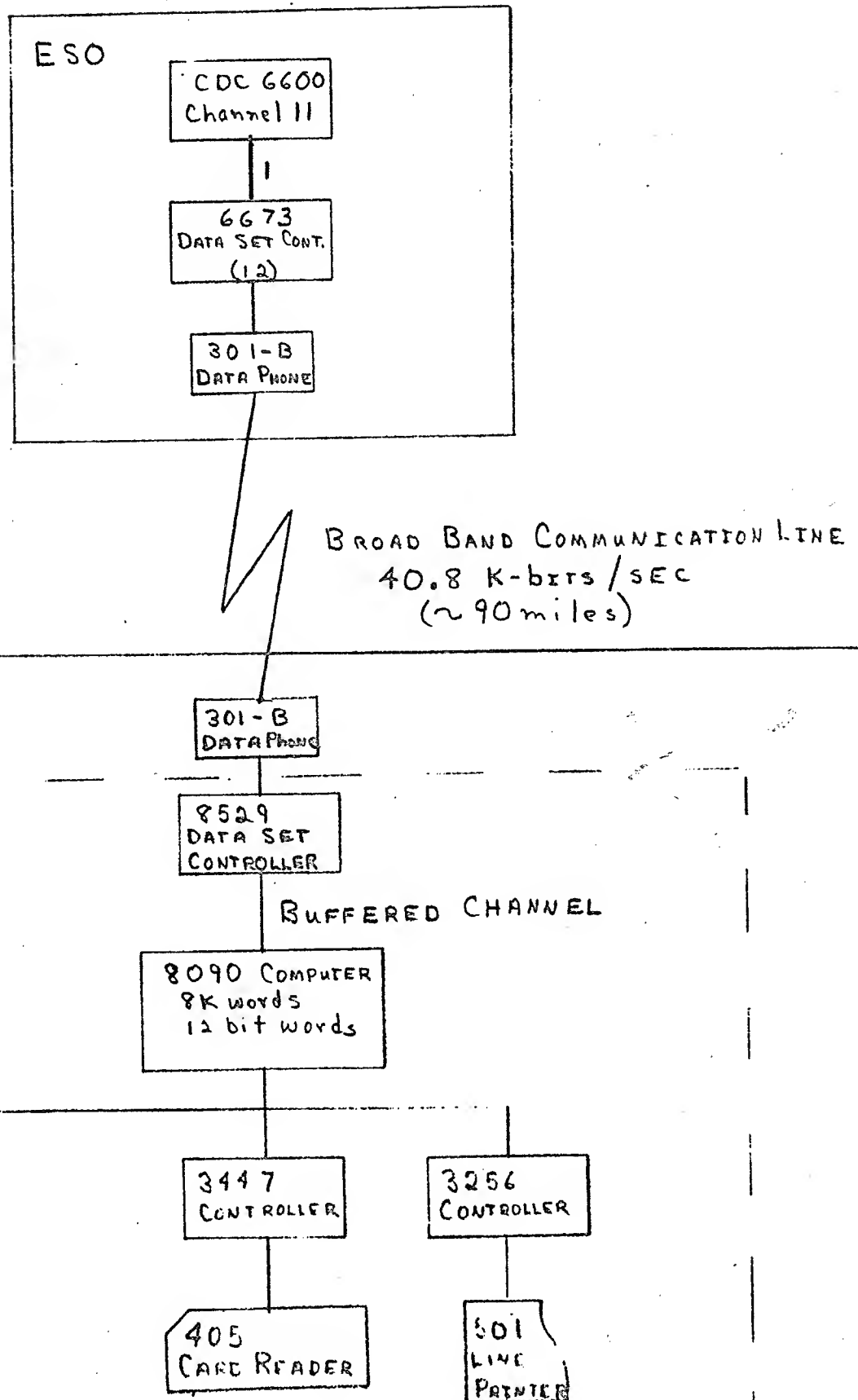
AEROSPACE CONFIGURATION

SBO REMOTE

3/1/68

RKV

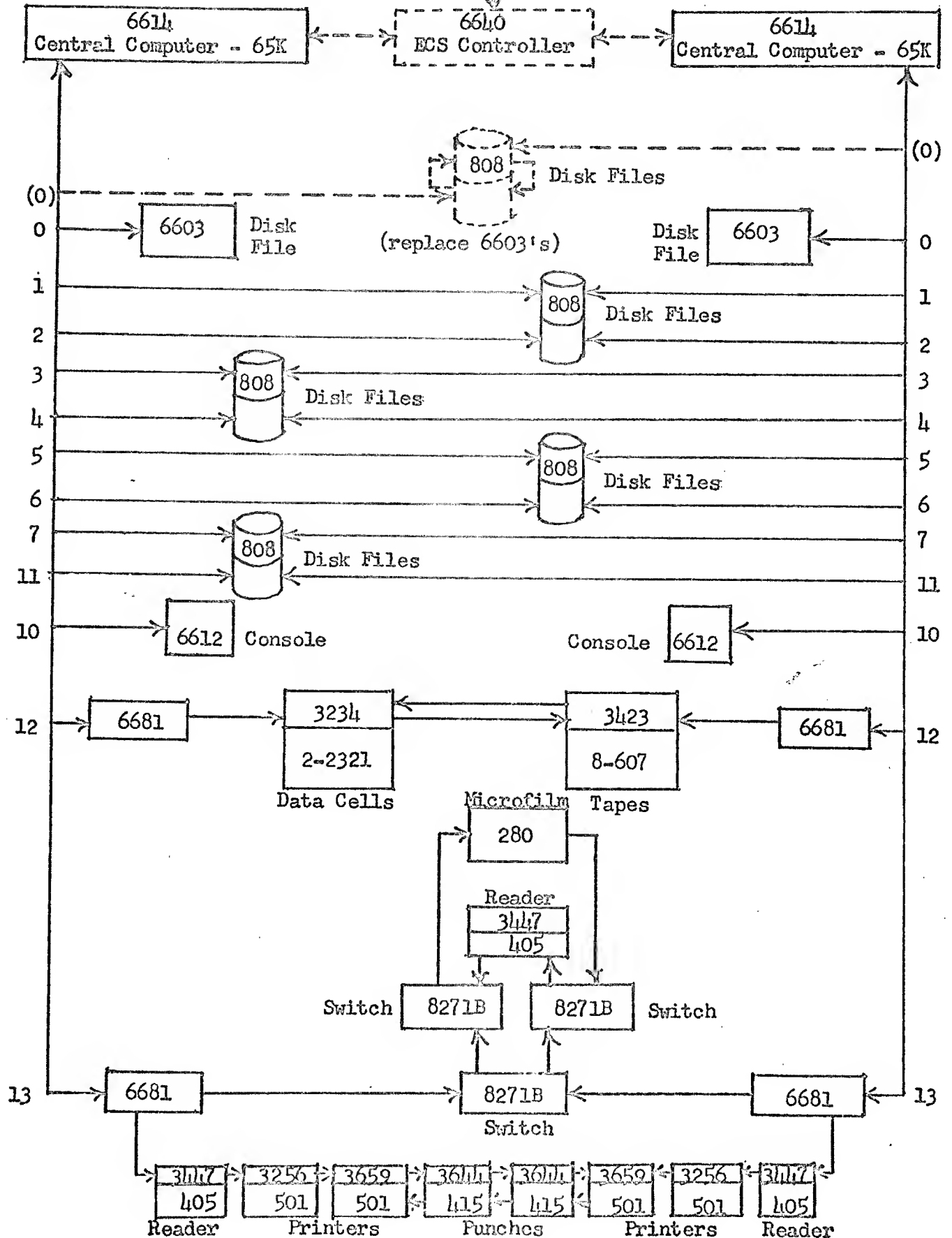
2 of 2



BETTS ATOMIC POWER LABORATORY

Attachment

6635 - EXTENDED CORE STORAGE
1000K



AEROSPACE CONFIGURATION

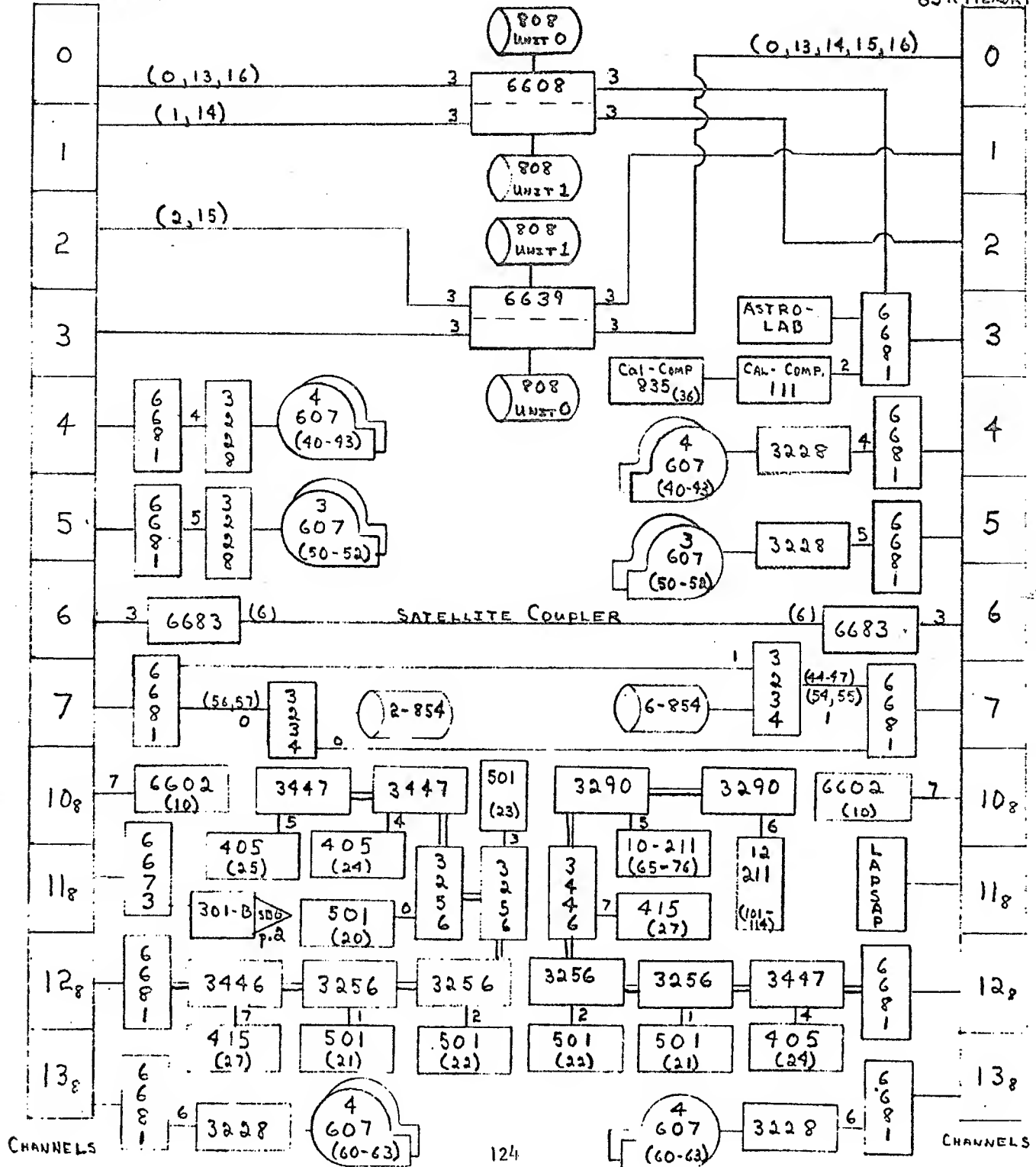
3/1/68
RKV
1 of 2

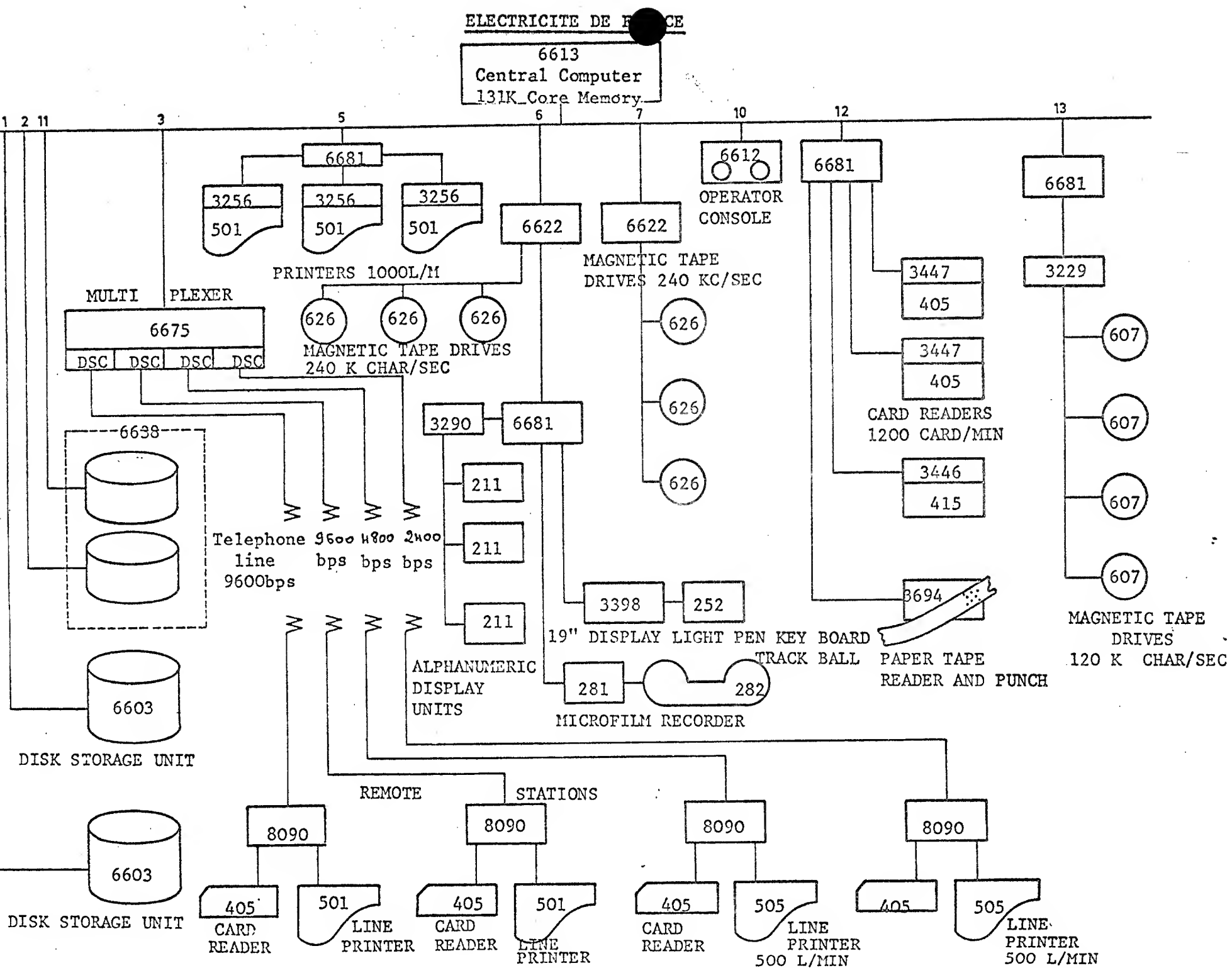
CDC
6400

65K MEMORY

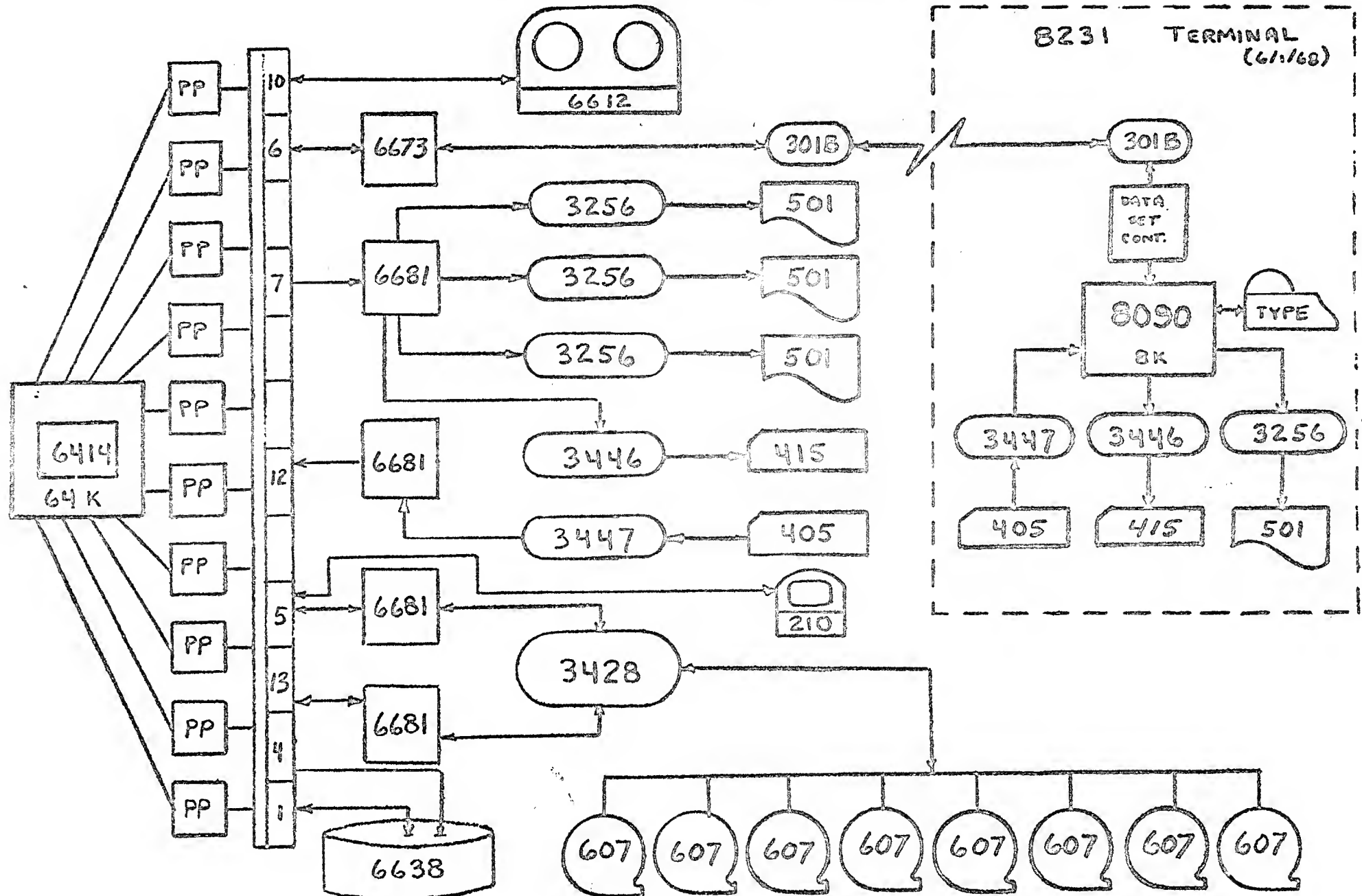
CDC
6600

131K MEMORY

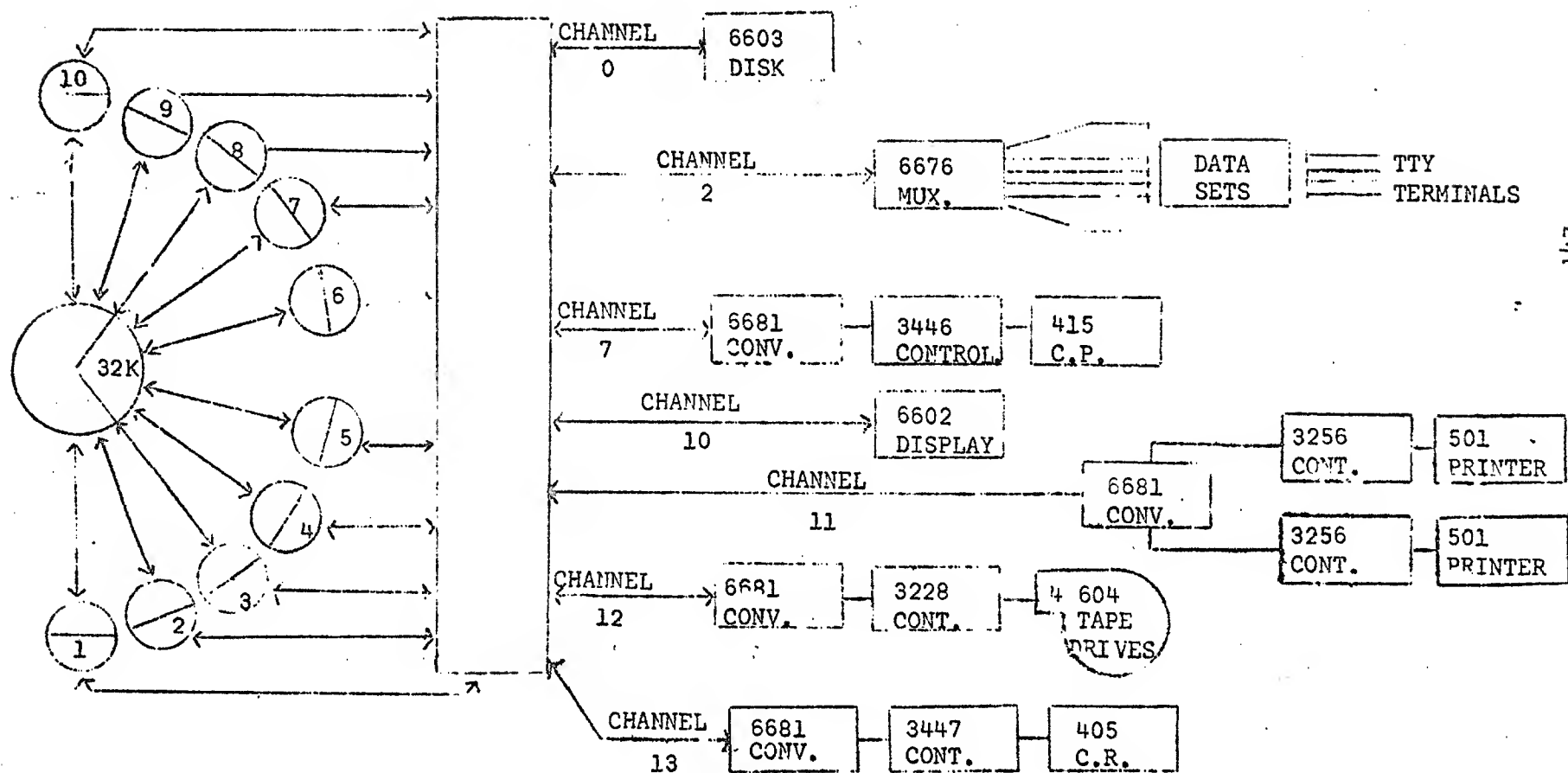




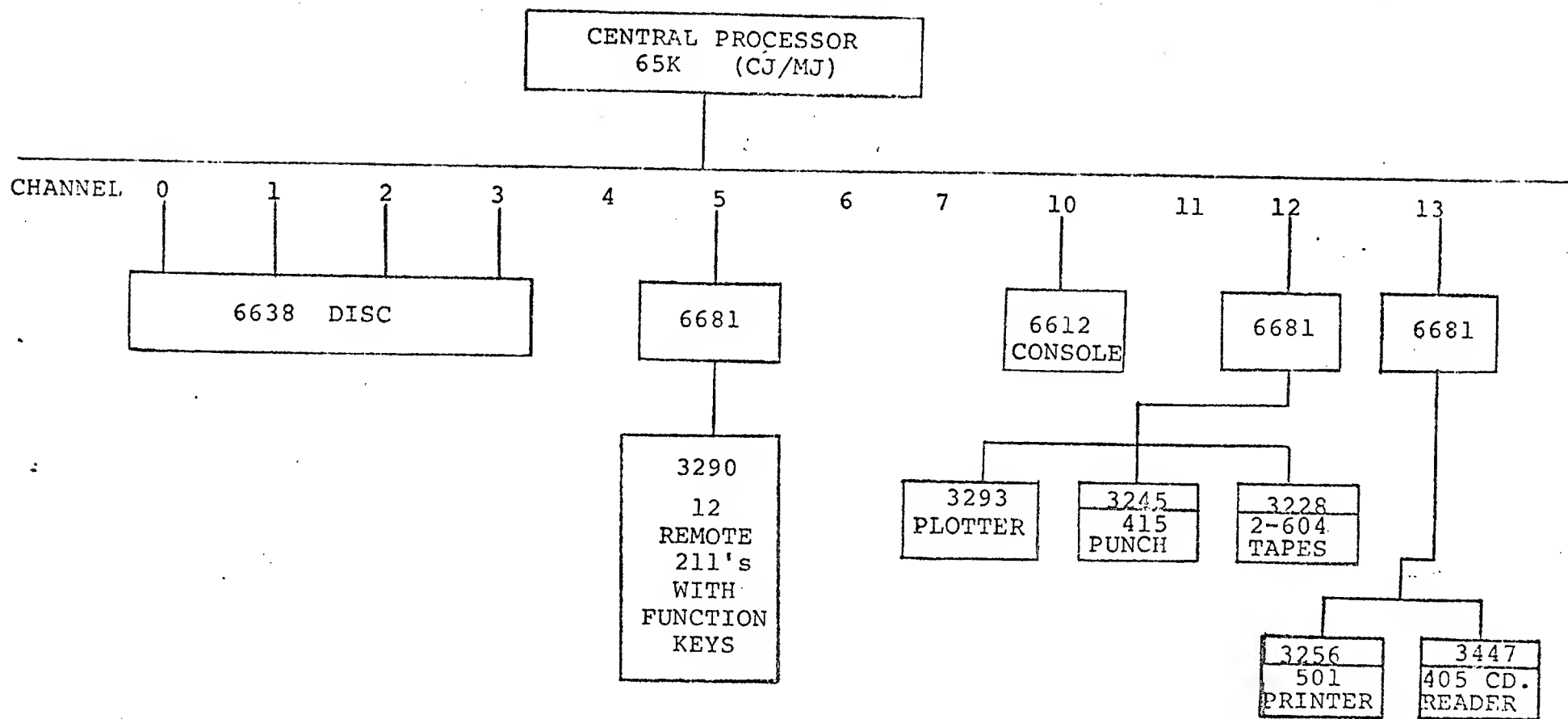
CDC 6400 Scientific Computer System

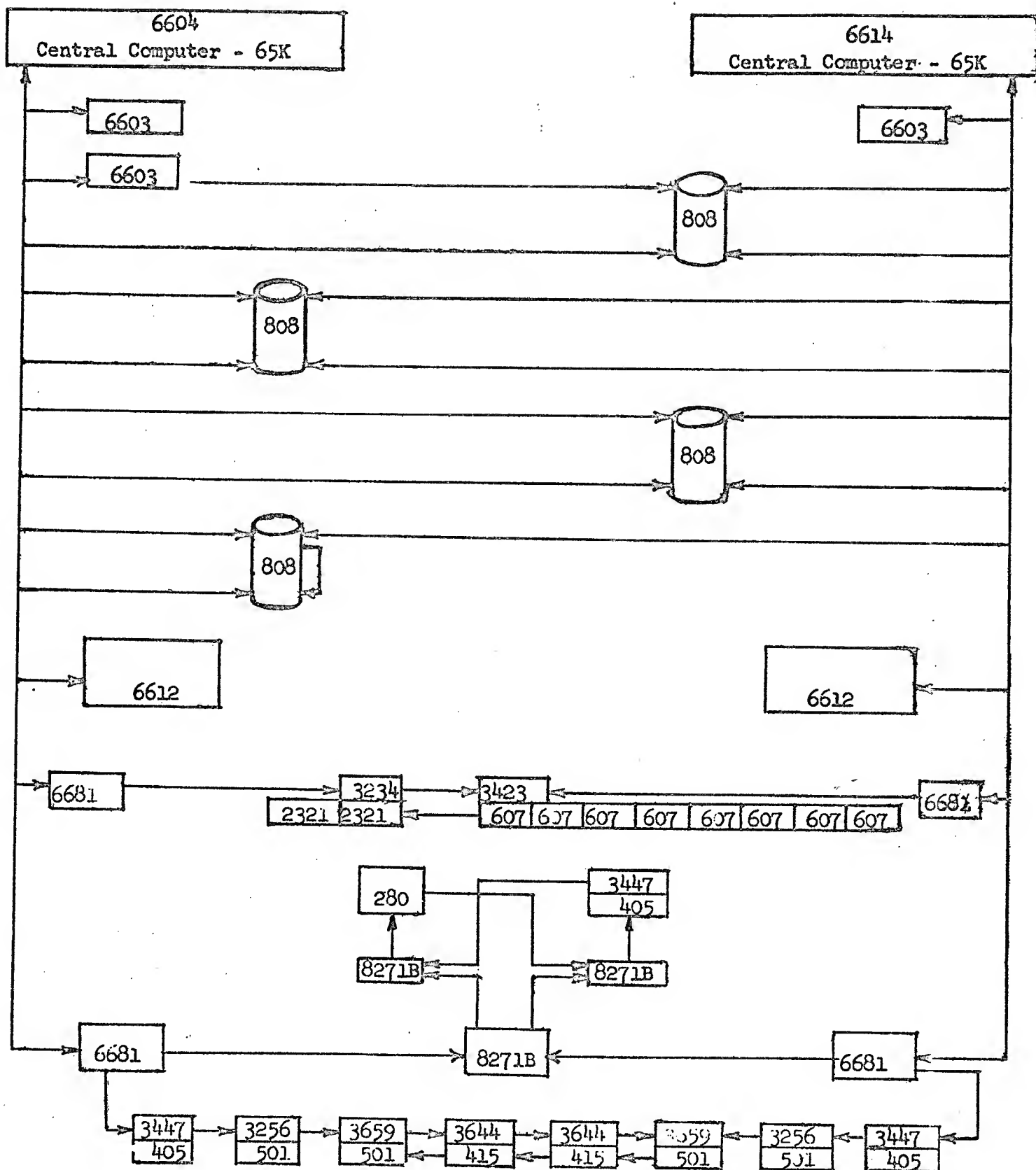


FSU CONFIGURATION



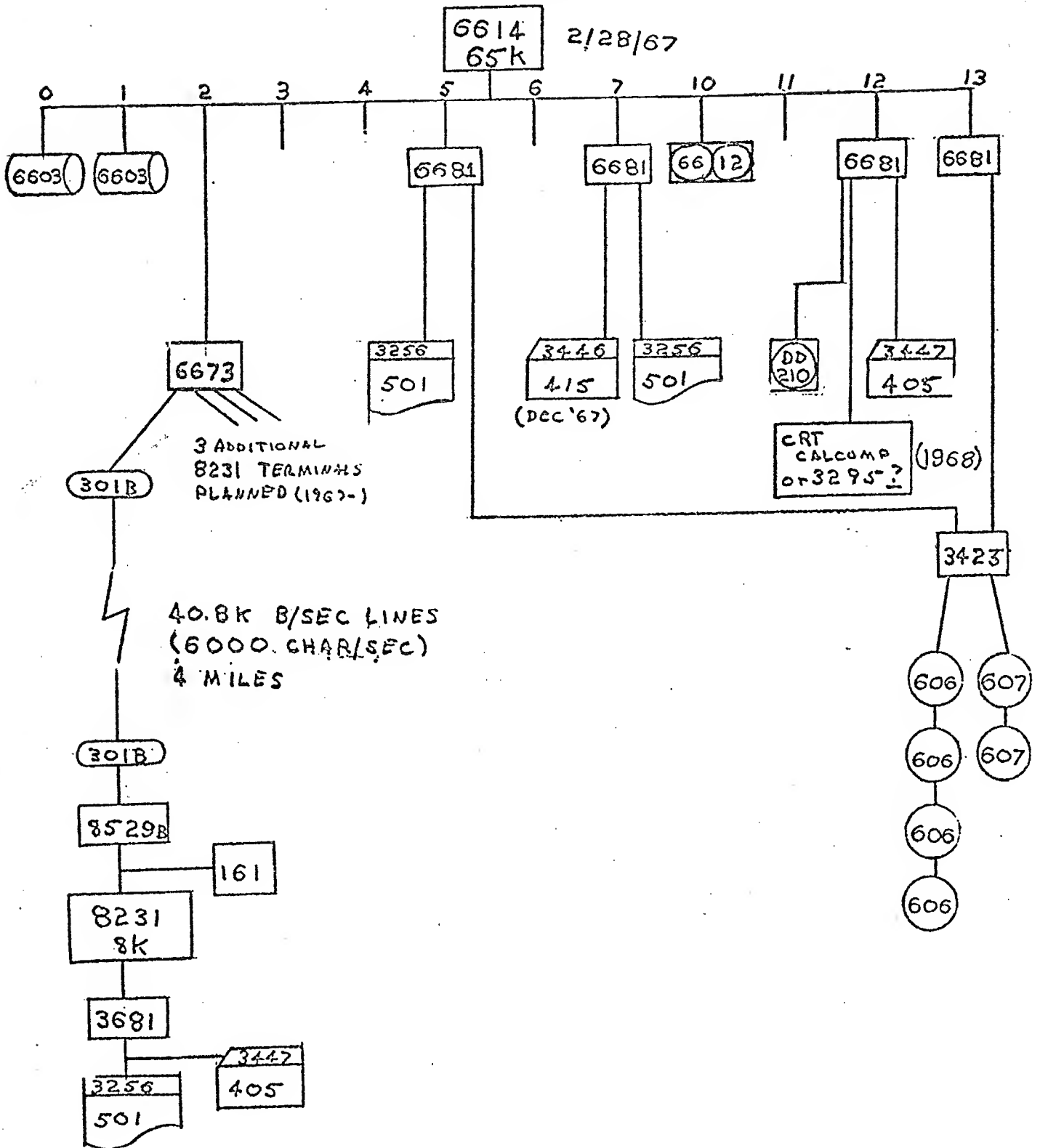
IDA COMPUTER FACILITY

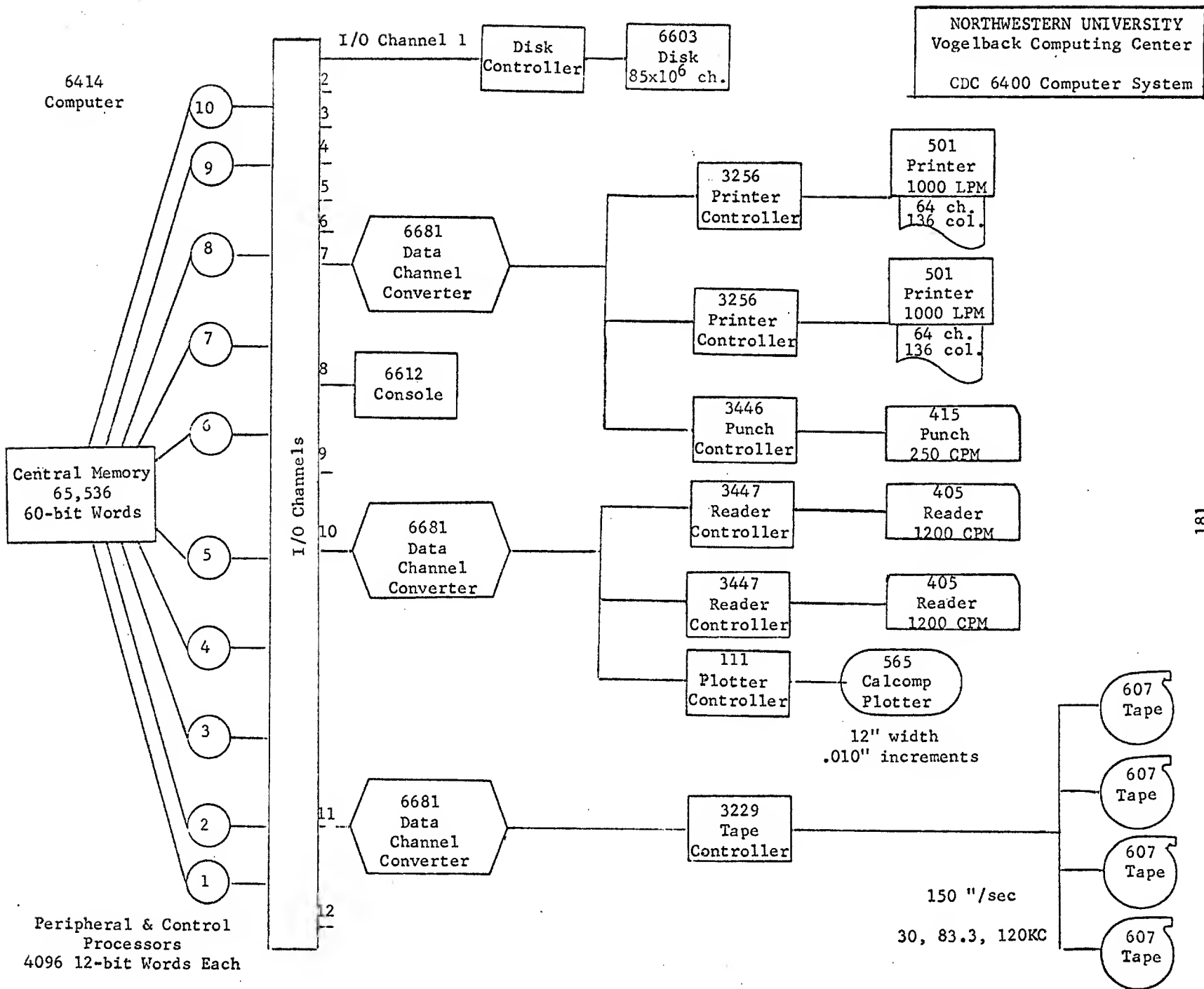




KNOLLS ATOMIC POWER LABORATORY CURRENT CONFIGURATION

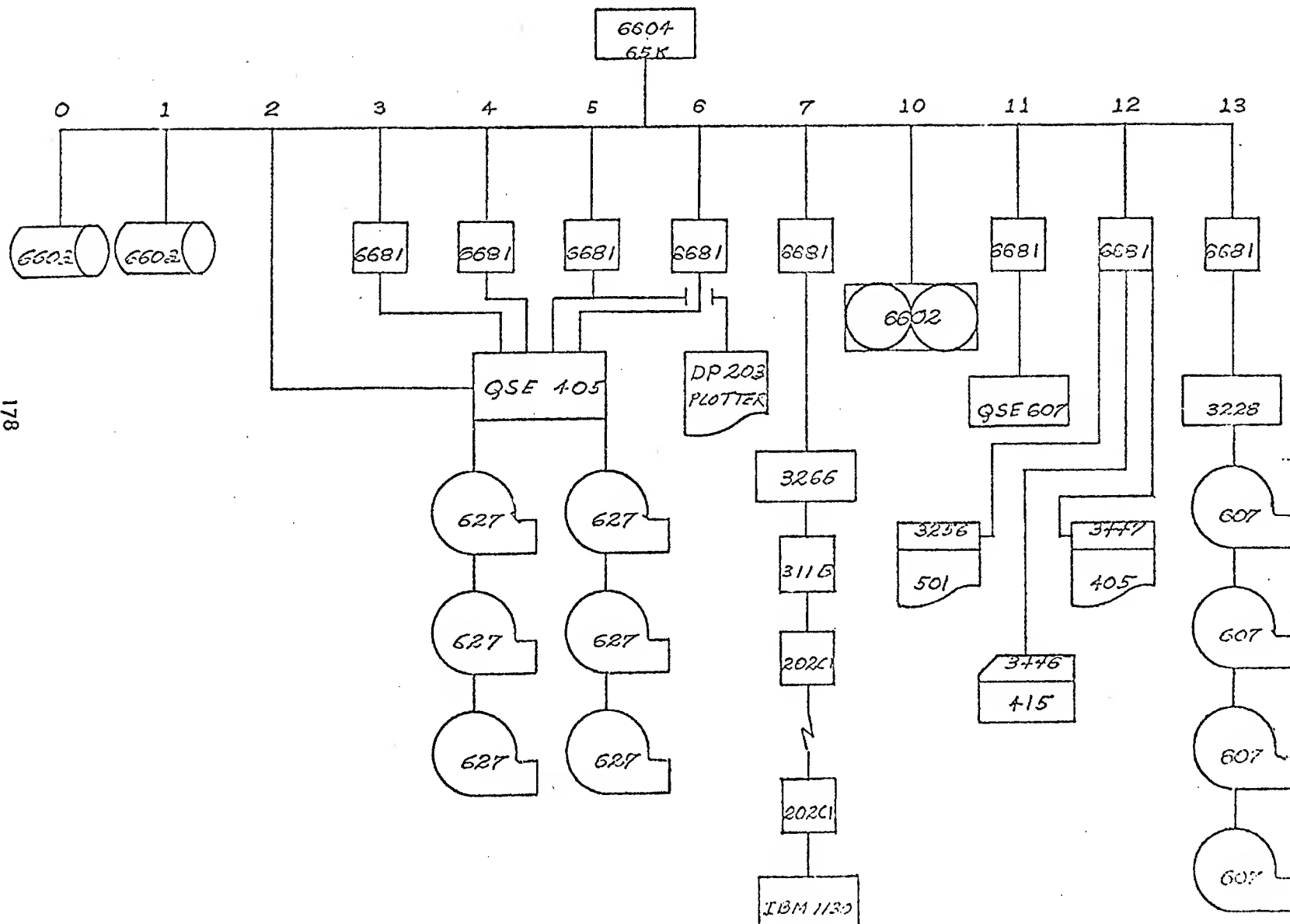
UNIVERSITY OF MINNESOTA
MINNEAPOLIS, MINNESOTA



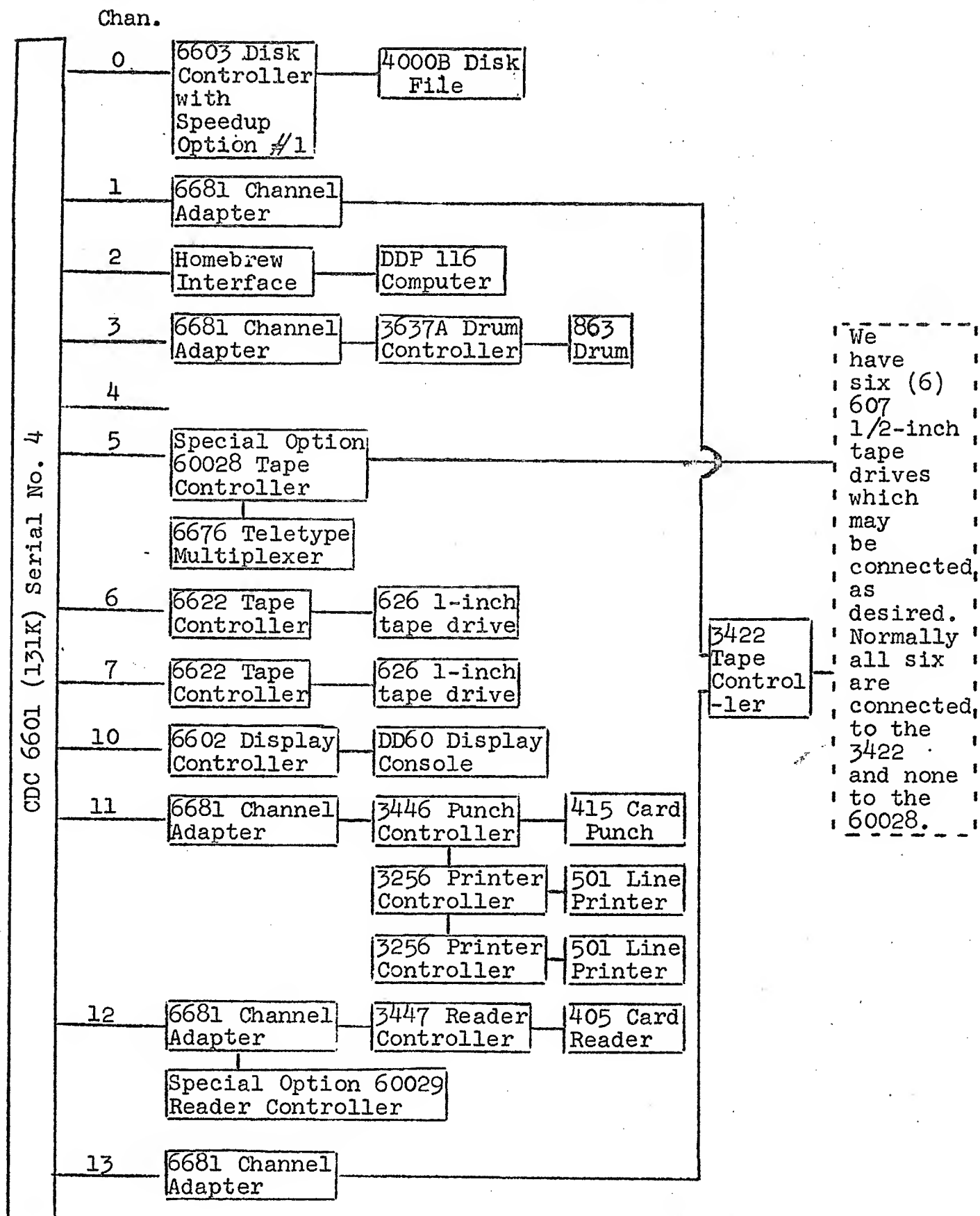


MOBIL OIL CORPORATION (MOGS)

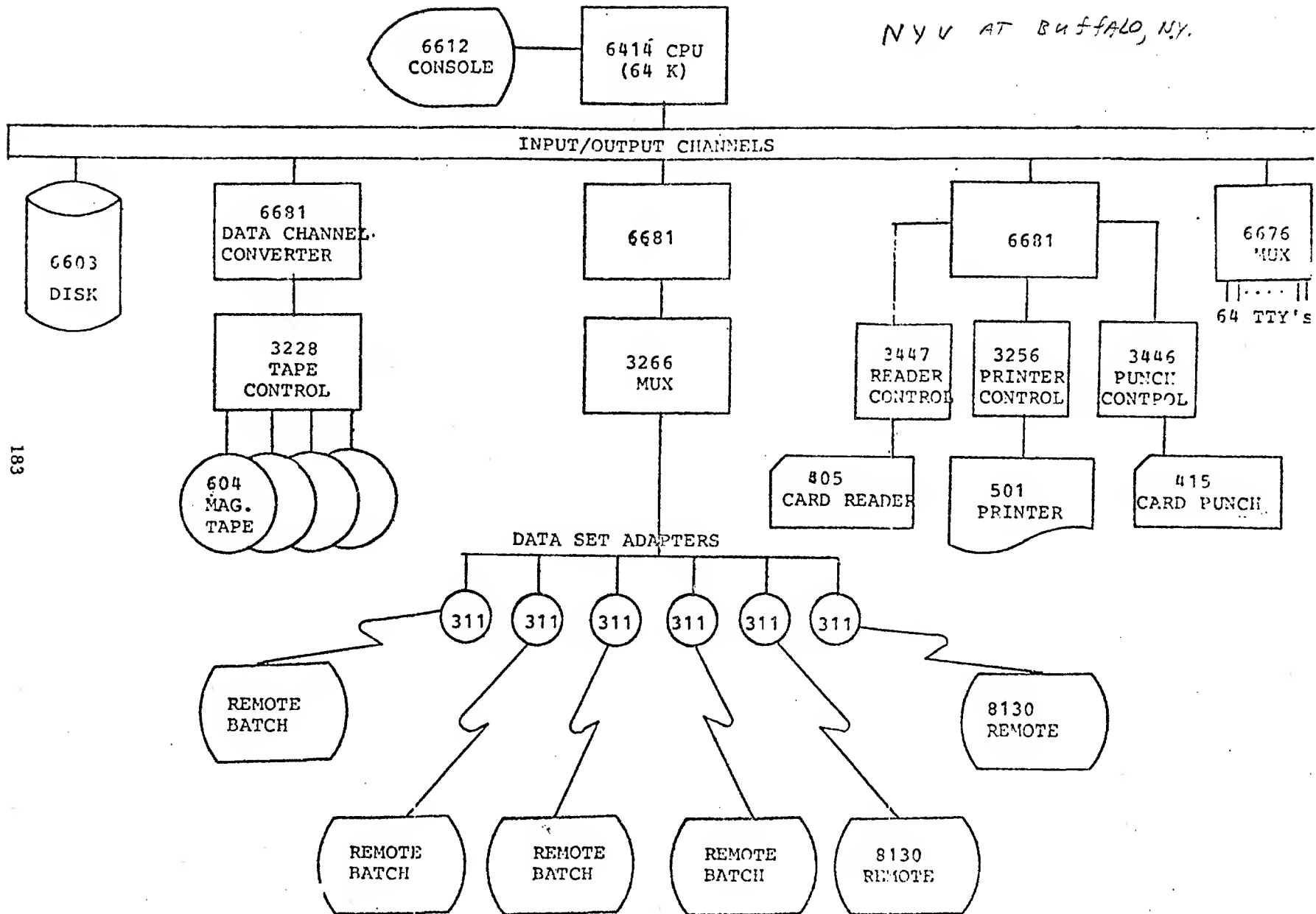
Geophysical Services Center
1930 Proctor Street
Dallas, Texas



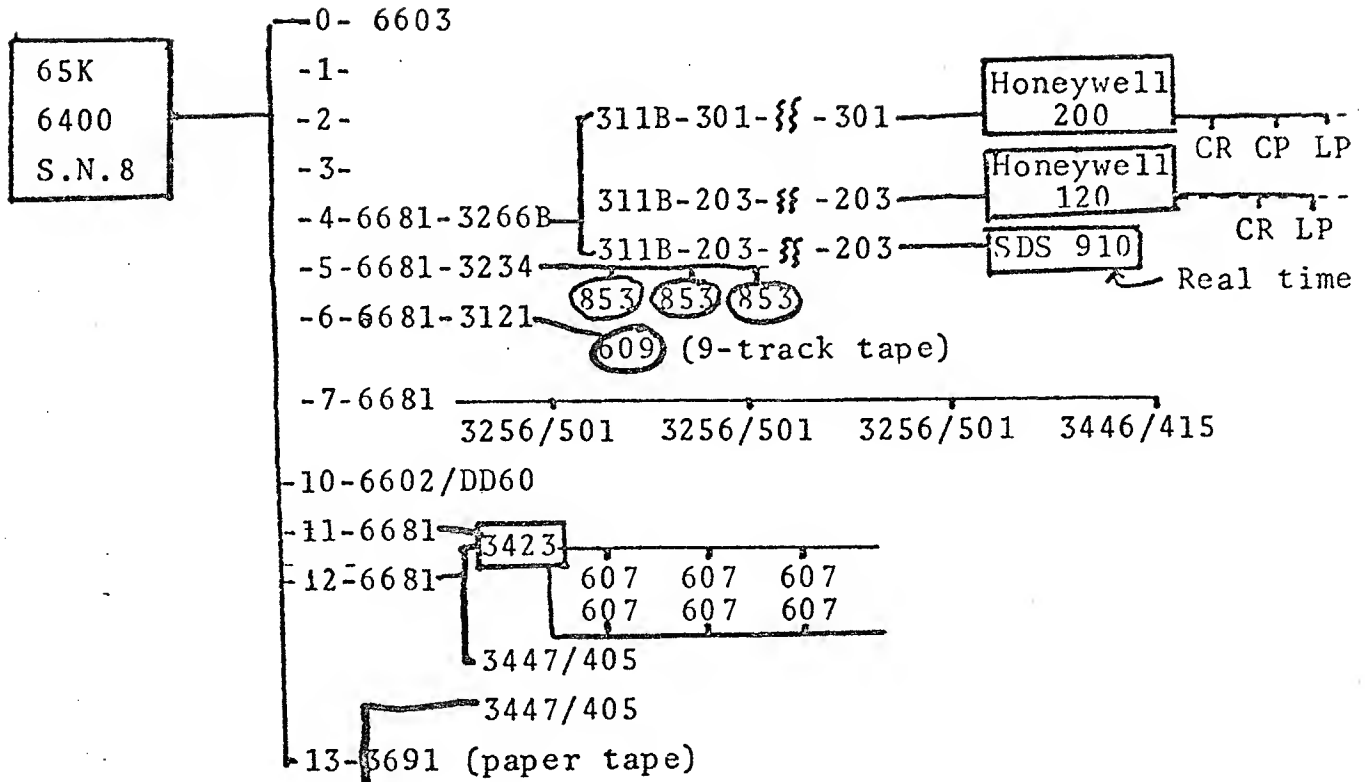
NYU -- Courant Institute CDC 6600 Channel Configuration (3/5/68)



NYU AT Buffalo, NY.



SAO CONFIGURATION



CDC INSTALLATION

